

Assignment 9

Problem 1) For $(\log n + 1)^2 + (\log n + 1)(n^2 + 1) = O(n^t)$ to hold true, n^t must be an upper bound of $(\log n + 1)^2 + (\log n + 1)(n^2 + 1)$. For this to be true, $\lim_{n \rightarrow \infty} \frac{(\log n + 1)^2 + (\log n + 1)(n^2 + 1)}{n^t} \leq d \mid d \in \mathbb{R} > 0$. This statement is true unless d is ∞ , so t must be an integer such that n^t grows at a faster rate than $(\log n + 1)^2 + (\log n + 1)(n^2 + 1)$. Since $(\log n + 1)^2 + (\log n + 1)(n^2 + 1)$ grows at a rate of $n^2(\log n)^2$, t must be ≥ 3 for the limit to approach 0, which would make the original statement hold true. For example, if $t = 2$, then the limit would approach ∞ and the original statement would not hold true. Thus, 3 is the least integer $t \mid (\log n + 1)^2 + (\log n + 1)(n^2 + 1) = O(n^t)$.

Problem 2) Counter-example: Let $f(n) = n+1$; $g(n) = n-1$; $h(n) = n$

First, prove $f(n)$ and $g(n)$ are both $\Theta(h(n))$.

$$c \leq \lim_{n \rightarrow \infty} \frac{n+1}{n} \leq d$$

$$c \leq \lim_{n \rightarrow \infty} 1 + \frac{1}{n} \leq d$$

$$c \leq \lim_{n \rightarrow \infty} 1 + 0 \leq d$$

$$c \leq 1 \leq d; c=1, d=1$$

$f(n) = \Theta(h(n))$ holds true.

$$c \leq \lim_{n \rightarrow \infty} \frac{n-1}{n} \leq d$$

$$c \leq \lim_{n \rightarrow \infty} 1 - \frac{1}{n} \leq d$$

$$c \leq \lim_{n \rightarrow \infty} 1 - 0 \leq d$$

$$c \leq 1 \leq d; c=1, d=1$$

$g(n) = \Theta(h(n))$ holds true.

Consider $(f-g)(n) = \Theta(h(n))$

$$c \leq \lim_{n \rightarrow \infty} \frac{n+1 - (n-1)}{n} \leq d$$

$$c \leq \lim_{n \rightarrow \infty} \frac{n+1 - n+1}{n} \leq d$$

$$c \leq \lim_{n \rightarrow \infty} \frac{2}{n} \leq d$$

$$c \leq 0 \leq d, \nexists c \text{ such that } c \in \mathbb{R} > 0.$$

As depicted in this counter-example, even if $f(n)$ and $g(n)$ are both $\Theta(h(n))$, $(f-g)(n) \neq \Theta(h(n))$.

Problem 3) Consider $f(n) = O(g(n))$ and $g(n) = O(f(n))$. This means that:

$$(i) \frac{f(n)}{g(n)} \leq a \mid a \in \mathbb{R} > 0 \text{ AND } (ii) \frac{g(n)}{f(n)} \leq b \mid b \in \mathbb{R} > 0 \text{ which implies } \frac{f(n)}{g(n)} \geq \frac{1}{b}$$

$$\frac{1}{3} \leq \frac{1}{2} \Rightarrow \frac{2}{3} \geq \frac{2}{2}; \quad 5, \frac{1}{5} \in \mathbb{R} > 0$$

Taking (i) and (ii) into account, we get $\frac{1}{b} \leq \frac{f(n)}{g(n)} \leq a$. If $b \in \mathbb{R} > 0, \frac{1}{b} \in \mathbb{R} > 0$

Since $a, \frac{1}{b} \in \mathbb{R} > 0$, $\frac{f(n)}{g(n)}$ is upper and lower bounded by constants. Therefore, if $f(n) = O(g(n))$ and $g(n) = O(f(n))$, then $f(n) = \Theta(g(n))$.