Project (10): Diagonal Difference

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The first algorithm (non-recursive)

Pseudocode:

Read n from user.

Create arr with size n*n.

For i = 0 to n-1

For j=0 to n-1

Read arr[i][j] from user.

LeftSum=0

Rightsum=0

For i=0 to n-1

LeftSum = leftSum + arr[i][i]

Rightsum=RightSum+ arr[i][n-i-1]

Diff= absolute value of LeftSum-RightSum

Prints Diff

Time complexity t(n):

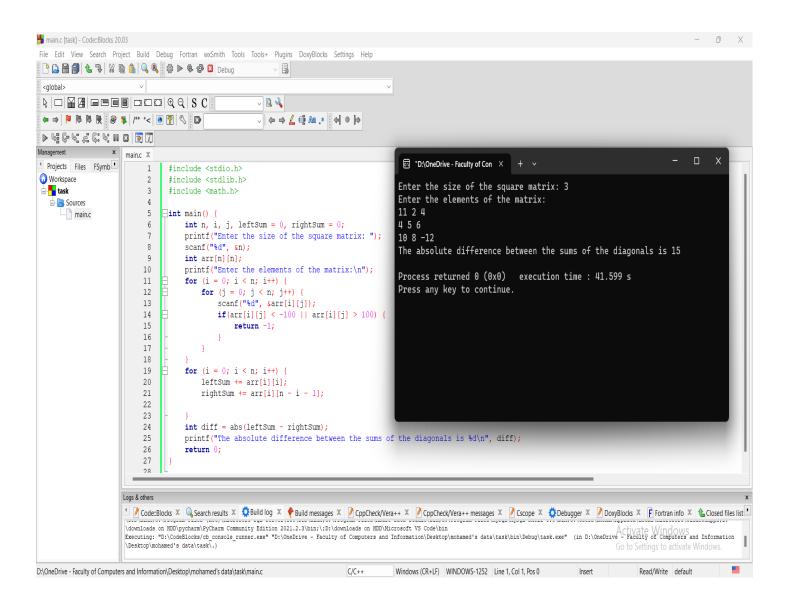
- Read the size of the square matrix from the user. This takes constant time and has a time complexity of O(1).
- Create a 2D array (arr) of size n x n. This takes O(n^2) time, since it needs to allocate memory for n x n elements.
- Read the elements of the matrix from the user and store them in the arr array. This takes O(n^2) time since it needs to iterate over each element of the array.
- checks if each element of the matrix is within the range [-100, 100].
 This check takes constant time per element, so the overall time complexity of this takes O(n^2).
- Initialize the variables leftSumDiagonal and rightSumDiagonal to 0.
 This takes constant time and has a time complexity of O(1).
- Iterate over the rows and columns of the matrix to calculate the sum of the left diagonal and the right diagonal. This takes O(n) time since it needs to iterate over each row and column of the matrix once.
- Calculate the absolute difference between the two sums using the absolute function. This takes constant time and has a time complexity of O(1).
- Print the result. This takes constant time and has a time complexity of O(1).

• Therefore, the overall time complexity of the code can be expressed as:

$$T(n) = O(1) + O(n^2) + O(n^2) + O(1) + O(n) + O(1) + O(1)$$

= $O(n^2)$

• The code



The second algorithm (recursive)

Pseudocode:

```
Define function diagonalSumDiff(matrix, n, i, leftsumdiagonal, rightsumdiagonal)
  if n = i
  returns abs(leftsumdiagonal - rightsumdiagonal)
  end if
  leftsumdiagonal += matrix[i][i]
  rightsumdiagonal += matrix[i][n - i - 1]
  i++
  return diagonalSumDiff(matrix, n, i, leftsumdiagonal, rightsumdiagonal)
end function
main function
user enters size of matrix n*n.
For i=0 to n
For j=0 to n
Read arr[i][j] from user.
If matrix[i][j] >=100 or matrix[i][j]<=-100
Prints Invalid matrix element.
Returns -1
intialize diff = diagonalSumDiff(matrix, n, 0, 0, 0);
prints diff
```

Time complexity t(n):

The time complexity of the diagonalSumDiff function in this algorithm is O(n), where n is the size of the matrix, because each recursive call of the function processes one element of each diagonal. The function makes a recursive call n times, once for each element in the main diagonal and the opposite diagonal, and performs a constant amount of work during each call.

The time complexity of the main function is O(n^2), where n is the size of the matrix, because it reads in n^2 elements from the input and checks their validity.

Therefore, the overall time complexity of the algorithm is $O(n^2)$.

The steps:

```
T(n) = T(n-2) + n-1 + n
T(n) = T(n-3) + n-2 + n-1 + n
T(n) = T(n-k) + kn - k(k-1)/2 ...(1)
```

For base case:

```
n - k = 1 so we can get T(1)

=> k = n - 1

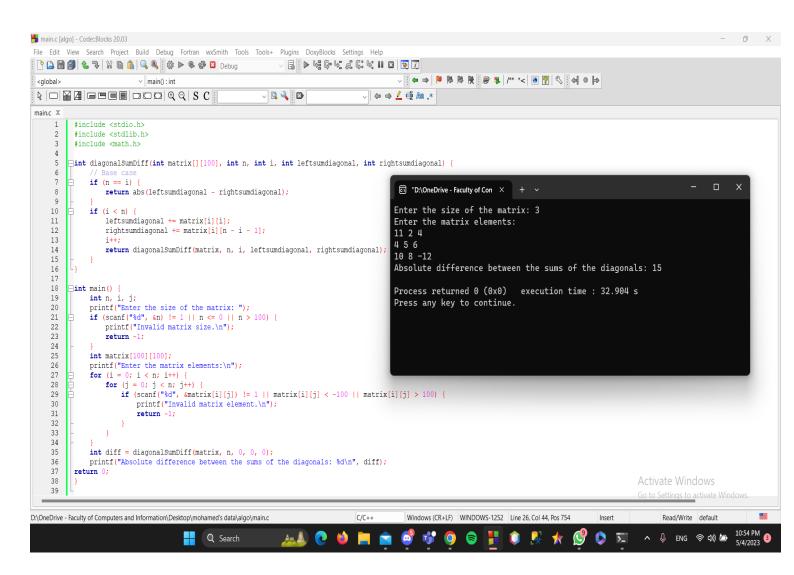
substitute in (1)

T(n) = T(1) + \frac{(n-1)n}{(n-1)(n-2)/2}

T(n) = T(1) + \frac{(n^2)}{(n-1)(n-2)}

So T(n) = o(n^2)
```

The code



The comparison between the two algorithms

comparison	The first algorithm(non-recursive)	The second algorithm(recursive)
The best case	O(1)	O(1)
	when the input matrix has a size of 1, which is also the smallest possible size for a square matrix.	when the input matrix has a size of 1, which is also the base case of the diagonalSumDiff function. In this case, the function will simply return 0 without performing any calculations.
The avg case	O(n^2)	O(n^2)
The worst	O(n^2)	O(n^2)
case		
	when the input matrix is a large	when the input matrix is a large
	square matrix of size n x n, where n	square matrix where the size of the
	is a very large number	matrix is a multiple of 2.
The most	both	
preferred		
algorithm		