

Hierarchical Leader Election Algorithm With Remoteness Constraint

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Outline

1 Introduction

- What's Distributed Systems ?
- Election Algorithms

2 Preliminaries

- System Model
- Modeling Asynchronous Dynamic Links
- Configurations and Executions
- Problem Definition

3 Leader Election Algorithm

- Informal Description
- Nodes, Neighbors and Heights
- Initial State

■ Goal Of The Algorithm

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What's Distributed Systems ?

A distributed system is a network that consists of autonomous computers that are connected using a distribution middleware. They help in sharing different resources and capabilities to provide users with a single and integrated coherent network.

The Bully Algorithm

As a first example, consider the bully algorithm devised by Garcia-Molina (1982). When any process notices that the coordinator is no longer responding to requests, it initiates an election. A process, P , holds an election as follows:

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- P sends an *ELECTION* message to all processes with higher numbers.
- If no one responds, P wins the election and becomes coordinator.
- If one of the higher-ups answers, it takes over. P 's job is done.

Election Algorithms

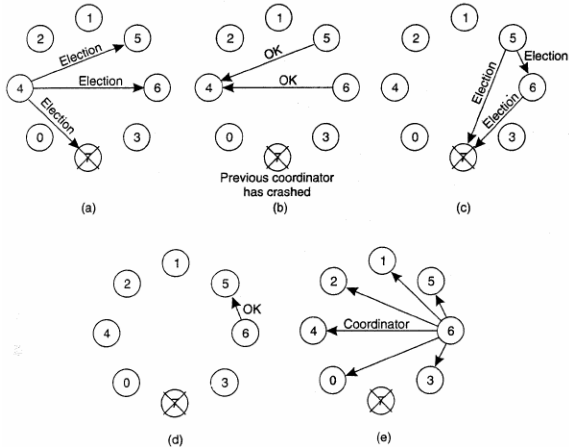


Figure: Bully Algorithm

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- We model the whole system as a set of (infinite) state machines that interact through shared events (a specialization of the IOA model [17]).

Asynchronous Dynamic Links' Model

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- and a queue mqueue_{uv} of messages.

Configurations & Executions

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- A configuration is a vector of node states, one for each node in P , and a vector of channel states, one for each channel in χ .

Problem Definition

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Each node u in the system has :

- a local variable lid_u to hold the id of the supreme leader;
- another local variable $slid_u$ to hold the identifier of the sub-leader whose remoteness towards u obeys the constraint.

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- First wave : initiated by one of the lost leader's neighbors looking for it;
- Second wave : initiated by the node located at the edge of the network if the search has hit a dead-end;
- Third wave : initiated by the same node which initiated the first wave updating the other nodes' heights.

Heights

The height for each node is a 7-tuple of integers

$((\tau, oid, r), \delta, (nlts, lid), id)$, where the first three components are referred to as the reference level (RL) and the fifth and sixth.

- τ , a non-negative timestamp which is either 0 or the value of the causal clock time when the current search for an alternate path to the leader was initiated.

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The Code triggered by Update Message

When node u receives $Update(h)$ from node $v \in forming \cup N$:

```

// if  $v$  is in neither forming nor  $N$ , message is ignored
1.  $height[v] := h$ 
2.  $forming := forming \setminus \{v\}$ 
3.  $N := N \cup \{v\}$ 
4.  $myOldHeight := height[u]$ 
5. if  $((nls^u, lid^u) = (nls^v, lid^v))$  // leader pairs are the same
6.   if (SINK)
7.     if  $(\exists (\tau, oid, r) \mid (\tau^w, oid^w, r^w) = (\tau, oid, r) \ \forall w \in N)$ 
8.       if  $((\tau > 0) \text{ and } (r = 0))$ 
9.         REFLECTREFLEVEL
10.      else if  $((\tau > 0) \text{ and } (r = 1) \text{ and } (oid = u))$ 
11.        ELECTSELF
12.      else //  $(\tau = 0)$  or  $(\tau > 0 \text{ and } r = 1 \text{ and } oid \neq u)$ 
13.        STARTNEWREFLEVEL
14.      end if
15.    else // neighbors have different ref levels
16.      PROPAGATELARGESTREFLEVEL
17.    end if
18.    // else not sink, do nothing
19.  else // leader pairs are different
20.    ADOPTLPIFPRIORITY( $v$ )
21.  end if
22. if  $(myOldHeight \neq height[u])$ 

```

Subroutines

ELECTSELF

1. $height[u] := (0, 0, 0, 0, -\mathcal{T}_u, u, u)$

REFLECTREFLEVEL

1. $height[u] := (\tau, oid, 1, 0, nlts^u, lid^u, u)$

PROPAGATELARGESTREFLEVEL

1. $(\tau^u, oid^u, r^u) := \max\{(\tau^w, oid^w, r^w) \mid w \in N\}$

2. $\delta^u := \min\{ \delta^w \mid w \in N \text{ and } (\tau^u, oid^u, r^u) = (\tau^w, oid^w, r^w) \} - 1$

STARTNEWREFLEVEL

1. $height[u] := (\mathcal{T}_u, u, 0, 0, nlts^u, lid^u, u)$

ADOPTLPIFPRIORITY(v)

1. if $((nlts^v < nlts^u) \text{ or } ((nlts^v = nlts^u) \text{ and } (lid^v < lid^u)))$

2. $height[u] := (\tau^v, oid^v, r^v, \delta^v + 1, nlts^v, lid^v, u)$

3. else send Update($height[u]$) to v

4. end if

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Where can I learn more?

Questions and Answers

Want to know more?

- Browse <http://web.mit.edu/smoot/history.htm>.

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- Smoot's Legacy http://alum.mit.edu/news/AlumniNews/Archive/smoots_legacy.

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- Smoot Salute!
<http://web.mit.edu/spotlight/smoot-salute>.