

Lattice and Heyting algebra objects in a topos

§ IV.8

- Recall:
- A lattice is a poset that has all binary products and coproducts (when viewed as a category)
 - ⑩ a set with 2 distinguished elements \perp, \top and two associative and commutative binary operations \wedge and \vee st.:

$$\begin{array}{llll} \text{meet} & \text{join} & x \wedge x = x & x \vee x = x \\ & & T \wedge x = x & \perp \vee x = x \\ & & x \wedge (y \vee x) = x = (x \wedge y) \vee x & \text{idempotent} \\ & & & \text{absorption} \end{array}$$

- A Heyting algebra is a Cartesian closed poset with all finite products and coproducts.
 - ⑩ a lattice with \perp and \top such that for any pair of elements x, y , there is an exponential $y^x := (x \Rightarrow y)$ characterized by (the adjunction)
- $$z \leq (x \Rightarrow y) \text{ iff } z \wedge x \leq y$$

Fact: The implication \Rightarrow in a Heyting algebra satisfies:

- (1) $(x \Rightarrow x) = \top$
- (2) $x \wedge (x \Rightarrow y) = x \wedge y$
- (3) $y \wedge (x \Rightarrow y) = y$
- (4) $x \Rightarrow (y \wedge z) = (x \Rightarrow y) \wedge (x \Rightarrow z)$

The "equational identities of a Heyting alg."

Now we will head in a different direction and think about lattices and Heyting algs that exist within a category other than Set.

- C: category with finite limits

def. A LATTICE OBJECT / INTERNAL LATTICE in C is $L \in \text{Ob}(C)$ together with 2 arrows:

$$\begin{array}{ll} \text{"meet"} & \wedge : L \times L \rightarrow L \\ \text{"join"} & \vee : L \times L \rightarrow L \end{array}$$

st. the identities in the equational def" of a lattice create a commutative diagram.

⑩ idempotent: $x \wedge x = x$

(e) The absorption law $x \wedge (y \vee x) = x = (x \wedge y) \vee x$ as a commutative diagram:

$$\begin{array}{c}
 \text{LHS} \\
 \begin{array}{ccccc}
 & \xleftarrow{\wedge} & & \xrightarrow{\wedge} & \\
 x & \downarrow P & & \uparrow P & x \\
 (x,y) & \xrightarrow{\delta \times 1} & (x,x,y) & \xrightarrow{1 \times \tau} & (x,y,x) \\
 & \downarrow & & \uparrow & \\
 & & (x,y) \vee x & &
 \end{array}
 \end{array}$$

$P := \pi_1$ (proj onto first factor)

$\tau: L \times L \rightarrow L \times L$

(twist map that interchanges the factors of the product)

$\delta: L \rightarrow L \times L$: diagonal map

- such a lattice object has T, \perp when there are arrows (where $1 \in C$ is terminal)

"top"

$$T: 1 \rightarrow L$$

"bottom"

$$\perp: 1 \rightarrow L$$

such that:

$$\begin{array}{l}
 x \vee \perp = x \quad \text{Really: } \perp(1) \text{ (the bottom arrow applied to 1)} \\
 \text{and} \\
 x \wedge T = x
 \end{array}$$

- (f) They make the following composites the identity:

$$\begin{array}{c}
 L \cong L \times 1 \xrightarrow[1 \times \perp]{(x,1)} L \times L \xrightarrow[V]{(x,\perp(1))} x \vee \perp(1) \\
 \curvearrowright 1
 \end{array}$$

$$\begin{array}{c}
 L \cong L \times 1 \xrightarrow[1 \times T]{(x,1)} L \times L \xrightarrow[\wedge]{(x,1)} L \\
 \curvearrowright 1
 \end{array}$$

def. L a lattice object.

Is an INTERNAL HETTING ALGEBRA IN C (or just "Hetting algebra") if \exists binary operation $\Rightarrow: L \times L \rightarrow L$ satisfying the diagrammatic version of the identities for \Rightarrow .

$$\Rightarrow: (x,x) \mapsto (x \Rightarrow x)$$

$T(1)$ = "one" in the lattice

1 element = "TOP"

$$T: 1 \rightarrow L$$

$$\begin{array}{c}
 (1) \quad 1 \xrightarrow{T} L(x \Rightarrow x) \\
 \uparrow \quad \nearrow \\
 L \times L \xrightarrow{(x,x)}
 \end{array}$$

$$x \Rightarrow x = T(1)$$

$$\begin{array}{c}
 (2) \quad x \wedge (x \Rightarrow y) = x \wedge y \quad (x,y) \xrightarrow[\wedge]{L \times L \rightarrow L} x \wedge y \\
 \delta \times 1 \downarrow \quad \uparrow \wedge \quad \uparrow \wedge \\
 L \times L \times L \xrightarrow[1 \times \Rightarrow]{(x,x,y)} L \times L \xrightarrow[\wedge]{(x,y)} x \wedge (x \Rightarrow y) \\
 (x,x,y) \mapsto (x, x \Rightarrow y)
 \end{array}$$

• For L a lattice object, can define corresponding partial order on L by:

$$x \leq y \text{ iff } x \wedge y = x$$

def. Subobject \leq_L of $L \times L$ is the equalizer:

$$\begin{array}{ccccc} \leq_L & \xrightarrow{e} & L \times L & \xrightarrow{\wedge} & L \\ \uparrow f! & \nearrow A & \downarrow P & & \\ D & & xxy \mapsto x & & \\ & & xxy \mapsto x \wedge y & & \end{array}$$

FACT (L, \leq_L) is an INTERNAL POSET • a poset that lives in the ambient cat/topos C

\leq_L reflexive:

$$\begin{array}{ccc} L & \xrightarrow{f} & L \times L \\ & \searrow & \uparrow \leq_L \\ & & \end{array}$$

• Transitivity of \leq_L :

• The subobject $\langle \pi_1 \circ v, \pi_2 \circ u \rangle : C \rightarrow L \times L$ factors through $e : \leq_L \rightarrow L \times L$ where C is the following pullback, with projections v and u :

$$\begin{array}{ccccc} C & \xrightarrow{u} & \leq_L & & ?? \\ \downarrow v & & \downarrow e & & \\ \leq_L & \xrightarrow{e} & L \times L & \xrightarrow{\pi_1} & L \\ & \downarrow & \downarrow \pi_2 & & \\ & & L \times L & \xrightarrow{\pi_2} & L \end{array}$$

(second defⁿ)

def. An INTERNAL HETTING ALGEBRA of a topos \mathcal{E} is an internal lattice of \mathcal{E} with an additional binary operation $\Rightarrow : L \times L \rightarrow L$ such that the 2 subobjects P and Q of $L \times L \times L$ defined by the 2 pullback squares below are equivalent subobjects:

$$\begin{array}{ccccc} P & \dashrightarrow & \leq_L & \dashleftarrow & Q \\ | & & \downarrow e & & | \\ \downarrow & & \downarrow & & \downarrow \\ L \times L \times L & \xrightarrow{\wedge \times 1} & L \times L & \xleftarrow{1 \times \Rightarrow} & L \times L \times L \end{array}$$

(ie) $\stackrel{\text{RHS}}{z \leq (x \Rightarrow y)}$ iff $\stackrel{\text{LHS}}{z \wedge x \leq y}$ (*)
 (equational defⁿ of Heyting alg.)

- This defⁿ is equivalent to the other one (equational identities for \Rightarrow)
 - equivalence proof: Yoneda applied to hom-sets $\text{Hom}(X, L)$

def. A HOMOMORPHISM OF LATTICES (or of HEYTING ALGEBRA OBJECTS) $L \rightarrow L'$

is $f: L \rightarrow L'$ that commutes with all operations.

$$\begin{array}{ccccc} 1 & \xrightarrow{T} & L & \xleftarrow{\wedge} & L \times L \\ \parallel & & f \downarrow & & \downarrow f \times f \\ 1 & \xrightarrow{T} & L' & \xleftarrow{\wedge} & L' \times L' \\ & \perp & & \vee, \Rightarrow & \end{array}$$

operation replacements

Thm 1
(EXTERNAL)

Let $A \in \text{Ob}(\mathcal{E})$.

The poset $\text{Sub}_{\mathcal{E}} A$ of subobjects of A has the structure of a Heyting algebra. This structure is natural in A .

↪ (e) The pullback along any morphism $k: A \rightarrow B$ induces a map k^{-1} of Heyting algebras as in:

$$\begin{array}{ccc} \text{Sub}_{\mathcal{E}} B & \xrightarrow{k^{-1}} & \text{Sub}_{\mathcal{E}} A \\ i_B \downarrow & & \downarrow i_A \\ \text{slices } I \rightsquigarrow \mathcal{E}/B & \xrightarrow{k^*} & \mathcal{E}/A \\ & \nearrow o & \end{array}$$

The "change-of-base" functor

$$\begin{array}{c} x' \\ f' \searrow \\ A \xrightarrow{k} B \\ \curvearrowright x \end{array}$$

(p.192) Consider any $k: A \rightarrow B$ in a topos \mathcal{E} .

Pullback along k then turns each object

$f: X \rightarrow B$ of \mathcal{E}/B into an object f' of \mathcal{E}/A :

$$\begin{array}{ccc} A \times_B B & =: X' & \longrightarrow X \\ (\text{the pullback}) & f' \downarrow & \downarrow f \\ & A & \xrightarrow{k} B \\ & \text{Use map } k: A \rightarrow B & \end{array}$$

defines a "change-of-base" functor / "PULLBACK FUNCTOR"

$$k^*: \mathcal{E}/B \rightarrow \mathcal{E}/A$$

take an element of \mathcal{E}/B

Proof of Thm 1

- $A \in \text{Ob}(\mathcal{E}) \rightsquigarrow \text{Sub}_\mathcal{E}(A)$ has lattice structure.
("factorization + images" section)

PROP 3 in § 6 Sub A is a lattice.

For each $k: A \rightarrow B$, the pullback along k is a morphism

$K^{-1}: \text{Sub}B \rightarrow \text{Sub}A$ of posets

$(\exists_k \dashv K^{-1})$ (ie, a functor). This functor has a left adjoint $\exists_K: \text{Sub}A \rightarrow \text{Sub}B$
 $S \mapsto (\text{image in } B \text{ under } k)$

- $S \rightarrow A, T \rightarrow A$: two subobjects in Sub A.

- Form their "intersection" (greatest lower bound in Sub A) by taking the pullback:

$$\begin{array}{ccc} S \cap T & \longrightarrow & T \\ \downarrow & \lrcorner & \downarrow \\ S & \longrightarrow & A \end{array}$$

toposes have
pullbacks.

- For their "union" (least upper bound), form the coproduct $(S+T)$ in the topos.
Then, $\exists! S+T \rightarrow A$ (by defⁿ of coproduct).
This has an image M (by result on factorizations + images)
 $\rightsquigarrow M \rightarrow A$ is a subobject that contains both S and T

$$\begin{array}{ccc} S+T & \xleftarrow{\quad} & T \\ \uparrow \text{epi} & \searrow \text{mono} & \downarrow \\ S & \xrightarrow{\exists! \text{ coproduct}} & A \end{array} \qquad M \rightarrow A \in \text{Sub } A$$

So we have $\wedge: \text{Sub } A \times \text{Sub } A \rightarrow \text{Sub } A$ and it is natural in A

(The pullback functor $K^{-1}: \text{Sub } B \rightarrow \text{Sub } A$ is sl. $K^{-1}(S \cap T) = K^{-1}(S) \cap K^{-1}(T)$)

Using the iso:

$$\text{Hom}(A, \wedge) \cong \text{Sub}(A)$$

(natural in A), we obtain op \wedge_A that makes the following commute:

$$\begin{array}{ccc}
 \text{Sub}(A) \times \text{Sub}(A) & \xrightarrow{\cap} & \text{Sub}(A) \\
 \parallel ? & & \downarrow ? \\
 \text{Hom}(A, \perp) \times \text{Hom}(A, \perp) & & \\
 \parallel ? & & \\
 \text{Hom}(A, \perp \times \perp) & \xrightarrow{\quad \wedge_A \quad} & \text{Hom}(A, \perp)
 \end{array}$$

Then \wedge_A is also natural in A so by Yoneda ($A := \perp \times \perp$, apply \wedge_A to id), then \wedge_A comes from a *uniquely determined map*

$$\wedge : \perp \times \perp \rightarrow \perp \quad (\text{via composition})$$

⑩ If subobjects S, T of A have char maps $s, t : A \rightarrow \perp$, then $S \cap T$ has char map

$$A \xrightarrow{s,t} \perp \times \perp \xrightarrow{\wedge} \perp$$

And we will use the following for the next thm:

For a fixed obj. A , the op.

$$\cap : \text{Sub}(A \times X) \times \text{Sub}(A \times X) \rightarrow \text{Sub}(A \times X) \quad \text{is } \underline{\text{natural}} \text{ in } X$$

By

$\text{Sub}(A \times X) \cong \text{Hom}(X, PA)$, we get q:

$$\text{Hom}(X, PA \times PA) \cong \text{Hom}(X, PA) \times \text{Hom}(X, PA) \xrightarrow{\wedge_X} \text{Hom}(X, PA)$$

(natural in X)

By Yoneda, we get a cov. map

$$\wedge : PA \times PA \rightarrow PA$$

This is *natural* in A : $\forall k : B \rightarrow A$, we have

$$\begin{array}{ccc}
 PA \times PA & \xrightarrow{\wedge} & PA \\
 \downarrow PK \times PK & & \downarrow PK \\
 PB \times PB & \xrightarrow{\wedge} & PB
 \end{array}
 \quad \text{commutes}$$

CLAIM: $\text{Sub}_{\mathcal{E}}(1)$ has exponentials

- The exponential U^V of two open objects $U \rightarrow 1$ and $V \rightarrow 1$ is again open.

$\rightsquigarrow U^V \rightarrow 1$ is monic by universal property of the exponential.

$\rightsquigarrow \text{Sub}_{\mathcal{E}}(1)$ has exponentials

Slice:

$$\text{Sub}_{\mathcal{E}}(A) \cong \text{Sub}_{\mathcal{E}/A}(1)$$

$\rightsquigarrow \text{Sub}_{\mathcal{E}}(A)$ is a Heyting algebra with the exponential in \mathcal{E}/A as its implication operator.

The structure is natural in A :

Let $k: A \rightarrow B$ be a morphism in \mathcal{E} . And consider the comm. square:

$$\begin{array}{ccc} \text{Sub}_{\mathcal{E}}(B) & \xrightarrow{k^*} & \text{Sub}_{\mathcal{E}}(A) \\ i_B \downarrow & & \downarrow i_A \\ \mathcal{E}/B & \xrightarrow{k^*} & \mathcal{E}/A \end{array}$$

i_A :
inclusion that identifies
subobjects of A in \mathcal{E}
with subobjects of 1 in
 \mathcal{E}/A

Then IV.7.2 (Thm 2 of previous section)

By thm in "slice of topos is a topos", then $\boxed{k^*}$ preserves exponentials, sums, and epimorphisms.

By lattice structure of $\text{Sub}_{\mathcal{E}}(B)$ and $\text{Sub}_{\mathcal{E}}(A)$, then K^{-1} is a Heyting alg. homomorphism. \square

Thm 2
(INTERNAL)

Let $A \in \text{Ob}(\mathcal{E})$.

The power object PA is an internal Heyting algebra.

(ie, the subobject classifier $\Omega := P1$ is an internal Heyting algebra).

The structure is natural in A :

⑩ Let $k: A \rightarrow B$ be a morphism in \mathcal{E} .

Then, the induced map $Pk: PB \rightarrow PA$ is a Heyting algebra homomorphism.

For each $X \in \text{Ob}(\mathcal{E})$, the internal structure on PA makes $\text{Hom}(X, PA)$ an external Heyting algebra so that the canonical isomorphism equational identities hold.

What do we NTS?

- PA is a lattice object in \mathcal{E} :

There are meet +
join arrows s.t.

equational identities hold.

- There is a binary op. $\dashv: PA \times PA \rightarrow PA$

$$\text{Sub}_{\mathcal{E}}(A \times X) \cong \text{Hom}_{\mathcal{E}}(X, PA)$$

is an isomorphism of external Heyting algebras.

[proof]

The proof starts with some ideas stated in Section 6: "Factorizations and images":

And we will use the following for the next thm:

For a fixed ob. A , the op.

$$\wedge: \text{Sub}(A \times X) \times \text{Sub}(A \times X) \rightarrow \text{Sub}(A \times X) \quad \text{is } \underline{\text{natural}} \text{ in } X$$

By

$\text{Sub}(A \times X) \cong \text{Hom}(X, PA)$, we get \wedge :

$$\text{Hom}(X, PA \times PA) \cong \text{Hom}(X, PA) \times \text{Hom}(X, PA) \xrightarrow{\wedge_X} \text{Hom}(X, PA)$$

(natural in X)

By Yoneda, we get a cov. map

$$\wedge: PA \times PA \rightarrow PA$$

This is natural in A : $\forall k: B \rightarrow A$, we have

$$\begin{array}{ccc} PA \times PA & \xrightarrow{\wedge} & PA \\ \downarrow Pk \times Pk & & \downarrow Pk \\ PB \times PB & \xrightarrow{\wedge} & PB \end{array}$$

commutes

So we have

$$\wedge : PA \times PA \longrightarrow PA$$

where for any $X \in \mathcal{E}$, the meet operation on $\text{Hom}_{\mathcal{E}}(X, PA)$ induced by composition corresponds to meet in the lattice

$$\text{Sub}_{\mathcal{E}}(A \times X) \quad \text{via the iso-}$$

$$\text{Hom}_{\mathcal{E}}(X, PA) \cong \text{Sub}_{\mathcal{E}}(A \times X)$$

- Other ops are defined similarly.

(eg) The (\Rightarrow) operation

For each $X \in \text{Ob}(\mathcal{E})$, $\text{Sub}_{\mathcal{E}}(A \times X)$ has an implication operator

$$\Rightarrow : \text{Sub}_{\mathcal{E}}(A \times X) \times \text{Sub}_{\mathcal{E}}(A \times X) \longrightarrow \text{Sub}_{\mathcal{E}}(A \times X)$$

that is natural in X .

$\rightsquigarrow \exists! \Rightarrow_X$ (natural in X) such that the following commutes:

$$\begin{array}{ccc}
 \text{Sub}_{\mathcal{E}}(A \times X) \times \text{Sub}_{\mathcal{E}}(A \times X) & \xrightarrow{\Rightarrow} & \text{Sub}_{\mathcal{E}}(A \times X) \\
 \parallel ? & & \parallel ? \\
 \text{Hom}_{\mathcal{E}}(X, PA) \times \text{Hom}_{\mathcal{E}}(X, PA) & & \\
 \parallel ? & & \parallel ? \\
 \text{Hom}_{\mathcal{E}}(X, PA \times PA) & \xrightarrow{\Rightarrow_X} & \text{Hom}_{\mathcal{E}}(X, PA)
 \end{array}$$

By naturality of \Rightarrow_X in X , then it must be induced (via composition) by a uniquely determined map

$$\Rightarrow : PA \times PA \longrightarrow PA$$

- Top and bottom elements: similar

• $Pk: PB \rightarrow PA$ is a homomorphism of Heyting algebras:

By the naturality of the Heyting alg. structure on $\text{Sub}_\mathcal{E}(A)$ and the comm. diagram:

$$\begin{array}{ccc} \text{Sub}_\mathcal{E}(B \times X) & \xrightarrow{\sim} & \text{Hom}_\mathcal{E}(X, PB) \\ \downarrow (k \times 1)^{-1} & & \downarrow \text{Hom}_\mathcal{E}(X, Pk) \\ \text{Sub}_\mathcal{E}(A \times X) & \xrightarrow{\sim} & \text{Hom}_\mathcal{E}(X, PA) \end{array}$$

Since $(k \times 1)^{-1}$ is a Heyting alg. hom, so is $\text{Hom}_\mathcal{E}(X, Pk)$ for each X in \mathcal{E} .

$\Rightarrow Pk$ is a hom of internal Heyting algebras
 (by definitions)

□

Rem: Internal Heyting alg. structure on Ω is the unique one such that

$\text{Sub}_\mathcal{E}(X) \xrightarrow{\sim} \text{Hom}_\mathcal{E}(X, \Omega)$ is a Heyting alg. isomorphism.