

Tutorial → 03

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Sect → 'F'

Roll → 14

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Ans: 01

```
while (low ≤ high)
```

```
{
```

```
    mid = low + (high - low) / 2;
```

```
    if (arr[mid] == key)
```

```
    {
```

```
        print("Found");
```

```
        print(count);
```

```
    }
```

```
    else if (arr[mid] > key)
```

```
        high = mid - 1;
```

```
    else
```

```
        low = mid + 1;
```

```
        count++;
```

```
}
```

```
print("Not found");
```

Ans: 02

Iterative Insertion Sort:

```
for (int i=1; i<n; i++)
```

```
{ j = i-1;
```

```
  x = arr[i];
```

```
  while (j > -1 and arr[j] > x)
```

```
  { arr[j+1] = arr[j];
```

```
    j--;
```

```
  }
```

```
  arr[j+1] = x;
```

```
}
```

Recursive:

```
void InsertionSort (int arr[], int n)
```

```
{
```

```
  if (n ≤ 1)
```

```
    return;
```

```
  insertionSort(arr, n-1);
```

```
  int last = arr[n-1];
```

```
  j = n-2;
```

```
  while (j ≥ 0 and arr[j] > last)
```

```
  {
```

```
    arr[j+1] = arr[j];
```

```
    j--;
```

```
  }
```

```
  arr[j+1] = last;
```

```
}
```

Insertion sort is online sorting because whenever a new element come, insertion sort define its right place.

Ans: 03

Bubble Sort $\rightarrow O(n^2)$

Insertion Sort $\rightarrow O(n^2)$

Selection Sort $\rightarrow O(n^2)$

Merge Sort $\rightarrow O(n \log n)$

Quick Sort $\rightarrow O(n \log n)$

Count Sort $\rightarrow O(n)$

Bucket Sort $\rightarrow O(n)$.

Ans: 04

Online Sorting \rightarrow Insertion Sort.

Stable Sorting \rightarrow Merge Sort, Bubble Sort.

Inplace Sorting \rightarrow bubble Sort, selection Sort.

Ans: 05

Iterative Binary Search.

```
while (low <= high)
```

```
{ mid = low + (high + low) / 2;
```

```
  if (arr[mid] == key)
```

```
    return true;
```

```
  else if (arr[mid] > key)
```

```
    high = mid - 1;
```

```
  else
```

```
    low = mid + 1; }  $O(n \log n)$ 
```

```
return false;
```

Recursive Binary Search.

```
if if (low <= high)
```

```
{
```

```
    int mid = low + (high + low) / 2;
```

```
    if (arr[mid] == key)
```

```
        return true;
```

```
    else if (arr[mid] > key)
```

```
        BinarySearch(arr, low, mid - 1);
```

```
    else
```

```
        BinarySearch(arr, mid + 1, high);
```

```
}
```

```
return false;
```

Ans $\Rightarrow 6$ $T(n) = \underbrace{T(n/2)} + \underbrace{T(n/2)} + C$

$T(n) = T(n/2) + 1$

Ans $\Rightarrow 7$ void (vector<int> arr, int key)

```
{
```

```
    sort(arr.begin(), arr.end());
```

```
    int low = 0, high = arr.size() - 1;
```

```
    while (low < high)
```

```
    {
```

```
        if (arr[low] + arr[high] == key) {
```

```
            print(low, high)
```

```
            break;
```

```
        }
```

```
        else if (arr[low] + arr[high] > key)
```

```
            high--;
```

```
        else
```

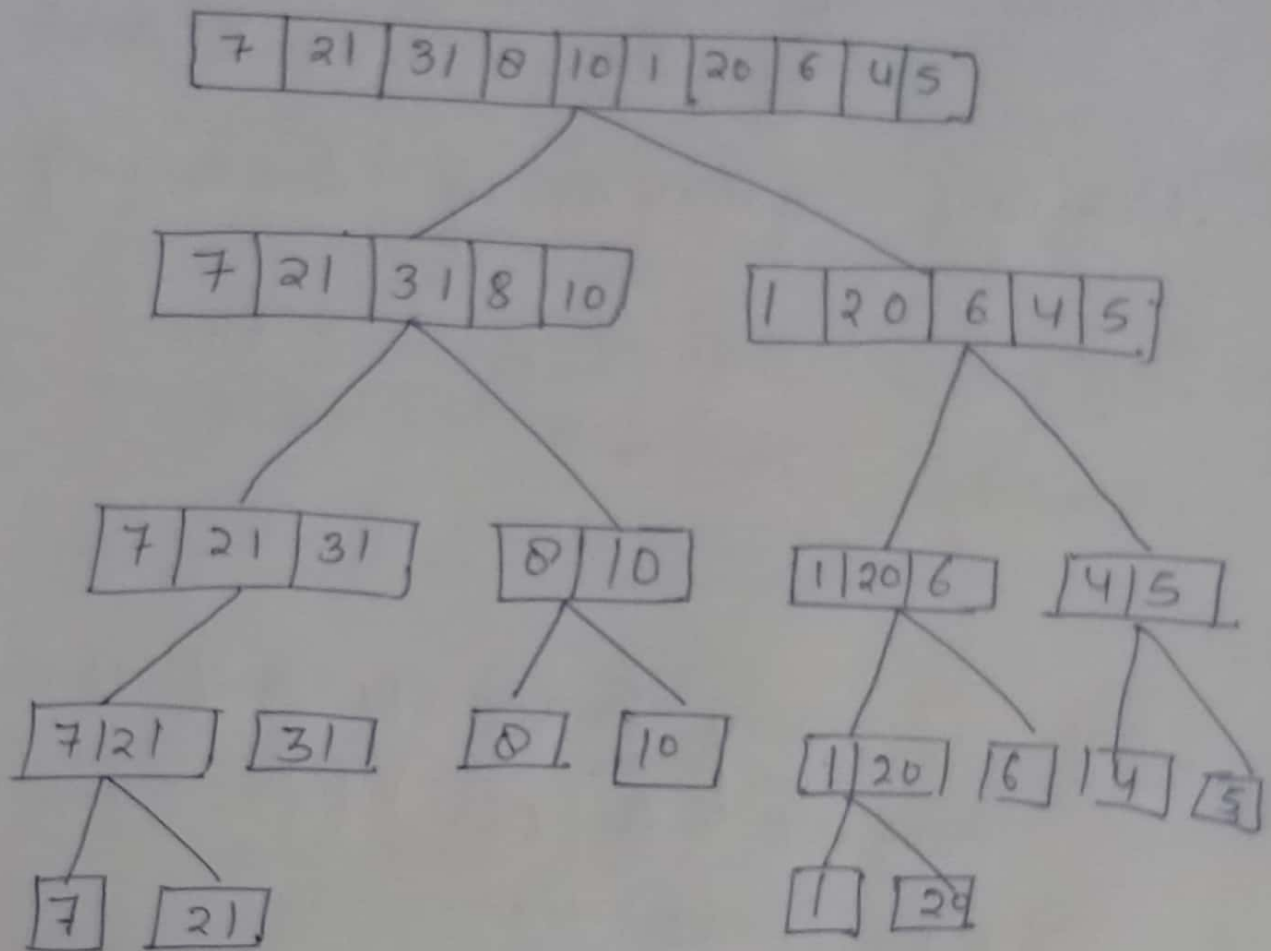
```
            low++;
```

```
    } print("No found");
```

```
}
```

Ans: 08 Quick Sort is the fastest general purpose Sort. In most practical solution, quicksort is the method of choice. If stability is important and space is available, merge sort might be best.

Ans: 09 Inversion indicates - how far are close the array is from being sorted.



Inversion = 31

Ques 9/11

Merge sort

Best Case $\rightarrow 2T(n/2) + O(n)$

Worst Case $\rightarrow 2T(n/2) + n$

Quick sort.

Best Case $\rightarrow 2T(n/2) + n$

Worst Case $\rightarrow T(n-1) + n.$

Basis	Quick sort	Mergesort.
• Position	Splitting done in any ratio.	Array divided in two equal parts.
• Works well on	Smaller array.	fine on any size of array.
• Sorting Method	Internal	External.
• Stability	Not Stable	Stable.

Ques 10 If array is already sort than Quick sort will take $O(n^2)$ time. otherwise it takes $(n \log n)$ time.

Ans 12:- Void StableSelectionSort (int arr[], int n)

```

{
    for (int i = 0; i < n - 1; i++)
    {
        int pos = i;
        int min = i;
        for (int j = i + 1; j < n; j++)
        {
            if (arr[j] < arr[min])
            {
                min = j;
            }
        }
        if (pos != i)
        {
            while (pos > i)
            {
                arr[pos] = arr[pos - 1];
                pos--;
            }
            arr[i] = arr[min];
        }
    }
}

```

Ques → 13 We will use the merge sort because we can divide the 4 GB data into 4 parts of 1 GB and sort them separately and combine them later.

Internal sorting → All the data is sorted in memory at all times sorting is in progress.

External sorting → all the data is stored outside memory and only load in small parts.