

Reedy categories and diagrams

Motivation:

Goal: To introduce the relevant terminology and ideas to understand the proof of the model category structure on Reedy diagrams.

Sources:

- "Model categories and their localizations" [Hirschhorn]
- "A primer on homotopy colimits" [Dugger]

§ 1. Preliminaries

Def C a small category.

C is a **REEDY CATEGORY** if it has two wide subcategories \vec{C} and $\overset{\leftarrow}{C}$ where each object is assigned a "**DEGREE**" $n \geq 0$ such that the following is satisfied:

- { (1) Every non-identity morphism of \vec{C} raises degree.
- (2) Every non-identity morphism of $\overset{\leftarrow}{C}$ lowers degree.

(3) Every morphism g of C has a unique factorization

$$g = \vec{g} \circ \overset{\leftarrow}{g} \quad \vec{g} \in \vec{C}, \quad \overset{\leftarrow}{g} \in \overset{\leftarrow}{C}$$

Remark :

- More general definition: degree function takes ORDINAL values

Some basic properties :

• C Reedy $\Rightarrow C^{\text{op}}$ is Reedy with $\overset{\rightarrow}{C^{\text{op}}} := (\overset{\leftarrow}{C})^{\text{op}}$ and $\overset{\leftarrow}{C^{\text{op}}} := (\vec{C})^{\text{op}}$

• C, D Reedy $\Rightarrow C \times D$ is Reedy with $\overset{\rightarrow}{C \times D} := \vec{C} \times \vec{D}$ and $\overset{\leftarrow}{C \times D} := \overset{\leftarrow}{C} \times \overset{\leftarrow}{D}$

Example: The cosimplicial and simplicial indexing categories.

Def

$$n \in \mathbb{Z}_{\geq 0}$$

$$[n] := \{0, 1, \dots, n\}$$

$$\text{cat } \Delta := \left\{ \begin{array}{l} \text{objects: } \{[n] \mid n \geq 0\} \\ \text{morphisms: } \Delta([n], [k]) := \{ \delta: [n] \rightarrow [k] \mid \delta(i) \leq \delta(j) \quad \forall 0 \leq i \leq j \leq n \} \end{array} \right.$$

• Δ is called the COSIMPPLICIAL INDEXING CATEGORY

• Δ^{op} is called the SIMPPLICIAL INDEXING CATEGORY

• M a category

\rightsquigarrow Functor $\Delta^{\text{op}} \rightarrow M$: A SIMPPLICIAL OBJECT IN M

\rightsquigarrow Functor $\Delta \rightarrow M$: A COSIMPPLICIAL OBJECT in M

Notation:

• X a simplicial object in M \rightsquigarrow denote $X_{[n]} =: X_n$

• X a cosimplicial object in M \rightsquigarrow denote $X_{[n]} =: X^n$

These are Reedy categories!

• Cosimplicial indexing category Δ

with $\deg([n]) := n$

$$\overset{\rightarrow}{\Delta} := \{ \text{injective maps} \}$$

$$\overset{\leftarrow}{\Delta} := \{ \text{surjective maps} \}$$

• Simplicial indexing category Δ^{op}

with $\deg([n]) := n$

$$\overset{\rightarrow}{\Delta^{\text{op}}} := \{ \text{op}^\circ \text{ of surjective maps} \}$$

$$\overset{\leftarrow}{\Delta^{\text{op}}} := \{ \text{op}^\circ \text{ of injective maps} \}$$

§ 2. Filtrations

Def

- \mathcal{C} is Reedy

- $n \in \mathbb{Z}_{\geq 0}$

- The n -FILTRATION $(F^n\mathcal{C})$ is the full subcategory of \mathcal{C} with all objects $\alpha \in \mathcal{C}$ s.t. $\deg(\alpha) \leq n$

④ 0-filtration of $\mathcal{C} = \left\{ \begin{array}{l} \text{obj : } \alpha \text{ s.t. } \deg(\alpha) = 0 \\ \text{mor : } \{1_\alpha \mid \deg(\alpha) = 0\} \end{array} \right\} \hookrightarrow \text{Reedy factorization}$

Proposition

\mathcal{C} Reedy. $n \in \mathbb{Z}_{\geq 0}$.

$\Rightarrow F^n\mathcal{C}$ is Reedy with:

$$\cdot \vec{F^n\mathcal{C}} := \vec{\mathcal{C}} \cap (F^n\mathcal{C})$$

$$\cdot \overleftarrow{F^n\mathcal{C}} := \overleftarrow{\mathcal{C}} \cap (F^n\mathcal{C})$$

and

$$\mathcal{C} = \bigcup_{n \geq 0} F^n\mathcal{C} \quad \text{where} \quad F^0\mathcal{C} \subset F^1\mathcal{C} \subset F^2\mathcal{C} \subset \dots$$

§ 3. Diagrams

In this section: \mathcal{C} is a Reedy category, M is a model category

Recall: The 0-filtration of a Reedy category has no non-identity maps.

\Rightarrow define a diagram $X: F^0\mathcal{C} \rightarrow M$ by choosing an object $X_\alpha \in M$ for each object $\alpha \in \mathcal{C}$ with $\deg \alpha = 0$.

Suppose we have $X: (F^{n-1}\mathcal{C}) \rightarrow M$

want to extend it to a diagram $X: F^n\mathcal{C} \rightarrow M$

Define $X: F^n \mathcal{C} \rightarrow M$ by:

(I) Objects

For each $\alpha \in \mathcal{C}$ with $\deg \alpha = n$, choose $X_\alpha \in M$.

(II) Morphisms

For each $\beta \in F^{n-1} \mathcal{C}$, and $\boxed{\beta \rightarrow \alpha}$ in $F^n \mathcal{C}$,

(*) Want a map $X_\beta \rightarrow X_\alpha$ s.t. if $\beta \rightarrow \beta' \in F^{n-1} \mathcal{C}$ and

$$\begin{array}{ccc} \beta & \xrightarrow{\quad} & \beta' \\ \downarrow & \searrow \alpha & \swarrow \\ & \alpha & \end{array} \text{commutes in } F^n \mathcal{C}$$

$$\Rightarrow \begin{array}{ccc} X_\beta & \xrightarrow{\quad} & X_{\beta'} \\ & \searrow & \swarrow \\ & X_\alpha & \end{array} \text{commutes in } M.$$

Let $i^n: F^{n-1} \mathcal{C} \rightarrow F^n \mathcal{C}$ be the inclusion functor.

Then, (*) is equivalent to choosing a map

$$\text{Lan}_{i^n} X(\alpha) = \text{colim}_{(i^n \downarrow \alpha)} X \rightarrow X$$

Dually, for each $f \in F^{n-1} \mathcal{C}$ and $\boxed{\alpha \rightarrow f}$ in $F^n \mathcal{C}$,

We need a map $X_\alpha \rightarrow X_f$ such that

If $\boxed{\alpha}$ commutes in $F^n \mathcal{C}$, then

$$\begin{array}{ccc} \alpha & & \\ \swarrow & \searrow & \\ f & \rightarrow & f' \\ \downarrow & & \\ X_\alpha & & \end{array}$$

$$\begin{array}{ccc} & & \\ & \searrow & \swarrow \\ & X_f & \rightarrow X_{f'} \\ & & \end{array} \text{commutes in } M.$$

This is equivalent to choosing a map

$$X_\alpha \longrightarrow \lim_{(\alpha \downarrow i^n)} X = \text{Ran}_{i^n} X(\alpha)$$

Key: The colimit and limit chosen above are independent of our choice of degree function on the Reedy category \mathcal{C} .

But, we need some conditions on the choices for these maps.

If $(\beta \rightarrow \gamma) \in F^{n-1}\mathcal{C}$ and

$$\begin{array}{ccc} & \alpha & \\ & \nearrow & \searrow \\ \beta & \longrightarrow & \gamma \end{array} \quad \text{commutes in } F^n\mathcal{C}$$

\Rightarrow

$$\begin{array}{ccc} X_\alpha & & \\ \nearrow & \searrow & \\ X_\beta & \longrightarrow & X_\gamma \end{array} \quad \text{must commute in } M.$$

Equivalent to:

$$\text{colim}_{(i^n \downarrow \alpha)} X \xrightarrow{\quad} \lim_{(\alpha \downarrow i^n)} X \quad (*)$$

- Theorem: (*) is enough to construct an extension of X from $F^{n-1}\mathcal{C}$ to $F^n\mathcal{C}$.
(Exercise)

§ 4. Latching and matching objects

$\alpha \in \mathbb{C}$

\mathbb{C} a Reedy category, $\alpha \in \mathbb{C}$ an object.

M a model category, $X: \mathbb{C} \rightarrow M$

- Def** • LATCHING CATEGORY OF \mathbb{C} at α , denoted $\partial(\vec{\mathbb{C}} \downarrow \alpha)$, is the full subcategory of $(\vec{\mathbb{C}} \downarrow \alpha)$ containing all the objects except 1_α .
- MATCHING CATEGORY OF \mathbb{C} at α , denoted $\partial(\alpha \downarrow \vec{\mathbb{C}})$, is the full subcategory of $(\alpha \downarrow \vec{\mathbb{C}})$ containing all the objects except 1_α .

Facts: ① $(\partial(\vec{\mathbb{C}} \downarrow \alpha))^{\text{op}} \cong \partial(\alpha \downarrow \vec{\mathbb{C}}^{\text{op}})$

② $(\partial(\alpha \downarrow \vec{\mathbb{C}}))^{\text{op}} \cong \partial(\vec{\mathbb{C}}^{\text{op}} \downarrow \alpha)$

$$\partial(\vec{\mathbb{C}} \downarrow \alpha) \xhookrightarrow{i} (\vec{\mathbb{C}} \downarrow \alpha)$$

↑

- Df** • LATCHING OBJECT OF X at α : $L_\alpha X := \underset{\partial(\vec{\mathbb{C}} \downarrow \alpha)}{\text{colim}} X = \text{Lan}_i X(\alpha)$
- LATCHING MAP OF X at α : $L_\alpha X \rightarrow X$

- MATCHING OBJECT OF X at α : $M_\alpha X := \underset{\partial(\alpha \downarrow \vec{\mathbb{C}})}{\lim} X = \text{Ran}_i X(\alpha)$
- MATCHING MAP OF X at α : $X_\alpha \rightarrow M_\alpha X$

Df Suppose $\alpha \in \mathbb{C}$ has degree n .

- $\partial(\alpha \downarrow F^n \mathbb{C}) := \left\{ \begin{array}{l} \text{full subcategory of } (\alpha \downarrow F^n \mathbb{C}) \text{ with objects the maps} \\ \alpha \xrightarrow{g} \beta \text{ s.t. } \exists \text{ a factorization } \alpha \xrightarrow{\tilde{g}} r \xrightarrow{\tilde{q}} \beta \text{ with } \tilde{q} \in \vec{\mathbb{C}}, \tilde{g} \in \vec{\mathbb{C}} \\ \text{and } \tilde{q} \neq 1_\alpha \end{array} \right\}$
- $\partial(F^n \mathbb{C} \downarrow \alpha) := \left\{ \begin{array}{l} \text{full subcategory of } (F^n \mathbb{C} \downarrow \alpha) \text{ with objects the maps } \beta \xrightarrow{g} \alpha \\ \text{s.t. } \exists \text{ factorization } \beta \xrightarrow{\tilde{q}} r \xrightarrow{\tilde{g}} \alpha \text{ with } \tilde{q} \in \vec{\mathbb{C}}, \tilde{g} \in \vec{\mathbb{C}}, \text{ and } \tilde{q} \neq 1_\alpha \end{array} \right\}$

objects used to analyze maps between Reedy diagrams: $\underset{\partial(F^n\mathcal{C} \downarrow \alpha)}{\text{colim}} X$ and $\underset{\partial(\alpha \downarrow F^n\mathcal{C})}{\lim} X$

Key fact: All colim are latching objects of X

All lim are matching objects of X .

$$X : F^n\mathcal{C} \rightarrow M$$

$$\begin{array}{c} [1] \rightarrow [2], \\ [0] \rightarrow [2] \end{array}$$

If $\alpha \in \mathcal{C}$ has degree n , $i^n : F^{n-1}\mathcal{C} \rightarrow F^n\mathcal{C}$ the inclusion functor. $M_\alpha X =$

(1) The LATCHING CATEGORY $\partial(\tilde{\mathcal{C}} \downarrow \alpha)$ is a right cofinal subcat. of both $(i^n \downarrow \alpha)$ and $\partial(F^n\mathcal{C} \downarrow \alpha)$

(2) The MATCHING CATEGORY $\partial(\alpha \downarrow \tilde{\mathcal{C}})$ is a left cofinal subcat. of both $(\alpha \downarrow i^n)$ and $\partial(\alpha \downarrow F^n\mathcal{C})$

$\mathcal{C}, M, \alpha \in \mathcal{C}$ as above.

Proposition Let $X : \mathcal{C} \rightarrow M$ (a \mathcal{C} -diagram in M), $i^n : F^{n-1}\mathcal{C} \rightarrow F^n\mathcal{C}$ the inclusion functor. There are natural isomorphisms:

$$\begin{aligned} L_\alpha X &\cong \underset{(i^n \downarrow \alpha)}{\text{colim}} X \\ &\cong \underset{\partial(F^n\mathcal{C} \downarrow \alpha)}{\text{colim}} X \end{aligned}$$

And

$$\begin{aligned} M_\alpha X &\cong \underset{(\alpha \downarrow i^n)}{\lim} X \\ &\cong \underset{\partial(\alpha \downarrow F^n\mathcal{C})}{\lim} X \end{aligned}$$

Summary: What do we have?

- \mathcal{C} Reedy, M a model category
- $X: F^{n \rightarrow \mathcal{C}} \rightarrow M$ diagram indexed by $(n \rightarrow)$ -filtration of \mathcal{C} .
- $\alpha \in \mathcal{C}$ with $\deg \alpha = n$.

\rightsquigarrow There is a natural map $L_\alpha X \rightarrow M_\alpha X$

\rightsquigarrow Extending X to a diagram $F^n \mathcal{C} \rightarrow M$ is equivalent to the following:

For each $\alpha \in \mathcal{C}$ of $\deg \alpha = n$, choose an object X_α and a factorization

$$\underline{L_\alpha X} \rightarrow \underline{X_\alpha} \rightarrow \underline{M_\alpha X}$$

of the natural map.

$$\begin{array}{ccc} & X_\alpha & \\ \nearrow & \downarrow & \searrow \\ \text{colim } X & & \lim X \\ (\iota^n \downarrow \alpha) & & (\alpha \downarrow i^n) \end{array}$$

$$\begin{array}{ccc} & X_\alpha & \\ \nearrow & \downarrow & \searrow \\ L_\alpha X & \rightarrow & M_\alpha X \end{array}$$

- This can be done independently for each object of $\deg \alpha = n$

§ 5. Maps between Reedy diagrams

Same situation as above: $\mathcal{C}, M, \alpha \in \mathcal{C}, \deg \alpha = n$

- $X, Y: \mathcal{C} \rightarrow M$ diagrams

Goal: Define $X \Rightarrow Y$ inductively on filtrations.

Recall: $F^n \mathcal{C}$ contains no non-identity maps.

$\rightsquigarrow f: X|_{F^n \mathcal{C}} \Rightarrow Y|_{F^n \mathcal{C}}$ is completely determined by choosing a map $X_\alpha \rightarrow Y_\alpha$ for every object $\alpha \in \mathcal{C}$ with degree 0.

Suppose $f: X|_{F^{n-1} \mathcal{C}} \Rightarrow Y|_{F^{n-1} \mathcal{C}}$

For every $\alpha \in \mathcal{C}$ with $\deg \alpha = n$, we have:

$$\begin{array}{ccccc}
 \text{colim } X & \longrightarrow & X_\alpha & \longrightarrow & \lim_{(\alpha \downarrow i^n)} X \\
 (\iota^n \downarrow \alpha) & & & & \\
 \downarrow & & & & \downarrow \\
 \text{colim } Y & \longrightarrow & Y_\alpha & \longrightarrow & \lim_{(\alpha \downarrow i^n)} Y
 \end{array}$$

Proposition (*) {

$$\begin{array}{ccccc}
 Ld X & \longrightarrow & X_\alpha & \longrightarrow & Md X \\
 \downarrow & & \text{dashed arrow} & & \downarrow \\
 Ld Y & \longrightarrow & Y_\alpha & \longrightarrow & Md Y
 \end{array}$$

$\left\{ \text{Extensions of } f \text{ to } F^n C \right\} \xleftrightarrow{\text{correspondence}} \left\{ \begin{array}{l} \text{choice, for each } \alpha \in C \text{ with } \deg n, \text{ of a} \\ \text{dashed arrow} \end{array} \right. \text{ that makes squares commute} \left. \right\}$

Suppose $A, B, X, Y : \mathcal{C} \rightarrow M$ are Reedy diagrams and the following commutes:

(1)

$$\begin{array}{ccc}
 A & \longrightarrow & X \\
 \downarrow & \nearrow h & \downarrow \\
 B & \longrightarrow & Y
 \end{array}$$

where h is only defined
on $B|_{F^n C}$

Then, for every $\alpha \in C$ with $\deg n$, we have an induced diagram:

$$\begin{array}{ccc}
 A_\alpha \underset{L_\alpha A}{\amalg} L_\alpha B & \longrightarrow & X_\alpha \\
 \downarrow & \nearrow h_\alpha & \downarrow \\
 B_\alpha & \longrightarrow & Y_\alpha \underset{M_\alpha Y}{\times} M_\alpha X
 \end{array}$$

(2)

$\left\{ h \text{ can be extended to } F^n C \text{ so that (1) commutes} \right\}$

\uparrow correspondence

$\left\{ \text{for every } \alpha \in C \text{ of deg } n, \text{ there is a map } B_\alpha \xrightarrow{h_\alpha} X_\alpha \text{ so that (2) commutes} \right\}$

- Have the ability to define the Reedy model structure.
 \Rightarrow statement and proof next week featuring Sayantan!