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Exact and Approximate Algorithms for Combinatorial Optimisation, The Traveling Salesman Problem as An Application

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Chapter 1

Problem Statement

1.1 History

The first use of the term 'traveling salesman problem' in mathematical circles may have been in 1931-32, as we shall explain below. But in 1832, a book was printed in Germany entitled *Der Handlungsreisende, wie er sein soll und was er zu thun hat, um Aufträge zu erhalten und eines glücklichen Erfolgs in seinen Geschäften gewiss zu sein. Von einem alten Commis-Voyageur* ("The Traveling Salesman, how he should be and what he should do to get Commissions and to be Successful in his Business. By a veteran Traveling Salesman").

Although devoted for the most part to other issues, the book reaches the essence of the TSP in its last chapter: 'By a proper choice and scheduling of the tour, one can often gain so much time that we have to make some suggestions.... The most important aspect is to cover as many locations as possible without visiting a location twice ...' [Voigt, 1831; MiMer-Merbach, 1983].

Chapter 2

Branch and Bound

2.1 Motivation

The direct method as we have seen is, despite the simplicity of its implementation, unrealistically slow for even very small instances.

Faster exact algorithms exist, but none of them is polynomial since TSP is NP-complete. In fact, under the assumption $P \neq NP$, no polynomial solution exists.

Branch and Bound is one such algorithm that we will dedicate the rest of the chapter to.

2.2 The idea of Branch and Bound

The idea of Branch and Bound is to eliminate certain branches from the search space to decrease runtime.

This is done by computing a *lower bound* and *upper bound* for every branch, and then pruning branches that are guaranteed to be worse than the best known solution.

2.3 The implementation