# OS Lab 11

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## Roles

Omar Harb: Multi-Level Feedback Queue,

Sara Mohamed: Round Robin,

Mohamed Shaalan: FCFS,

### "Process" Struct Attributes

- PID
  - to keep track of the processes
- Burst Time
  - o needed for all scheduling algorithms, to properly move the time t
- Remaining Time
  - need to check when process gets done executing
- Arrival Time
  - o needed to keep track of what processes can be scheduled, according to the current time t
- Response
- Waiting
- Turnaround

## Dependencies

```
Process *get_n_processes(int n)
{
    Process *processes = new Process[n];
    for (i to n-1)
    {
        processes[i].pid = i + 1;
        processes[i].burst_time = rand(1-100);
        processes[i].arrival_time = rand(0-100);
        processes[i].remaining_time = processes[i].burst_time;
    }
    sort(processes);    //sort processes by arrival time
    return processes;
}
```

### Pseudocode: Round Robin

```
void getRRwaitingtime(Process processes[], int n, int waitingtime[], int turningtime[], int
responsetime[], int quantum)
    int remaining burstTimes[n];
    for (i = 0 \text{ to } n-1)
        remaining burstTimes[i] = processes[i].burst time;
    int t = 0; int finishedProcesses = 0; //initializing time t and num of finished processes
    for (;;) {
        int processesStarted = 0; //number of processes started in this for loop over the processes
        for (i = 0 \text{ to } n-1) {
             if (arrivaltime(i) <= t & remaining burstTimes(i) > 0) //if process arrived and not finished
                      if (first time process is running)
                          responsetime[i] = t;  //evaluation of process response time
                     processesStarted++; //increment num of processes started
                      if (remaining burstTimes[i] > quantum)
                          t += quantum;
                          remaining burstTimes[i] -= quantum;
```

### Pseudocode: Round Robin

```
// if remaining burst time less than/equal to quantum
                else
                    t = t + remaining burstTimes[i];
                    waitingtime[i] = t - processes[i].arrival time - processes[i].burst time;
                    remaining burstTimes[i] = 0;
                    finishedProcesses++;
     //end of inner for loop
  if (finishedProcesses == n) //if all processes done, break from the infinite for loop
       break;
   if (processesStarted == 0) //if none of the processes started in the past for loop, increment t
       t++;
for (int i = 0; i < n; i++) //turnaround time is simply wait time * burst time
  turnarountime[i] = waitingtime[i] + processes[i].burst time;
```

### Pseudocode: FCFS

```
void run fcfs(Process processes[], int n, int waitingtime[], int turnaroundtime[], int
responsetime[] ){
    currentTime = 0;
    waitingtime[0] = 0; // first process has no waiting time
    if (processes [0].arrival == 0) { //if first process arrives at time 0, add its burst time to currentTime
         currentTime = currentTime + processes[0].burst time;}
    else { // if first process doesn't immediately arrive, jump to its arrival time + burst time
         currentTime = currentTime + processes[0].arrival_time + processes[0].burst_time;}
    for (i to n-1)
        if (processes[i].arrival time < t) { // if process arrives before currentTime, it has to wait
             waitingtime[i] = t - processes[i].arrival time; //it waited from when it arrived till currentTime
             t += processes[i].burst time; }
         else { // if process arrives after or on currentTime, it has no waiting time
             waitingtime[i] = 0; //process did not wait
             t = processes[i].arrival_time + processes[i].burst_time;}
```

### Pseudocode: FCFS

### Pseudocode: Multi-Level Feedback Queue

```
void run mlq(vector<Proc Queue> &queues, processes[], mlq waiting time[], mlq turnaround time[],
mlq response time[], n)
    m = queues.size();
    t = \hat{0}; // current time, starting at 0
                                                                                  struct Proc Queue
    // all processes initially assigned to first queue.
    for (i = 0 \text{ to } n-1)
        queues[0].processes.push(processes[i]);
                                                                                       int time quantum;
    for (i = 0 \text{ to } m-1)
                                                                                       queue<Process> processes;
                                                                                  };
        if (i == m - 1 \text{ and } m > 1)
             sort queue(queues[i]);
        min arrival = INF;
        num moved = 0;
        while (!queues[i].processes.empty())
            &process = queues[i].processes.front();
            if (process.arrival time <= t)
                 // so it doesn't get overwritten when assigned to lower gueue
                 if (i == 0)
                     process.response = t - process.arrival time;
                 int quantum = min(process.remaining time, queues[i].time quantum);
                 t += quantum;
                process.remaining time -= quantum;
                 queues[i].processes.pop();
```

### Pseudocode: Multi-Level Feedback Queue

```
if (process done executing)
                process.turnaround = t - process.arrival time;
                process.waiting = process.turnaround - process.burst time;
             else
                if (i < m - 1)
                  push process into next queue
       else
             move process to end of queue
             min arrival = min(min arrival, process.arrival time);
             num moved++;
             if (num moved == queues[i].processes.size()) // process is now in original order
                 t = min arrival;
                 num moved = 0;
                 min arrival = INF;
get results (processes, mlq waiting time, mlq turnaround time, mlq response time, n);
```

## Sample Test Case

Process ID	Arrival Time	Burst Time
1	2	5
2	8	2
3	8	2
4	9	9
5	10	1

## **Testing Round-Robin**

Time-quantum = 4

#### **Gantt Chart:**

ı	X	P1		X	P2	P3	P4	P5	P1	P4	P4	
0		2	6	8	10	) 12	2 16	5 17	7 18	3 2	2 2	25

Processes	Burst time	Turn around time	Response time	Waiting time
1	5	5	0	0
2	2	2	0	0
3	2	4	2	2
4	9	13	3	4
0 5	1	7	6	6
DC C 111	24.00			VOC 1 1 1 1 1 44. []

## Testing FCFS

#### **Gantt Chart:**

ı	Χ	P1	ı	X	P2	P3	P4	P5	
0		2	7	8	10	) 12	2	1 2	2

Processes	Burst time	Turn around time	Response time	Waiting time
1	5	5	0	0
2	2	2	0	0
3	2	4	2	2
4	9	12	3	3
5	1	12	11	11

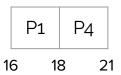
## Testing MLFQ

Time-quantum = {3, 5, INF}

#### Queue 1

	X	P1	X	P2	P3	P4	P5	
0	2	2 5	5 8	10	) 12	2 1	5 16	

#### Queue 2



#### Queue 3



#### Results (Computations)

PID	Turnaround	Response	Waiting
1	18 - 2 = 16	2 - 2 = 0	16 - 5 = 11
2	10 - 8 = 2	8 - 8 = 0	2 - 2 = 0
3	12 - 8 = 4	10 - 8 = 2	4 - 2 = 2
4	24 - 9 = 15	12 - 9 = 3	15 - 9 = 6
5	16 - 10 = 6	15 - 10 = 5	6 - 1 = 5

PS C:\User	s\Mohamed Sh	aalan\Desktop\Uni\	Course Material	\OS Lab\Lab11> ./main
Processes	Burst time	Turn around time	Response time	Waiting time
1	5	16	0	11
2	2	2	0	0
3	2	4	2	2
4	9	15	3	6
5	1	6	5	5

### Run Details

Processes: 100 - 2000

Arrival time range: **0 - 100** 

Burst time range: 1 - 100

RR time quantum: 20

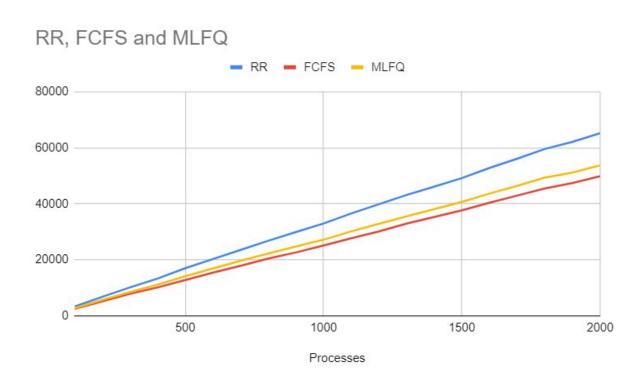
MLFQ Queue 1 time quantum: 3

MLFQ Queue 2 time quantum: 5

MLFQ Queue 3: FCFS

## Waiting Time Graph

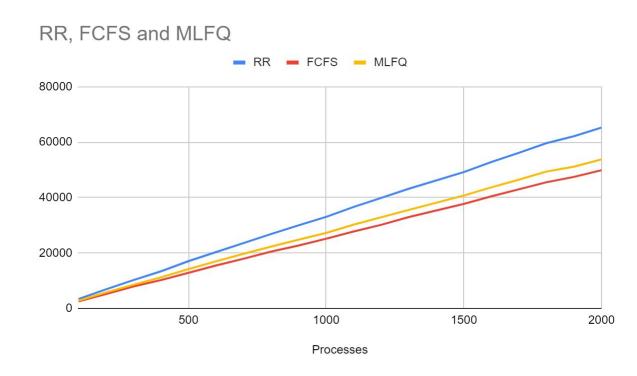
 understandably, rr takes most waiting time, followed by MLFQ and then FCFS



## Turnaround Time Graph

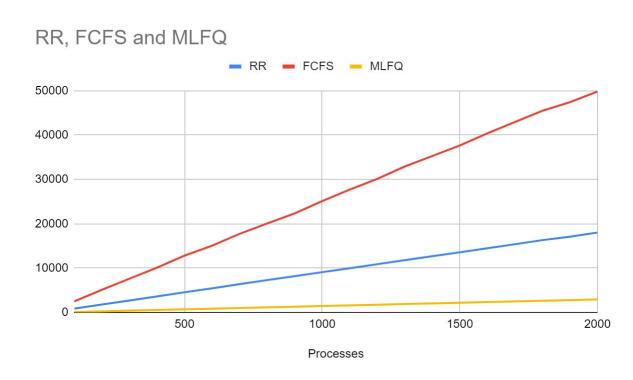
 understandably, rr takes most turnaround time, followed by MLFQ and then FCFS

 turnaround very similar to waiting, since it is evaluated by adding waiting time to burst time



## Response Time Graph

 understandably, MLFQ takes the least response time, since the time quantum assigned to its first rr queue is smaller than that of the rr algorithm; followed by rr and then fcfs



## Results & Analysis

 round robin always takes least response time, and differs depending on the time quantum assigned to it

 FCFS always leads to the shortest waiting and turnaround time, this is due to a relatively short burst time (less convoy effect)

 FCFS is always assigned last in the MLFQ, because a process is guaranteed to complete execution in that algorithm; so we finish off the MLFQ with FCFS to conclude the execution of all remaining processes