Lect 11

Adder and Multiplier

CS221: Digital Design

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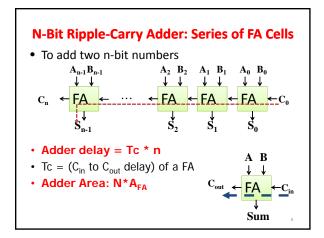
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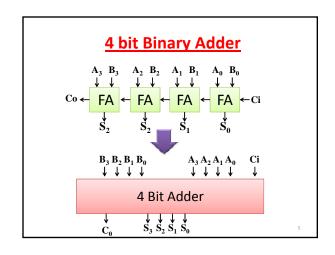
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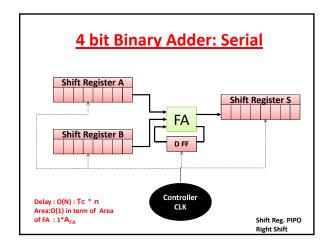
Outline

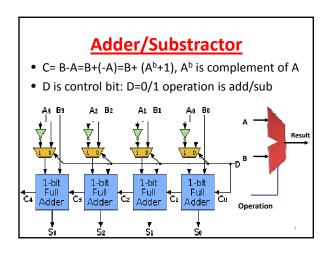
- Combinational Block
- Adder, Substractor, BCD Adder
- Efficient : Adder Design
 - RCA, CS_kA, CS_lA, CLA
- Binary Multiplier
 - Array, Sequential, Booth
- Floating Point

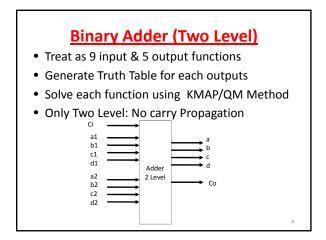
Adding Two One-bit Operands • One-bit Half Adder: A B Sum Cout 0 0 0 $Sum = A \oplus B$ 0 1 0 HA Cout = A.B1 0 0 • One-bit Full Adder: C_{in} A B Sum Cout 0 0 0 0 0 0 0 1 0 $Sum = A \oplus B \oplus Cin$ 0 1 0 0 Cout = A.B + B.Cin0 1 1 0 C_{out} FA + A.Cin 1 0 0 1 0 1 0 1 0 Şum 1 1 0 0 1 1 1 1





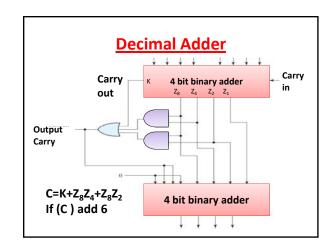


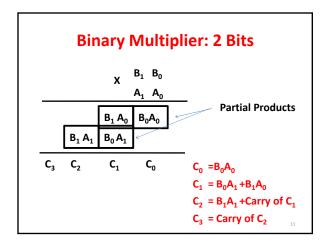


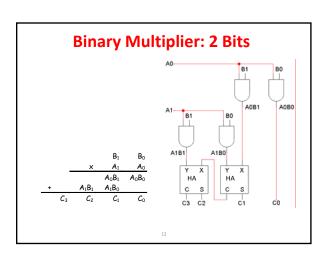


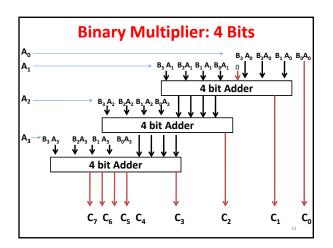
Decimal Adder

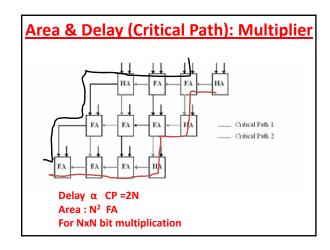
- Decimal numbers are represented with BCD code.
- When two BCD digits A and B are added
 - if A+B<10 result is a valid BCD digit
 - if A+B>9 result will not be valid BCD digit. It must be corrected by adding 6 to the result
- If A+B >9 add 6 to solve this issue

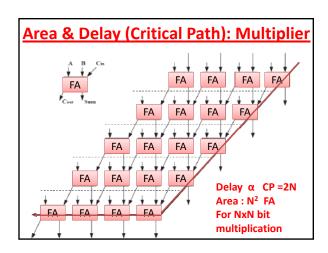














Adder Universal Use

- Adder : A = B + C
- Substractor: A = B + (-C), 2's complement
- Compare: C = A> B ? 1:0, (A-B > 0) ? 1:0
 Special case of compare with 0
- Multiply
- Divide
- Mod
- Floating point: Add/sub/mul...

Adding Two One-bit Operands • One-bit Full adder Cout Cin A B Sum 0 0 1 | 1 \dot{Sum} 0 1 $Sum = A \oplus B \oplus Cin$ 1 0 1 Cout = A.B + B.Cin0 0 1 + A.Cin 1 1 1

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Addition of Two N-Bit numbers

$$x + y + c_{in} = 2^n c_{out} + s$$

The solution:

$$s = (x + y + c_{in}) \mod 2^n$$

 $c_{out} = 1 \text{ if } (x + y + c_{in}) \ge 2^n \text{ else } 0$

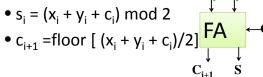
Example

- $011110 + 101101 = 1 (x 2^6) + 001011$
- X=30, Y=45
- $30 + 45 = 75 = 2^6 \times 1 + 11$
- Solution
 - $-S=(30+45+0)\%2^6=11$
 - Cout= 1 if (30+45+0 >= 2^6) else 0 = 1

Primitive module FA

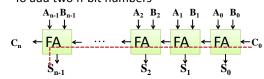
$$x_i + y_i + c_i = 2 c_{i+1} + s_i$$

with solution

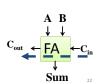


N-Bit Ripple-Carry Adder: Series of

• To add two n-bit numbers

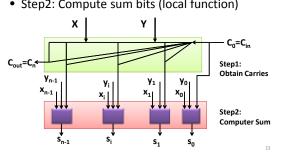


- Adder delay = Tc * n
- Tc = $(C_{in} \text{ to } C_{out} \text{ delay}) \text{ of a FA}$



Adder Schemes

- Step1: Obtain carries
 - (Carry at i depends on j < i), Non-trivial to do fast
- Step2: Compute sum bits (local function)



Mathematically: C_i & S_i

- $C_i = FuncC (x_{i-1}, ..., x_0, y_{i-1}, ..., y_0, c_{in})$
- $S_i = FuncS(x_i, y_i, c_i)$ $= (x_i + y_i + c_i) \mod 2$