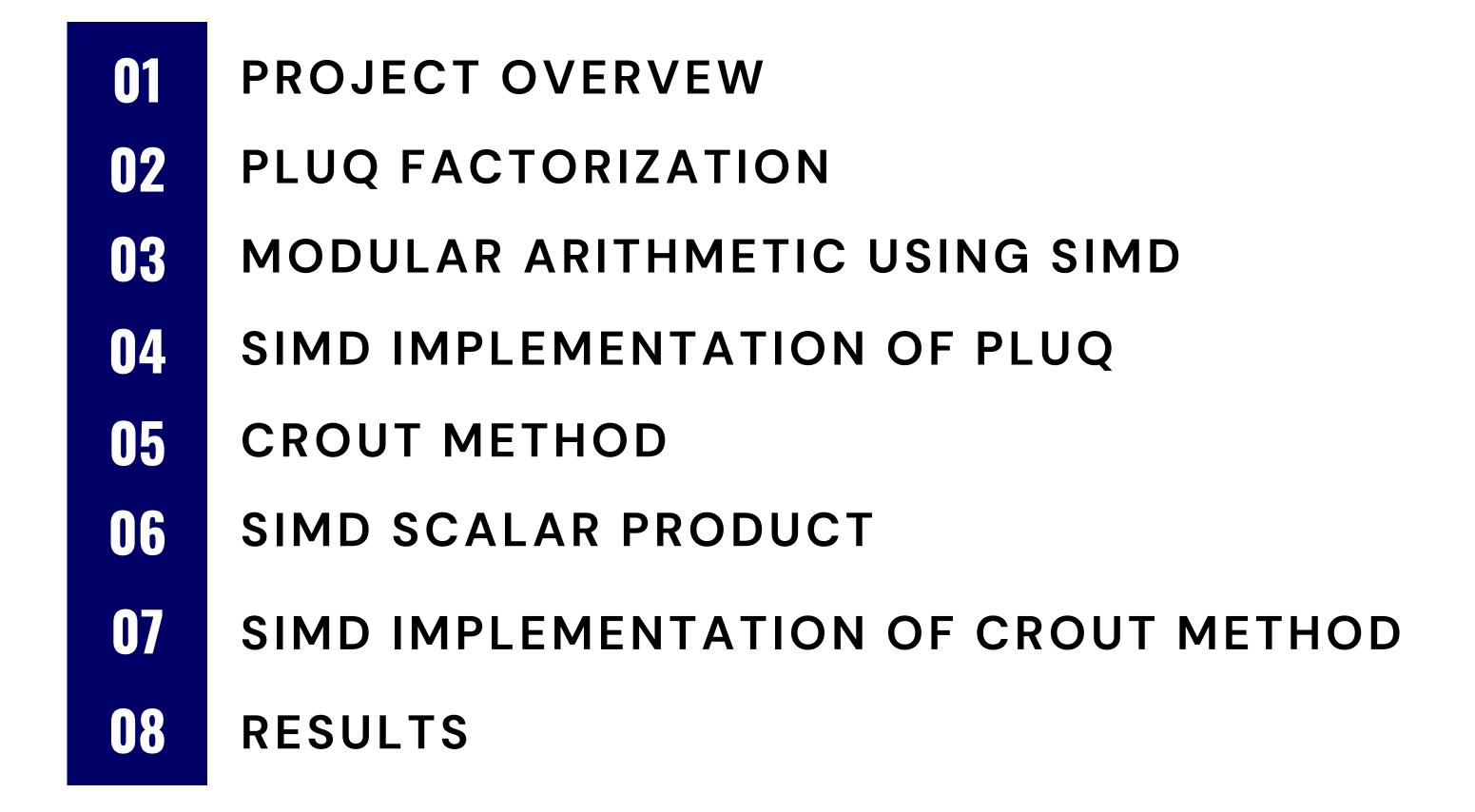
Gaussian Elimination High-Performance Implementation

Mohamed Imad Eddine Ghodbane May, 2024



CONTENT



PROJECT OVERVEW

- Gaussien elimination over a finite field Fp.
- p prime, less than 30 bits in length.
- Libraries include FFLAS-FFPACK, Flint known for high performance.
- Room for improvement for matrices of intermediate dimensions.
- AVX2 vectorization for enhanced performance.
- Aim for superior performance compared to existing libraries.

Input:

• Matrix A of size m×n with entries in the field Fp.

Output:

LU decomposition of A, with permutation matrices P and Q.

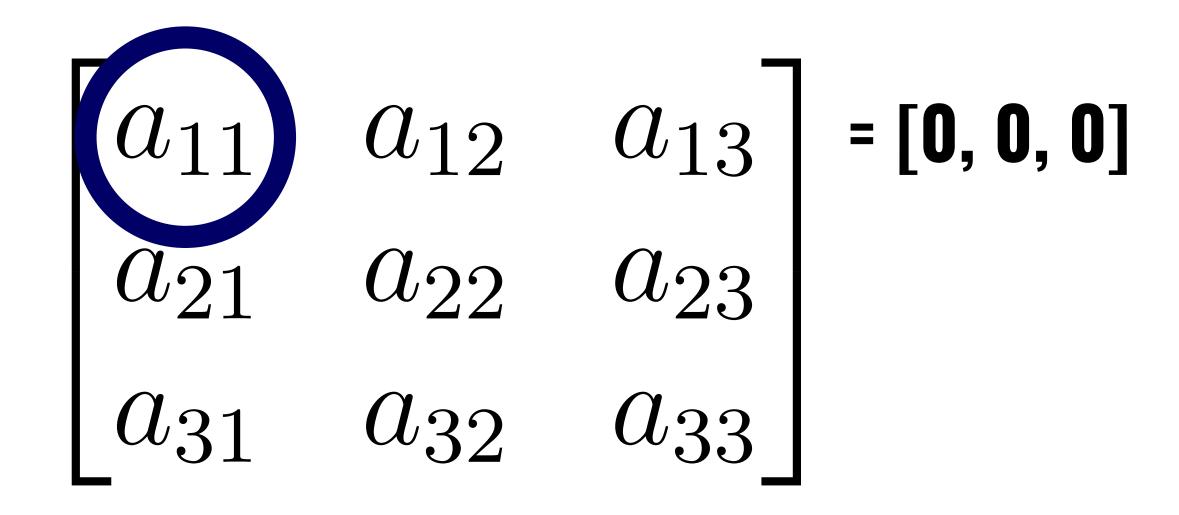
Operations:

- Pivoting: Involves row and column rotations.
- Row Reduction: Process of reducing rows to obtain the LU decomposition.

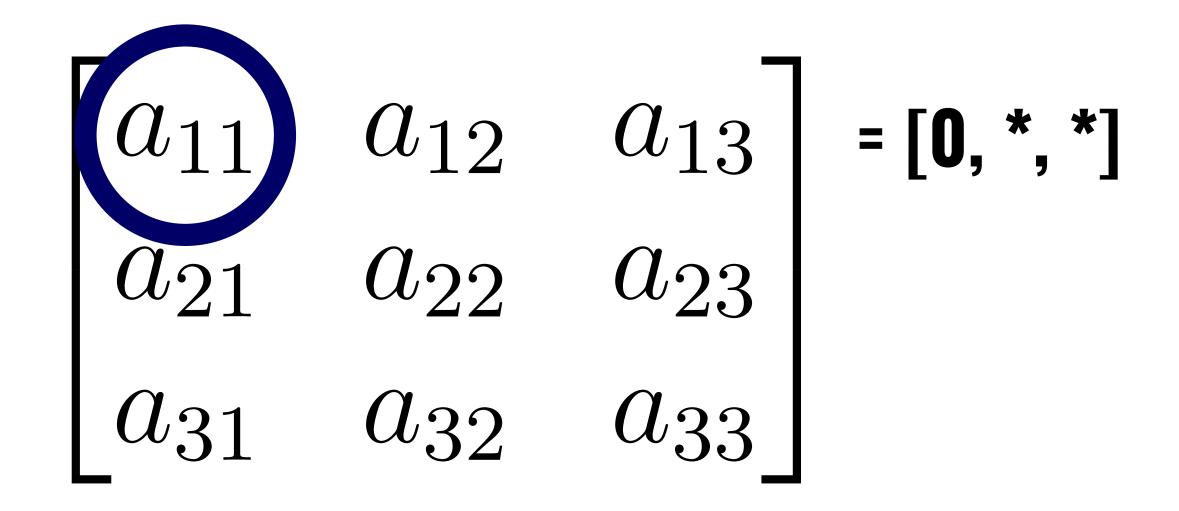
Input

$$m = n = 3$$

$$A = \begin{bmatrix} a_{11} & a_{12} & a_{13} \\ a_{21} & a_{22} & a_{23} \\ a_{31} & a_{32} & a_{33} \end{bmatrix}$$



$$egin{bmatrix} a_{21} & a_{22} & a_{23} \ a_{31} & a_{32} & a_{33} \ a_{11} & a_{12} & a_{13} \ \end{bmatrix}$$



$$egin{bmatrix} a_{21} & a_{11} & a_{13} \ a_{22} & a_{21} & a_{23} \ a_{32} & a_{31} & a_{33} \end{bmatrix}$$

Row Reduction

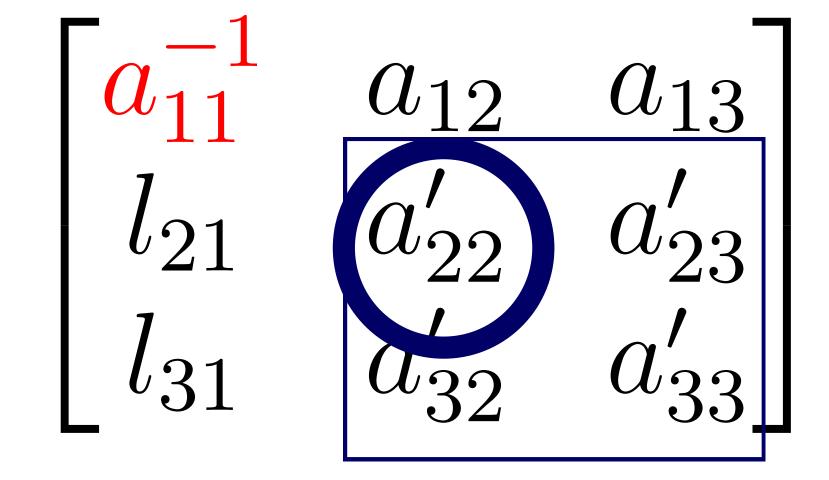
$$a_{11}^{-1} \cdot \begin{pmatrix} a_{21} \\ a_{31} \end{pmatrix} = \begin{bmatrix} a_{11}^{-1} & a_{12} & a_{13} \\ l_{21} & a_{22} & a_{23} \\ l_{31} & a_{32} & a_{33} \end{bmatrix}$$

Row Reduction

$$egin{bmatrix} a_{11} & a_{12} & a_{13} \ l_{21} & a_{22} & a_{23} \ l_{31} & a_{32} & a_{33} \ \end{bmatrix} = egin{bmatrix} a_{22} \ a_{23} \ a_{33} \ \end{bmatrix} = egin{bmatrix} a_{32} \ a_{33} \ \end{bmatrix}$$

$$= \begin{pmatrix} a_{22} \\ a_{23} \end{pmatrix} - l_{21} \cdot \begin{pmatrix} a_{12} \\ a_{13} \end{pmatrix}$$
$$= \begin{pmatrix} a_{32} \\ a_{33} \end{pmatrix} - l_{31} \cdot \begin{pmatrix} a_{12} \\ a_{13} \end{pmatrix}$$

Row Reduction



Output

$$P \cdot egin{bmatrix} u_{11} & u_{12} & u_{13} \ l_{21} & u_{22} & u_{23} \ l_{31} & l_{32} & u_{33} \end{bmatrix} \cdot Q$$

PLUQ IMPLEMENTATION

```
Pivoting
*/
// Row Reduction
for (int k = matrixRank + 1; k < m; k++) {
   A->data[k * n + matrixRank] = mult(
    A->data[k * n + matrixRank], inv, p);
    for (int j = matrixRank + 1; j < n; j++)
        A->data[k * n + j] = sub(
        A->data[k * n + j],
        mult(A->data[k * n + matrixRank],
        A->data[matrixRank * n + j], p), p);
```

Subtraction

```
int sub(int a, int b, int p) {
   int r = a - b;
   return r < 0 ? r + p : r;
}</pre>
```

Subtraction

_mm_sub_epi32

 a_2 a_3 a_1 a_4 *b*₃ b_2 $a_1 - b_1 \mid a_2 - b_2 \mid a_3 - b_3 \mid a_4 - b_4$

Subtraction

mm_cmplt_epi32

Subtraction

_mm_add_epi32

Subtraction __mm_blendv_epi8

Subtraction

```
__m128i sub_avx2(__m128i a, __m128i b, __m128i vp) {
    __m128i result = _mm_sub_epi32(a, b);
    __m128i mask = _mm_cmplt_epi32(result, _mm_setzero_si128());
    __m128i adjusted_result = _mm_add_epi32(result, vp);
    result = _mm_blendv_epi8(result, adjusted_result, mask);
    return result;
}
```

```
int mult(int a, int b, int p) {
   long long r = (long long)a * b;
   return r % p;
}
```

$$h = ab = cp + e$$
 h_1 h_2 h_3 h_4
 $u = \frac{1}{p}, d = hu$
 u_1 u_2 u_3 u_4
 u_4
 u_1 u_2 u_3 u_4
 u_4

$$h = ab = cp + e$$

$$u = \frac{1}{p}, d = hu$$

$$c = \lfloor d \rfloor$$

$$\begin{bmatrix} d_1 \end{bmatrix} \begin{bmatrix} d_2 \end{bmatrix} \begin{bmatrix} d_3 \end{bmatrix} \begin{bmatrix} d_4 \end{bmatrix}$$
 $\begin{bmatrix} C_1 \end{bmatrix} \begin{bmatrix} C_2 \end{bmatrix} \begin{bmatrix} C_3 \end{bmatrix} \begin{bmatrix} C_4 \end{bmatrix}$

$$h = ab = cp + e$$
 $u = \frac{1}{p}, d = hu$
 $c = \lfloor d \rfloor$
 $e = h - cp$
 h_1
 h_2
 h_3
 h_4
 h_5
 h_6
 h_7
 h_8
 h_8
 h_8
 h_9
 h_9

Multiplication

```
__m256d mul_mod_p(__m256d a, __m256d b, __m256d u, __m256d p) {
    __m256d h = _mm256_mul_pd(a, b);
    __m256d d = _mm256_mul_pd(h, u);
    __m256d c = _mm256_floor_pd(d);
    __m256d e = _mm256_fnmadd_pd(c, p, h);
    return e;
}
```

The resulting product often exceeds the 52-bit length allocated for the mantissa in double floating-point representation. This can lead to a loss of precision in the final result!!

```
_m256d mul_mod_p(__m256d a, __m256d b, __m256d u, __m256d p) {
  m256d h = mm256 mul pd(a, b);
  m256d l = mm256 fmsub pd(x, y, h);
  _{m256d} d = _{mm256} mul_pd(h, u);
  m256d c = mm256 floor pd(d);
  _{m256d b} = _{mm256_fnmadd_pd(c, p, h);}
  _{m256d} = _{mm256_add_pd(b, 1);}
  _{m256d} t = _{mm256} sub_pd(e, p);
  e = _mm256 blendv_pd(t, e, t);
  t = _mm256 add pd(e, p);
  return mm256 blendv pd(e, t, e);
```

SIMD IMPLEMENTATION OF PLUQ

```
rows elimination avx2(int *A data, int n, int matrixRank, int c, int p,
__m256d vp, __m256d vu ,__m128i vp_128, int k) {
   _{m256d} vc = _{mm256} set1_pd(c);
   m256d tmp;
   int i;
    for (i = matrixRank + 1; i + 3 < n; i += 4) {
       __m128i v1 = _mm_loadu_si128((__m128i *)&A_data[matrixRank*n+i]);
       __m128i v2 = _mm_loadu_si128((__m128i *)&A_data[k*n+i]);
       m256d vDouble = _mm256_cvtepi32_pd(v1);
       tmp = mul mod p(vc, vDouble, vu, vp);
        _m128i resultInt = _mm256_cvttpd_epi32(tmp);
        m128i result = sub avx2(v2, resultInt, vp 128);
       _mm_storeu_si128((__m128i *)&A_data[k * n + i], result);
    } // loop handles elements that don't fit into chunks of 4
```

RESULTS

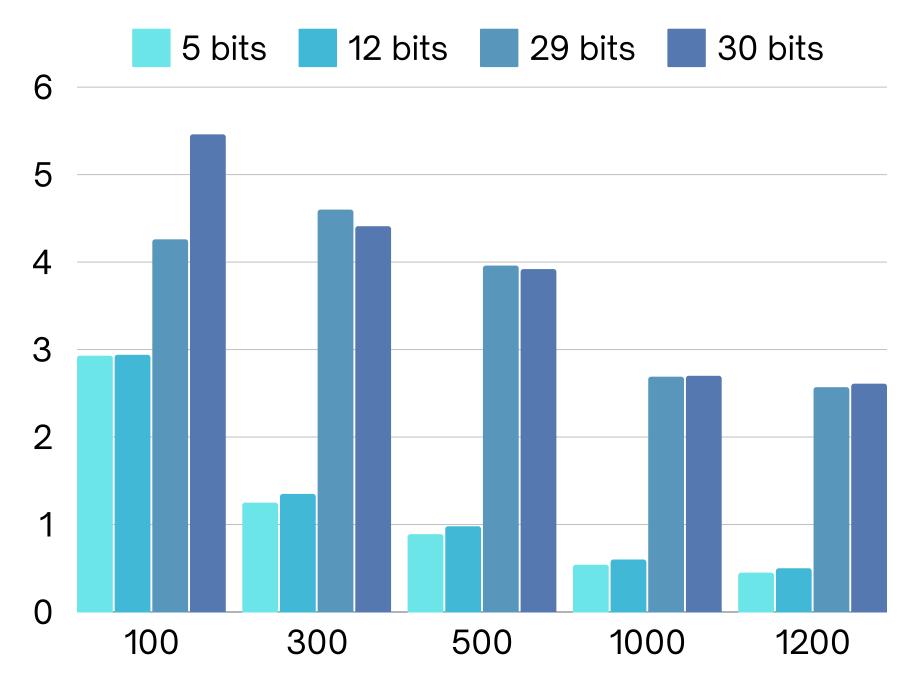
AVX2 vs Basic PLUQ

| Sizes | 100 | 300 | 500 | 1000 | 1200 |
|------------|------|-------|-------|--------|---------|
| Basic (ms) | 0.69 | 18.70 | 86.57 | 694.06 | 1199.92 |
| AVX2 (ms) | 0.15 | 3.90 | 17.80 | 139.98 | 242.56 |
| Speedup | 4.60 | 4.79 | 4.86 | 4.95 | 4.95 |

PLUQ Speedups using AVX2 and 12 Bits Length Prime

RESULTS

AVX2 vs FLINT



AMD Ryzen™ 7 PRO 7840U w/ Radeon™ 780M Graphics × 16