

## Time Complexity (TM)

For any  $M$ , its time complexity

$$tc_M: \mathbb{N} \rightarrow \mathbb{N} \text{ s.t. } tc_M(n)$$

is the maximum number of transitions taken by a computation ("run") of  $M$  on an input of length  $n$ .

$$t_{cm} : \mathbb{N} \rightarrow \mathbb{N}$$

input  
length

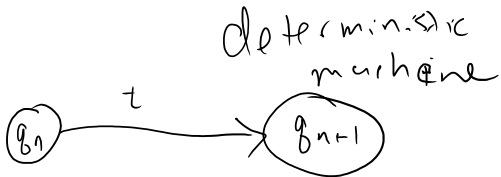
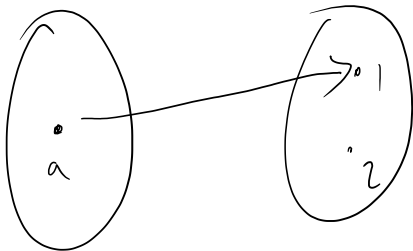
max  
transitions  
taken

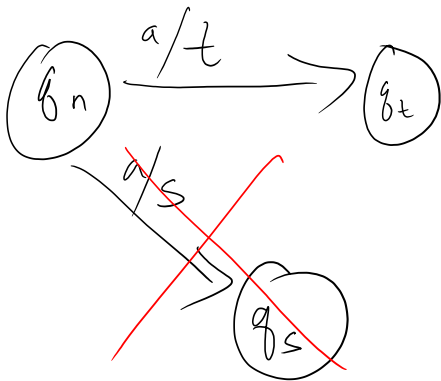
$$\underline{t_{cm}(n)}$$

$$\left( f(n) \in O(n^2) \right)$$

TM is deterministic if  
 $\delta$  is a (partial) function.

Intuitively, this means for any  
state and symbol read from  
the tape, there is exactly one  
transition.





A language  $L$  is accepted in deterministic time, if there is a single-tape deterministic Turing Machine for the language.

$$t_{cm} \in O(f(n))$$

↑  
complexity  
class.

Example: Palindrome Example (script)

palindrome: forward - reading =  
backward - reading

racecar

a's and b's

aa

ab a

abba

ababa

" - trivial pluridromae

$aa$  or  $bb$

$\uparrow \uparrow$

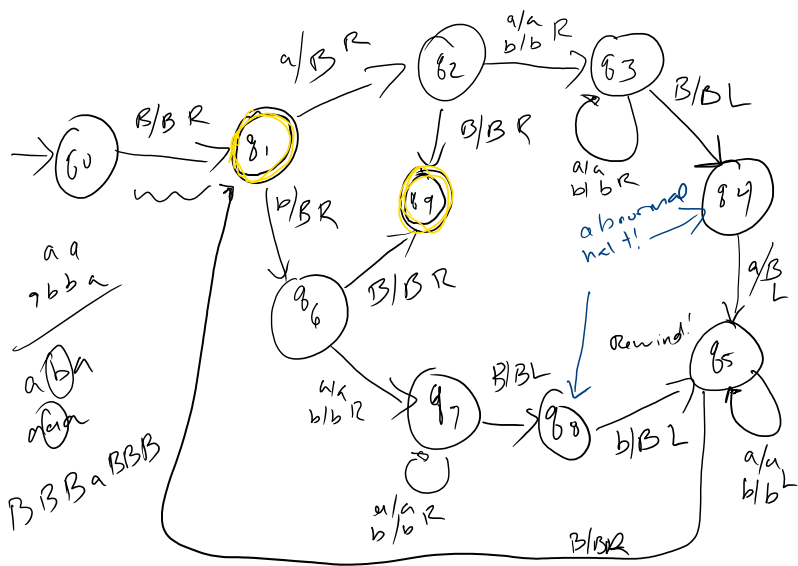
$a \underline{u} a$

$\uparrow \quad \uparrow$

$b \underline{u} a$

where  $u$  is  
a palindrome.







$$4 + 3 + 2 + 1$$

$$\sum_{i=0}^n i = \frac{n^2 + n}{2}$$

$$n^2 + n + 1$$

exact  
complexity

$$T_{cm} = \underline{n^2 + n + 1} \in \underline{O(n^2)} \quad \square$$