OS Lab 9

Sara Mohamed 900203032 Omar Harb 900201063 Mohamed Shaalan 900201539

Roles

Omar Harb: threads_compute, bonus functions, results and graphs

Mohamed Shaalan: mmap_compute, bonus functions, results and graphs

Sara Mohamed: threads_compute, bonus functions, results and graphs

Inputs and Outputs of mmap

Inside out mmap_compute() function, we initialize the output "result" using the mmap() function

- *addr = NULL, so kernel can place mapping wherever it sees fit
- length = sizeof(int), as result is a shared
 int
- prot = PROT_READ | PROT_WRITE, so all processes sharing this space can read/write
- flags = MAP_SHARED | MAP_ANONYMOUS, so all processes mapped to this object share the space; and so that the mapping is anonymous, i.e. not connected to any files
- fd = 0
- offset = 0, to start from the beginning of the file open on the fd descriptor

pipes_compute():

```
double pipes compute(int (*f)(int, int),
char* filename, int n proc)
    if (n proc == 0)
                                            int fdesc arr[n proc][2];
                                            int nperprocess = N/n proc;
        exit;
                                            int extraN = nperprocess + (N%n proc);
    int N = getcountinfile(filename);
    int* arr = getArray(filename, N);
                                            for (i=0 to n proc-1)
    start clock;
                                                if (pipe(fdesc_arr[i]) == -1)
    pid t parentid = getpid();
                                                    print error message;
                                                    exit(1); // initializing pipes, one per
    int count = 0, index = 0;
                                                              child process
    int tempread, tempwrite;
```

```
while (fork()!=-1 && count < n proc)
    if (getpid() != parent id)
         close(fdesc arr[count][0]);  // closes read side of pipe
         if (count == 0) // if first element, get result using index and extraN
              tempwrite = childResult(index, index + extraN - 1, f);
              write(fdesc arr[count][1], &tempwrite, sizeof(tempwrite));
              close(fdesc arr[count][1]); // close write side after done
                            // else, get result using index and nperprocess
         else
              tempwrite = childResult(index, index + nperprocess - 1, f);
              write(fdesc arr[count][1], &tempwrite, sizeof(tempwrite));
              close(fdesc arr[count][1]); // close write side after done
         break;
    else
         if (count == 0)
              index += extraN;
         else
              index += nperprocess;
                                                 child that will use that index
         count++;
```

This branch eliminates the need for another pipe- the parent increments the index properly before forking the

```
pipes_compute():
       if (getpid() != parent id)
           exit; // children exit program after their computations
      count = 0;
       if (getpid() == parent id)
           close(fdesc arr[count][1]); // close write side of first pipe
           read(fdesc arr[count][0], &tempread, sizeof(tempread));
           result = tempread;
           count++; // read first result, store in tempread, increment count
           while (count < n proc)
                close(fdesc arr[count][1]);
                read(fdesc arr[count][1], &tempread, sizeof(tempread));
                result = f(result, tempread);
                count++; // read remaining results, combine together
      end clock;
      print result;
       free arr;
      return (end clock - start clock);
```

sequential_compute():

```
double sequential compute (int
(*f)(int, int), char* filename)
    int. N =
                                              else
getcountinfile(filename);
                                                  result = arr[0];
    int* arr = getArray(filename,
                                                  for (i = 1 \text{ to } N-1)
N);
                                                       result = f(result, arr[i]);
    start clock;
                                                  stop clock;
    int result = 0;
                                              free arr;
    if (N==0)
                                              print result to screen;
         result = 0;
         stop clock;
                                              return (end clock - start clock);
    else if (N==1)
         result = arr[0];
         aton aloak.
```

mmap_compute():

```
const int N = getcountinfile(filename);
    int* arr = getArray(filename, N);
    initialize result mmap;
    if(result == MAP FAILED) {
        printf("map failed!\n");
        return 0.0;
    *result = 0;
    Start clock;
    pid t parent id = getpid();
```

```
int count = 0, index = 0;
int nperprocess = N / n_proc;
int extraN = nperprocess + (N % n_proc);
```

//identical to pipes_compute(), we find the number of elements each child process will compute the result of

```
while (fork() != -1 \&\& count < n proc)
        if (getpid() != parent id)
             if (count == 0) // if first element, get result using index and extraN
                   *result = f(*result, childResult(arr, index,
                   index+extraN-1, f)); //add to result
             else //else, get result using index and nperprocess
                   *result = f(*result, childResult(arr, index,
                   index+nperprocess-1, f));
              break;
        else
             if(count == 0)
                 index += extraN;
             else
                 index += nperprocess;
             count++;
```

```
mmap_compute():
```

```
if(getpid() != parent id)
        exit; //children exit
    else
        wait for all children to exit;
    end clock;
    print result;
    int err = munmap(result, sizeof(int)); //unmap shared memory
    if(err != 0) {
        printf("unmapping failed\n");
        return -1.0;
    return start clock - end clock;
```

Dependencies: thread_result(), thread_data struct, global variables

```
void* thread result(void* arg) {
                                                       // global variables
    thread data* td = (thread data*) arg;
                                                       int tresult = 0;
    if(td->sindex > td->eindex)
                                                       pthread mutex t mutex1;
        return NULL;
                                                       typedef struct {
                                                           int* arr;
    else
                                                           int sindex;
        pthread mutex lock(&mutex1);
                                                           int eindex;
                                                        } thread data;
        // critical section
        for(int i = td->sindex to td->eindex) {
            tresult = f(tresult, td->arr[i]);
        pthread mutex unlock(&mutex1);
```

thread_compute():

```
double single process thread compute(int (*f)(int, int), char* filename,
int n thread) {
    int N = getcountinfile(filename);
    int* arr = getArray(filename, N);
    start clock;
    int nperprocess = N / n thread;
    int extraN = nperprocess + (N % n thread);
    pthread t tid[n thread];
    thread data tdata[n thread]; //array of data associated to each thread
```

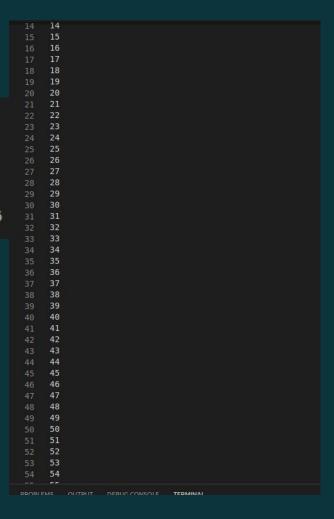
```
int error, index = 0;
    for (int i = 0 to n thread -1) {
        tdata[i].arr = arr; // "assign" each thread the array
        tdata[i].sindex = index; // "assign" start index = index
        if(i == 0)
            tdata[i].eindex = index+extraN-1;
            index += extraN;
        else
            tdata[i].eindex = index+nperprocess-1;
            index += nperprocess;
        error = pthread create(&(tid[i]), NULL, thread result, (void*)
&tdata[i]);
        if(error != 0)
            printf("Thread can't be created: [%s]\n", error);
            i--; revert index value; // try again since we have to create n_thread
threads
    for (int i = 0; i < n thread; i++)
                                           end clock, free arr, print result;
        pthread join(tid[i], NULL);
                                           return (end clock - start clock);
    // join all threads
```

Run Example

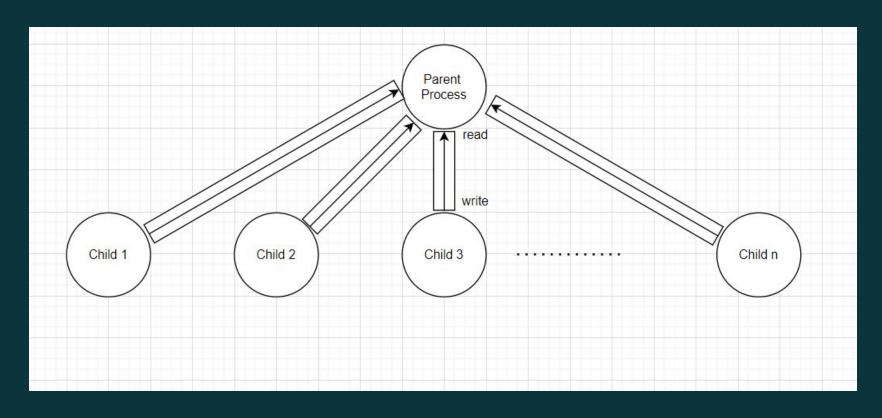
This run enters 500 elements (1, 2, 3, ..., 500) to file and then computes them using the four functions

result =
$$\sum_{i=1}^{n} i = 1 + 2 + 3 + ... + n = \frac{n(n+1)}{2}$$

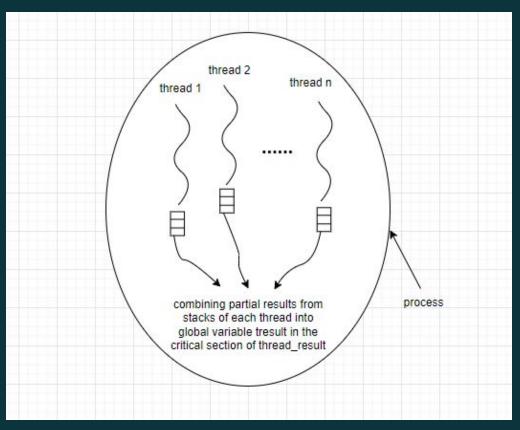
= 250(501) = 125250.



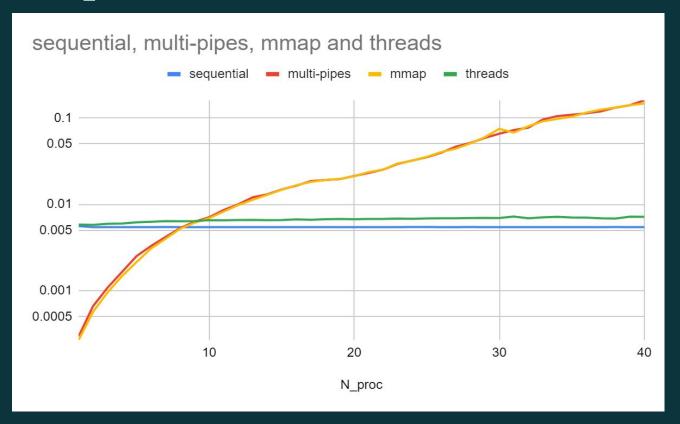
Diagrams: Pipes



Diagrams: Threads

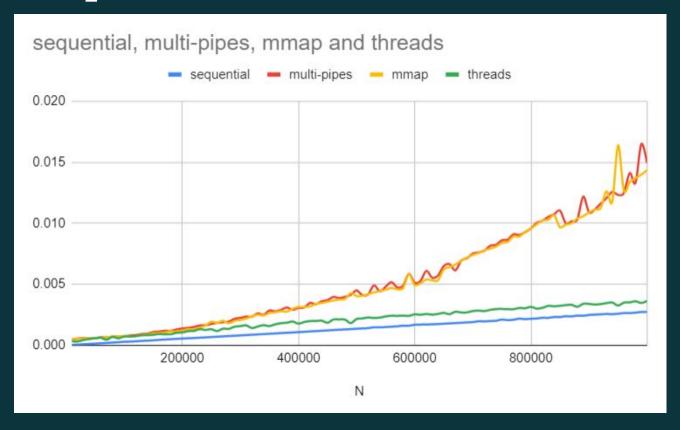


Graphs



The graph contains data with N fixed at 2M, and n_proc ranging from 1 to 40. Tests were conducted on a PC with 8 CPU cores. Intersection occurs at n_proc = 8. Graph presented in log scale for clearer differentiation of results.

Graphs

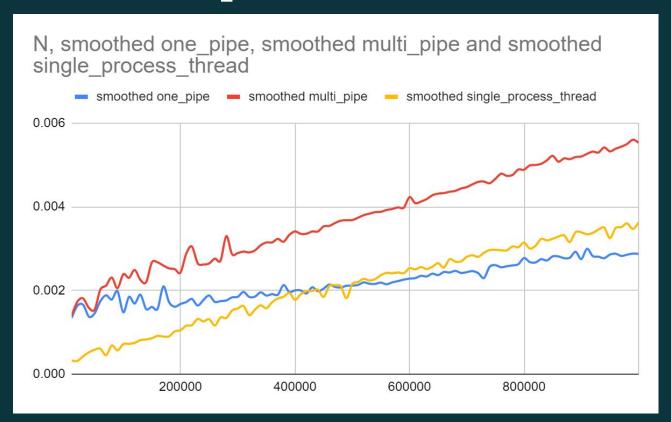


The graph contains data with n_proc and n_thread fixed at 8, and N ranging from 10K to 1M. Tests were conducted on a PC with 8 CPU cores.

Bonus Functions

- created multi_process_thread_compute, single_process_thread_compute,
 multi_pipe_compute, and one_pipe_compute functions (attached in zipped file)
- For comparison graph, fixed n_thread and n_proc at 8, as testing was conducted on a PC with 8 CPU cores
- Multi_process_thread_compute generated error, so did not include it in our generated graph (but still attached in zipped file)

Bonus Graph



The graph contains data with n_proc and fixed at 8, and N ranging from 10K to 1M. Tests were conducted on a PC with 8 CPU cores.

Ranges and Results

- Testing and runs were conducted on a PC in the Systems Engineering Lab with 8 CPU cores
- Smoothing of graphs was conducted by getting the median of three readings.
- Fixing n_proc and n_thread to 8, we see that sequential compute always takes the least time, as it has no overhead tied to it, and thread_compute() nears it with increasing values of N.
- Fixing N to 2M, we see that with n_proc and n_thread less than 8, sequential and thread_compute take longer time, but the behaviour reverses for values on n_proc and n_thread greater than 8.
- Ranges of tests chosen to both reflect full graph view while maintaining details and information for the smaller values.