

TCP for Large Congestion Windows

High Speed TCP(HSTCP)

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TICN Presentation

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- ▶ Problem Statement
- ▶ What is High Speed TCP
- ▶ How is TCP different from HSTCP
- ▶ How is $a(cwnd)$ and $b(cwnd)$ is calculated
- ▶ Algorithm
- ▶ Comparison between standard TCP and HSTCP
- ▶ Problems with HSTCP
- ▶ References

Problem Statement

- ▶ TCP's average congestion window is roughly $1.2/\sqrt{p}$.
where

$$p = \frac{\text{number of packets marked or number of packets dropped}}{\text{total number of packets sent}}$$

p is known as packet drop rate

- ▶ Maintaining an average cwnd of at least 1.2×10^k packets requires a packet loss/corruption rate of at most 10^{2k}
- ▶ A Standard TCP connection with:
 - ▶ 1500-byte packets
 - ▶ a 100 ms round-trip time
 - ▶ a steady-state throughput of 10 Gbpswould require:
 - ▶ an average congestion window of 83,333 segments
 - ▶ and at most one drop (or mark) every 5,000,000,000 packets (or equivalently, at most one drop every 1 2/3 hours).
- ▶ **this is not possible!!**

What is High Speed TCP

- ▶ Just like Standard TCP when cwnd is low
- ▶ More aggressive than Standard TCP when cwnd is high.
- ▶ important question is when is cwnd is low and at what point it is going to differ with standard TCP.
- ▶ if the size of cwnd is 38 segments then it is considered as low cwnd and standard TCP is followed.
- ▶ but if size of cwnd is more than 38 segments then HSTCP is followed.

How is TCP is different from HSTCP

- ▶ A standard TCP flow in congestion avoidance increases its window by at most **one packet** per round-trip time and
- ▶ a multiplicative decrease of **halving** the current congestion window
but in case of HSTCP :
- ▶ the congestion window increase by $a(cwnd)$ segments per round-trip time in the absence of congestion.

$$cwnd = cwnd + a(cwnd)$$

where $a(w)$ is 1 for standard TCP

- ▶ In case of congestion the congestion window is reduced by $b(cwnd)$ as

$$cwnd = (1 - b(cwnd)) * cwnd$$

where $b(w)$ is 0.5 for standard TCP

How is a(cwnd) and b(cwnd) is calculated: Algorithm

- ▶ important variables used in the algorithm are:
 - ▶ cwnd = current congestion window size.
 - ▶ Low_Window = 38 segments, used to differ from standard TCP.
 - ▶ High_Window = 83000 segments.
 - ▶ High_P = 10^{-7}
 - ▶ High_Decrease = 0.1 (ie decrease by 10%)

Algorithm: Sudo Code

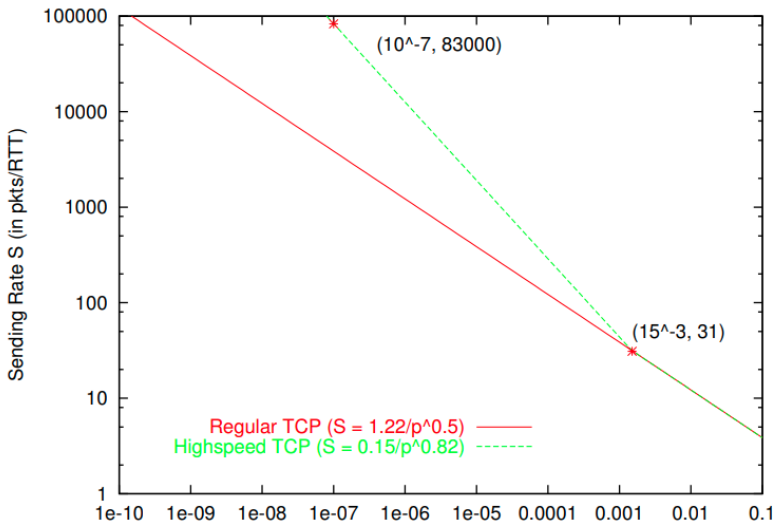
- ▶ Remember 2 imp formula :
 - ▶ for increasing cwnd, $cwnd = cwnd + a(cwnd)$
 - ▶ during congestion, $cwnd = (1 - b(cwnd)) \times cwnd$
- ▶ if $cwnd \leq Low_Window$
use standard TCP , $a(cwnd)=1$ and $b(cwnd)=0.5$
- ▶ if $(cwnd > Low_Window) \wedge (cwnd \leq High_Window)$
 $b(cwnd) =$
 $(High_Decrease - 0.5) \frac{(\log(cwnd) - \log(Low_Window))}{(\log(High_Window) - \log(Low_Window))} + 0.5$
 $a(cwnd) = w^2 \times p(cwnd) \times 2 \times b(cwnd) / (2 - b(cwnd))$
 $p(cwnd) = \frac{0.078}{w^{1.2}}$

Algorithm: Contd..

- ▶ if $cwnd > High_Window$
 $b(cwnd) = High_Decrease$
 $a(cwnd) = High_Window^2 \times High_P \times 2 \times b(w) / (2 - b(w))$

Comparison between standard TCP and HSTCP

HighSpeed TCP: the modified response function.



► Unfair

- it is unfair to the devices which will be using standard TCP as in HSTCP the rate at which we are increasing the window size is more and
- also during the congestion standard TCP reduce congestion window by half while in HSTCP we reduce by small amount as comparison to standard TCP.
- it is also unfair to new HSTCP connection if already a older HSTCP data packets are present on the same bandwidth.

- ▶ RFC 3649 : <https://www.ietf.org/rfc/rfc3649.txt>
- ▶ Home page HSTCP
: <http://www.icir.org/floyd/hstcp.html>
- ▶ Wiki page: <https://en.wikipedia.org/wiki/HSTCP>