

Big O notation

Sutkamp book - chap. 14.2  
Script chap 2.

$$\mathbb{N} = \{0, 1, 2, \dots\}$$

$$\mathbb{Z} = \{\dots, -2, -1, 0, 1, 2, \dots\}$$

$$\mathbb{R} = \text{reals}$$

## Big O notation

Let's take  $f, g: \mathbb{N} \rightarrow \mathbb{N}$

We say  $f$  is of order  $g$  if

there is a constant  $c > 0 \in \mathbb{N}$  and

$n_0 \in \mathbb{N}$  s.t.  $f(n) \leq c \cdot g(n)$

$\forall n \geq n_0$ .

$\uparrow$  threshold.

$O(g) := \{ f \mid f \text{ is of order } g \}$

"Big O of g"

if  $f \in O(g)$  we say


g provides an asymptotic  
upper bound  
for f.

if  $f \in O(g)$   $g \in O(f)$

then we say they are of the same  
order.

$f = O(g)$  instead  $f \in O(g)$

↑  
abuse of  
institution.

 correct

$$p(n) = n^2 + o(n)$$

$$f(n) = n^2 + g(n), \quad \begin{matrix} g(n) \in O(n) \\ \quad \quad \quad \in o(n) \end{matrix}$$

Example

$$f(n) \sim n^2$$

$$g(n) \sim n^3$$

show that  $f \in O(g)$

$$c = 1$$

$$n > 1 \text{ i.e. } n_0 \geq 2$$

$$f(n) \leq g(n)$$

$$c, n_0 \text{ s.t.}$$

$$f(n) \leq c \cdot g(n)$$

$$\forall n \geq n_0$$

$n$	$f(n)$	$g(n)$
1	1	1
2	4	8
3	9	27

## Example

$$f(n) = n^2$$

$$g(n) = n^3$$

show that

$$g \notin O(f)$$

Assume that  
 $g \in O(f)$

$$\exists c, n_0 \text{ s.t.}$$

$$g(n) \leq c \cdot f(n) \quad \forall n \geq n_0.$$

$$\text{choose } n_1 = 1 + \max\{c, n_0\}$$

$$n_1^3 = n_1 \cdot n_1^2 > c \cdot n_1^2 \quad \text{where } n_1 > n_0.$$

3.

$$n_1 = 1 + \max\{c, n_0\}$$

$$n_1 = 1 + c \quad \text{or} \quad 1 + n_0$$

$$c > n_0$$

$$n_0 > c$$

$$1 + c > c$$

$$1 + n_0 > n_0$$

$$n_1 > c$$

$$n_1^3 = \underline{n_1 \cdot n_1^2} > \underline{c \cdot n_1^2}$$

we have a contradiction

$$\exists n \geq n_0 \text{ s.t.} \\ g(n) \geq c f(n)$$

this is the opposite of  
our definition,

Thus we may conclude  
that  $g \notin O(f)$ .



## Proof by Contradiction.

Show  $P$ .

$\therefore$  Assume  $\neg P$ .

if there arises a contradiction in  $\neg P$ ,

we conclude  $P$ .

Exhaustion.

Induction.

Construction.

Contraposition.

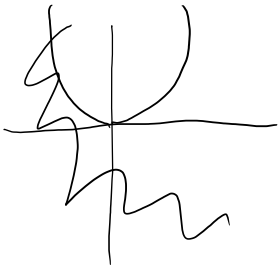


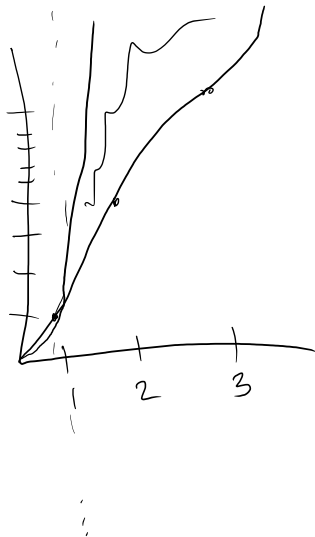
Be familiar

$$f = n^2$$

$$g = n^3$$

$$\mathbb{N} \rightarrow \mathbb{N}$$





$$f(n) = n^2 + 2n + 5$$

$$g(n) = n^2$$

show  $g \in O(f)$

$$\text{if } c=1 \\ n > 0$$

$$n^2 \leq n^2 + 2n + 5$$

$$0 \leq 2n + 5$$

show  $f \in O(g)$

$$f(n) = n^2 + 2n + 5$$

$$g(n) = n^2$$

$$\begin{aligned} n^2 + 2n + 5 &\leq n^2 + 2n^2 + 5n^2 \\ &\leq 8n^2 \end{aligned}$$

$$C = 8$$

$$\underline{\forall n \geq 1}$$