## Certifying Automated Reasoning

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## Automated reasoning

Significant progress in last couple of decades on combinatorial solvers

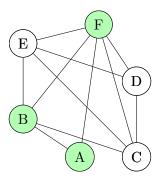
- Boolean satisfiability (SAT) & modulo theories (SMT), solving and optimization [Biere, Heule, van Maaren, and Walsh, 2021]
- Constraint programming [Rossi, van Beek, and Walsh, 2006]
- Pseudo-boolean (0-1 integer linear programming) [Elffers and Nordström, 2020].

scheduling learning DAGs kidney matching cancer treatment allocation of education hardware and software verification bounded model checking allocation of work air traffic control

## Example problem

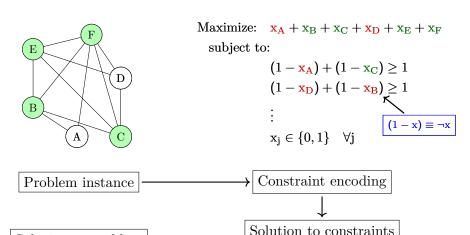
Maximum Clique

Decision problem: Is there clique of size 3? yes Optimization problem: What is the size of the largest clique?



## Automated Reasoning

for solving maximum clique

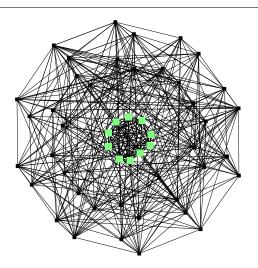


Solution to problem |

Green=1, Red=0

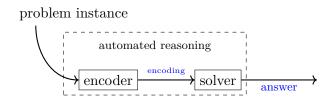
## The main question of the day

Can we trust the answer?



## Automated reasoning

in general



3 main approaches toward trustworthiness:

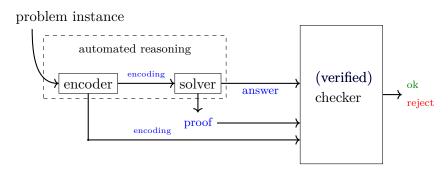
testing

formal verification

proof logging

## **Proof Logging**

[Järvisalo, Heule, and Biere, 2012; Wetzler, Heule, and Jr., 2014; Heule, 2021; van Doornmalen, Eifler, Gleixner, and Hojny, 2023; Bogaerts, Gocht, McCreesh, and Nordström, 2022]



#### Desiderata of proof format

- powerful
- simple

## Redundance-based proofs

#### Concrete Constraints

Propositional Logic, SAT, MaxSAT

- Instance:
  - ▶ Set of clauses, (CNF formula)
  - ▶ a linear objective function cost
- Find assignment  $\tau$  that:
  - satisfies all clauses and
  - minimizes cost

$$\tau(y) = \tau(b_1) = \tau(b_3) = 1$$
  
 $\tau(x) = \tau(b_2) = 0$ 

$$F = \{(b_1 \lor \mathbf{x}), (\neg \mathbf{x} \lor \mathbf{b_2}), (\mathbf{b_2} \lor \mathbf{y}), (\neg \mathbf{y}, \mathbf{b_3})\}$$

$$cost \equiv 2b_1 + 4b_2 + b_3$$

$$cost(\tau) = 3$$

## Clause Redundancy

[Järvisalo, Heule, and Biere, 2012; Heule, Kiesl, and Biere, 2020; Ihalainen, Berg, and Järvisalo, 2022]

#### Definition

Clause C is redundant for formula F and objective cost if

$$minimum-cost(F) = minimum-cost(F \land C)$$

(wrt. cost)

equisatisfiability a special case

#### Example:

$$(x \vee b_1) \wedge (\neg x \vee b_2)$$

 $(\neg b_2)$  is redundant

$$cost = b_1 + 2b_2$$

 $(\neg b_1)$  is **not** redundant

## Characterizing redundancy

[Heule, Kiesl, and Biere, 2020; Ihalainen, Berg, and Järvisalo, 2022]

#### (informal) Theorem

C redundant for F and cost iff there exists a set of literals  $L_C$  that fixes any solution  $\tau$  of F that falsifies C without increasing its cost.

#### Example

$$F = (x \vee b_1) \wedge (\neg x \vee b_2 \vee b_1) \wedge F' \xleftarrow{F' \text{ does not}} \\ cost = b_1 + 2b_2 \\ C = (\neg b_2) \text{ is redundant} \\ L_C = \{\neg b_2, b_1\}$$

$$F' \xleftarrow{F' \text{ does not}} \\ contain b_1, b_2, \neg b_1 \text{ or } \neg b_2$$

$$then assigning b_1 = 1, b_2 = 0 \text{ and} \\ the rest according to } \\ then assigning b_1 = 1, b_2 = 0 \text{ and} \\ the rest according to } \\ then assigning b_1 = 1, b_2 = 0 \text{ and} \\ the rest according to } \\ then assigning b_1 = 1, b_2 = 0 \text{ and} \\ then assigning b_1 = 1, b_2 = 0 \text{ and} \\ the rest according to } \\ then assigning b_1 = 1, b_2 = 0 \text{ and} \\ the rest according to } \\ then assigning b_1 = 1, b_2 = 0 \text{ and} \\ the rest according to } \\ then assigning b_1 = 1, b_2 = 0 \text{ and} \\ the rest according to } \\ then assigning b_1 = 1, b_2 = 0 \text{ and} \\ the rest according to } \\ then assigning b_1 = 1, b_2 = 0 \text{ and} \\ the rest according to } \\ then assigning b_1 = 1, b_2 = 0 \text{ and} \\ the rest according to } \\ then assigning b_1 = 1, b_2 = 0 \text{ and} \\ the rest according to } \\ then assigning b_1 = 1, b_2 = 0 \text{ and} \\ the rest according to } \\ then according to } \\$$

Note: adding redundant clauses might change the set of solutions

## A (very simplified) redundancy-based proof

e.g. [Heule, Kiesl, and Biere, 2020; Bogaerts, Gocht, McCreesh, and Nordström, 2022]

A proof for  $F = \{C_1, \dots, C_n\}$  and cost is a sequence:

$$C_1, C_2, \dots, C_n, C_{n+1}, \dots$$
 = empty clause

s.t. each  $C_{n+t}$  is either:

- redundant wrt.  $C_1 \wedge \ldots \wedge C_{n+t-1}$ , or
- cost < cost $(\tau)$  for a solution  $\tau$  of  $C_1 \wedge \ldots \wedge C_{n+t-1}$ .

The proof establishes:

- optimality if  $C_{n+t} = cost < cost(\tau)$  for some t
- infeasibility of constraints, otherwise

redundancy-based proof systems are strong need to be careful with deletion

## Redunancy for simulating solver reasoning

#### Subsumed Literal Elimination

[Berg, Saikko, and Järvisalo, 2016; Korhonen, Berg, Saikko, and Järvisalo, 2017]

#### Assume:

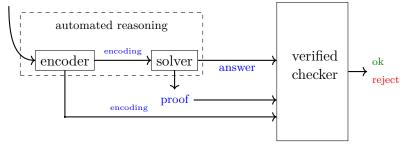
- i)  $b_2$  appears at least in the same clauses as  $b_1$ .
- ii) the coefficient of  $b_2$  in cost is at most the coefficient of  $b_1$ . Then fix  $b_2 = 0$  and simplify.

## Redundancy Reasoning $(x \vee b_1) \wedge (\neg x \vee b_2 \vee b_1)$ $cost = b_1 + 2b_2$ $(x \lor b_1) \land (\neg x \lor b_2 \lor b_1)$ add $(\neg b_2)$ $(x \vee b_1) \wedge (\neg x \vee b_1)$ add $(b_1)$ $(b_1)$ remove $(\neg x \lor b_1 \lor b_2)$

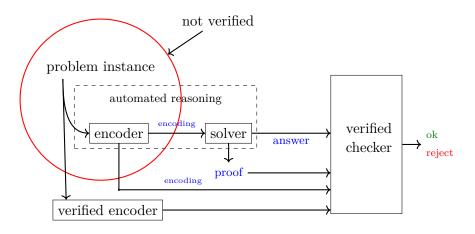
What do we need to trust?

## Recap: Certified Automated Reasoning

#### problem instance



## What about the encoding?



• problem-specific verified encoder can prove the right properties of the encoding

## What are these right properties?

[Gocht, McCreesh, Myreen, Nordström, Oertel, and Tan, 2024; Ihalainen, Oertel, Tan, Berg, Järvisalo, Myreen, and Nordström, 2024]

$$\begin{array}{l} \text{is\_clique } vs \; (v,e) \stackrel{\text{def}}{=} \\ vs \subseteq \; \{ \; \mathsf{0,1,...,}v - \mathsf{1} \; \} \; \land \\ \forall \, x \; y. \; x \in vs \land y \in vs \land x \neq y \Rightarrow \mathsf{is\_edge} \; e \; x \; y \\ \mathsf{max\_clique\_size} \; g \; \stackrel{\mathsf{def}}{=} \; \mathsf{max_{set}} \; \{ \; \mathsf{card} \; vs \; | \; \mathsf{is\_clique} \; vs \; g \; \} \end{array}$$

#### What are we trusting now?

- e.g. HOL model of verified checkers and correspondence to real system
- HOL4 theorem prover, including logic, implementation, and execution environment [Slind and Norrish, 2008]

# Proof logging in the Constraint Reasoning and Optimization Group

#### Earlier

• Fundamentals of redundancy notions in boolean decision problems (SAT) [Järvisalo, Heule, and Biere, 2012]

#### Currently

- Fundamentals of redundancy notions in boolean optimization (MaxSAT) [Berg and Järvisalo, 2019; Ihalainen, Berg, and Järvisalo, 2022]
- Certifying solvers and preprocessors [Ihalainen, Oertel, Tan, Berg, Järvisalo, Myreen, and Nordström, 2024; Berg, Bogaerts, Nordström, Oertel, and Vandesande, 2023]
- Multiobjective optimization [Jabs, Berg, Ihalainen, and Järvisalo, 2023]

#### Conclusion

Proof logging in automated reasoning:

- Guarantees correctness of results
- Supports development of increasingly complex reasoning into solvers.
- Provides audibility to third parties without access to the solver.

### Open Challenges

- Practical scaling.
- Proof logging e.g. PSPACE-complete problems.
- Proving bounds on the proof systems used.

I am hiring someone to work on these kinds of topics!

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