

# Types of Algorithms

# Types of Algorithms

- Brute force Algorithm
  - Recursive Algorithm
  - Backtracking Algorithm
  - Divide and Conquer Algorithm
  - Greedy Algorithm
  - Dynamic Programming Algorithm
- 
- Searching Algorithm
  - Sorting Algorithm

# Brute force Algorithm

- Check all possible ways
- All possible solution
- Pros
  - Guaranteed way to find the correct solution
  - Ideal for solving small and simpler problems
- Cons
  - Inefficient
  - Often goes above  $O(n!)$
- Find smallest and largest  $n$  digit numbers with sum digits equal to  $k$

# Recursive Algorithm

- Function calls itself directly or indirectly
- Calling a copy of itself and solving smaller sub-problems
  - Breaking the problem into sub problems of the same type
  - Base condition
- Properties of Recursion
  - Performing the same operations multiple times with different inputs
  - In every step we solve smaller inputs and smaller problem
  - Base condition to stop recursion
- Factorial, Fibonacci, Tower of Hanoi, DFS, ...

# Backtracking Algorithm

- Solving the problem in an incremental way
  - Recursively
- Build solution one piece at a time
- Remove solutions that fail to satisfy the constraints
  
- Hamiltonian Cycle
- Sudoku
- N queen problem

# Divide and Conquer Algorithm

- Solve the problem in three steps
  - Divide: Divide the problem into sub problems
  - Conquer: Solve sub problems independently
  - Combine: Add the combined results

# Greedy Algorithm

- Solution is built part by part
- Based on immediate benefit
- Make local optimal choice and hope of finding a global optimum
- Problem properties:
  - Greedy choice property: Choosing the best solution in each phase leads to global optimal
  - Optimal substructure: optimal solution contains the optimal solution to the subproblems
- Dijkstra, Prim, Kruskal, ...

# Example for greedy proof (prim)

- Greedy solution is always part of some optimum solution
  - $T'$  is a subtree of a MST for  $G$ . then there is a complete MST for  $G$  that contains  $T'$  and contains the edge  $(u, v)$  that is the minimum weight edge from a vertex in  $T'$  to a vertex not in  $T'$
  - Consider  $T^*$  (MST with  $T'$  and without  $(u, v)$ )
  - Consider the unique path between  $u$  and  $v$
  - The unique path leaves  $T'$  with some edge  $(y, z)$ 
    - $y$  is in  $T'$  and  $z$  isn't
  - Remove  $(y, z)$  and add  $(u, v)$ 
    - New graph is a tree
    - $\text{weight}(u, v) \leq \text{weight}(y, z)$



# Dynamic Programming Algorithm

- Optimization over recursion
  - Recursion with repeated calls for the same input
- Store the results of sub problems
- Prevent re computing the same solution
- Reduces time complexity
  - Recursive Fibonacci:  $O(2^n)$
  - DP Fibonacci:  $O(n)$

# Sorting algorithms

- Selection sort
- Bubble sort
- Insertion sort
- Merge sort
- Heap sort
- Counting sort
- ...

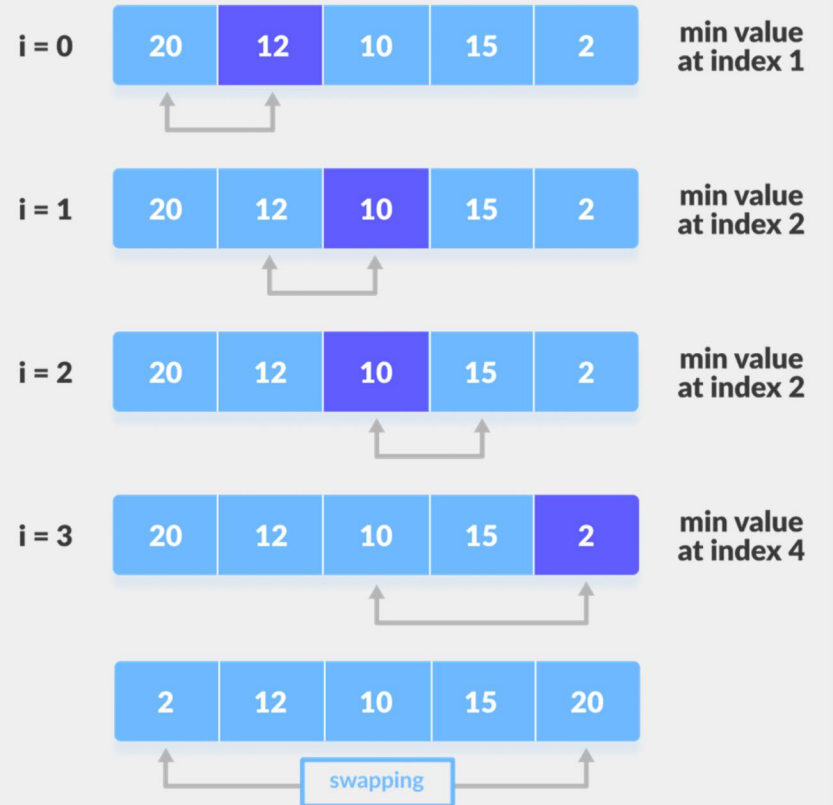
# Selection sort

- Each time find the current minimum
- Swap minimum with first element
- Time complexity:  $O(n^2)$

```
for i in range(len(A)):
    # Find the minimum element in remaining
    # unsorted array
    min_idx = i
    for j in range(i + 1, len(A)):
        if A[min_idx] > A[j]:
            min_idx = j

    A[i], A[min_idx] = A[min_idx], A[i]
```

step = 0



# Bubble sort

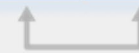
- Compare adjacent elements
- Swap if the order is incorrect
- Time complexity:  $O(n^2)$

```
for i in range(len(A)):  
    for j in range(0, len(A) - i - 1):  
        if A[j] > A[j + 1]:  
            A[j], A[j + 1] = A[j + 1], A[j]
```

step = 0

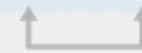
i = 0

-2	45	0	11	-9
----	----	---	----	----



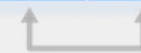
i = 1

-2	45	0	11	-9
----	----	---	----	----



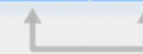
i = 2

-2	0	45	11	-9
----	---	----	----	----



i = 3

-2	0	11	45	-9
----	---	----	----	----



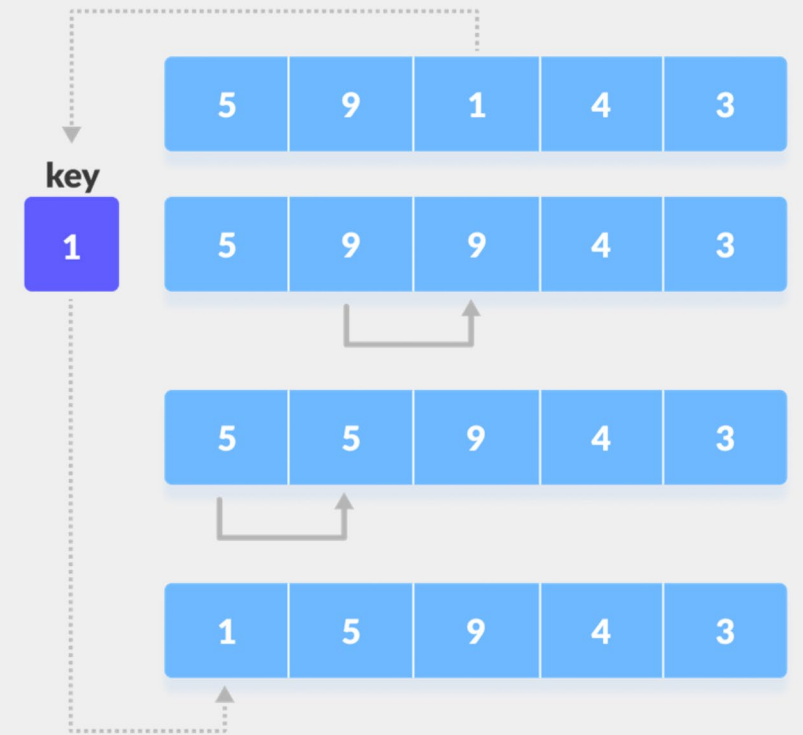
-2	0	11	-9	45
----	---	----	----	----

# Insertion sort

- Each time insert l'th element in sorted first part of the array
- Finds new data location and moves data to its place
- Shift all other cells
- Time complexity:  $O(n^2)$

```
for step in range(1, len(A)):  
    key = A[step]  
    j = step - 1  
  
    while j >= 0 and key < A[j]:  
        A[j + 1] = A[j]  
        j = j - 1  
  
    A[j + 1] = key
```

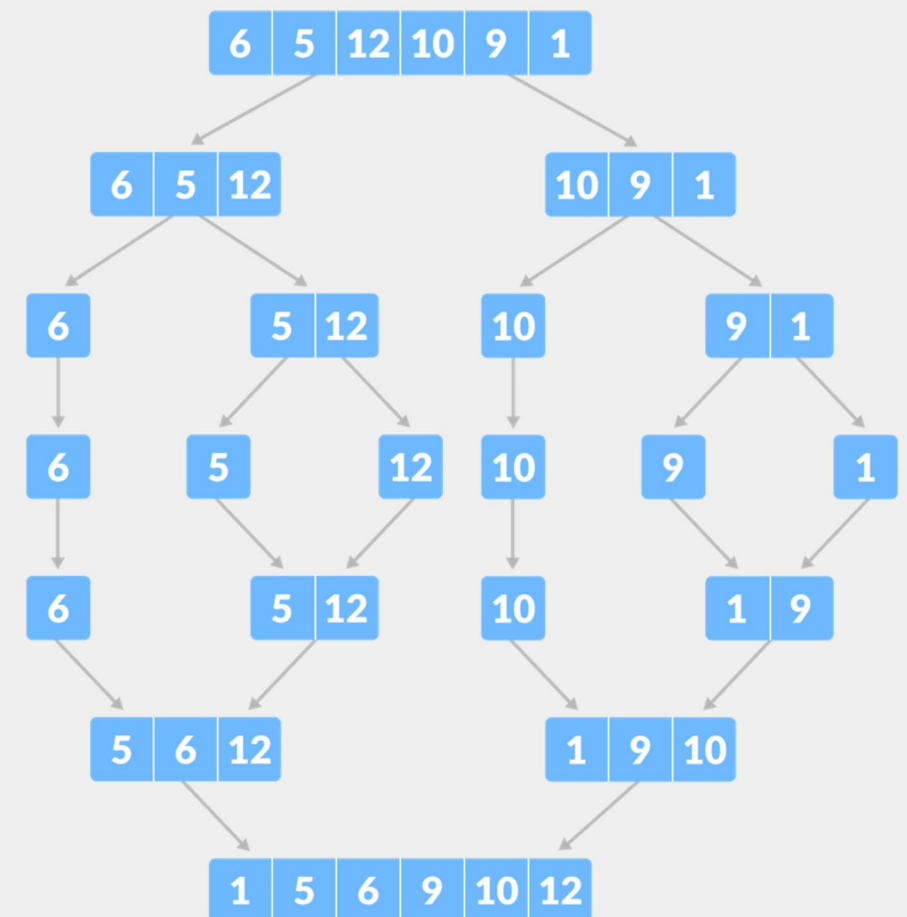
step = 2



# Merge sort

- Split array into two halves
- Sort each half recursively
- Merge sorted arrays
- Time complexity:  $O(n \log n)$

```
def mergeSort(array):  
    if len(array) <= 1:  
        return array  
  
    mid = len(array)//2  
  
    L = mergeSort(array[:mid])  
    R = mergeSort(array[mid:])  
    return combine_sorted(L, R)
```



# Merge sort

```
def mergeSort(array):  
    if len(array) <= 1:  
        return array  
  
    mid = len(array)//2  
  
    L = mergeSort(array[:mid])  
    R = mergeSort(array[mid:])  
    return combine_sorted(L, R)
```

```
def combine_sorted(L, R):  
    i = j = 0  
    array = []  
    while i < len(L) and j < len(R):  
        if L[i] < R[j]:  
            array.append(L[i])  
            i += 1  
        else:  
            array.append(R[j])  
            j += 1  
  
    return array + L[i:] + R[j:]
```

# Counting sort

- Not based on comparison
- Find the maximum element
- Generate a new list of zeros with  $\text{len} = \text{max}(\text{Array})$
- Count each value
- Time complexity =  $O(n + m)$
- Linear complexity
- More memory usage

```
def countingSort(array):  
    count = [0] * max(array)  
    res = []  
  
    for value in array:  
        count[value] += 1  
  
    for i in range(len(count)):  
        res = res + ([i] * count[i])
```



# honorable mentions

- Quick sort
  - Divide the array based on a pivot
- Pancake sort
  - Sort only by flipping array from 0 to 1
  - Find max, flip twice to place it at the end
- Bogo sort
  - Random permutation until success
- Sleep sort
  - Generate new thread for each value and sleep it for time = value
  - Print the value of each thread after sleep