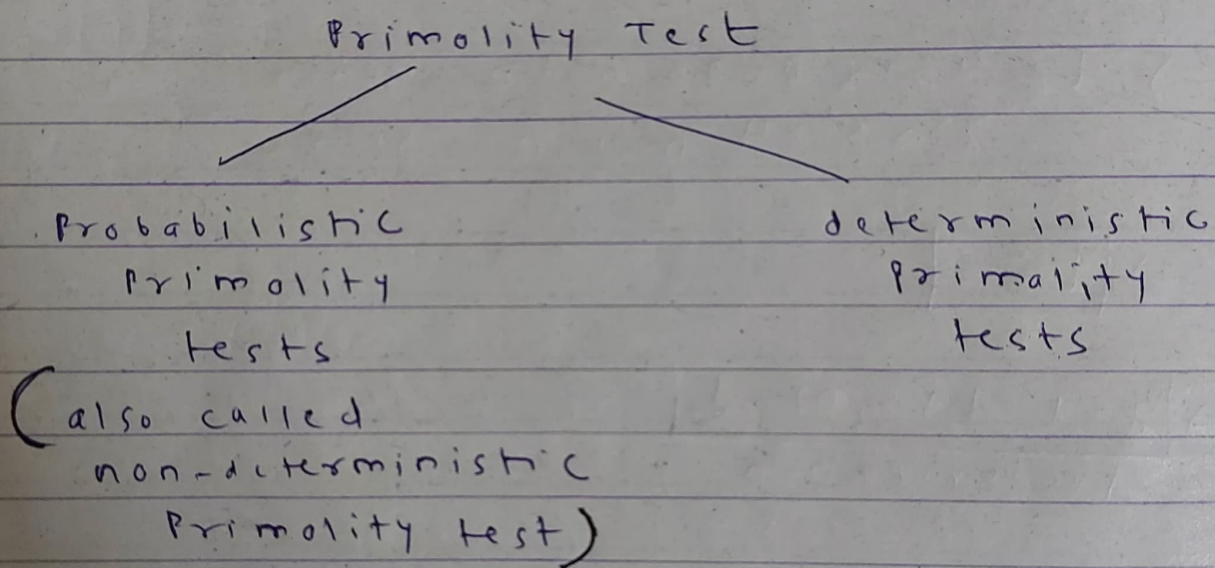


## BTP Thesis Part-1

### Primality Test :

A primality test is a function that determines if a given integer greater than 1 is prime or composite.



### Deterministic Primality Tests:

If we put an integer  $n$  in the primality test, then the output will be "yes" if the integer is prime, "no" if integer is not prime (composite integer).

## Non-deterministic Primality Test:

This test takes an integer  $n$  and returns "no" if  $n$  is not prime, or returns "may be a prime".

A non-deterministic primality test also called probabilistic primality test returns either that the inputted integer is not a prime or that is "probably a prime" to some given degree of likelihood.

## Primality Test Algorithms:

### ① Sieve of Eratosthenes Primality Test:

This is the simplest Algo. to test whether a given integer is prime or not. It also finds the prime numbers less than or equal to given integer  $n$ .

~~Algorithm~~

① Input:  $n \in \mathbb{N}$

②  $a[1 \dots n]$  = integer array

③



→ To check prime

Algorithm: (Input  $n \in \mathbb{N}$ )  
↳ natural numbers

①  $a[1 \dots n] = \text{integer array}$

② for  $j = 1$  to  $n$  do  
{

③  $a[j] \leftarrow j$ }

④  $i \leftarrow 2$

⑤ while  $i^2 \leq n$  do  
    if  $a[i] \neq 0$  then  
         $t \leftarrow 2 \cdot i$   
        while ( $t \leq n$ ) do  
             $a[t] \leftarrow 0$   
             $t \leftarrow t + i$

~~complete~~  
 $i = i + 1$

⑥ for  $j = 2$  to  $n$  do  
    if  $a[j] \neq 0$  then return  $a[j] = \text{prime}$

Ex:  $n = 10$

primes  $\leq 10$  (Sieve of Eratosthenes)

Algo:

~~1~~   2   3   ~~4~~   5   ~~6~~   7   ~~8~~   9   ~~10~~  
          ↑           ↑           ↑           ↑           ↑

~~1~~   2   3   ~~4~~   5   ~~6~~   7   ~~8~~   ~~9~~   ~~10~~  
          ↑           ↑           ↑

~~1~~   2   3   ~~4~~   5   ~~6~~   7   ~~8~~   ~~9~~   ~~10~~

⇓

2, 3, 5, 7 are primes

⇒ 10 not a prime

Complexity Analysis:

Sieve method has  $O(10^{\log_{10} n})$  operations to prove primality is  $n$  has  $\log_{10} n$  digits.

We can say this algorithm is of exponential time in terms of input length.



## ② Trial division Algorithm:

### Algorithm:

For  $k = 2, 3, 4 \dots \sqrt{n}$

Test if (for any  $k$ )

$n \equiv 0 \pmod{k}$  output composite

else

return  $n = \text{prime}$

### Complexity Analysis:

Trial division runs in time  $O(\sqrt{n} \log^2(n))$ , but this running time is exponential in the input size since the input represents  $n$  as binary number with  $\lceil \log_2(n) \rceil$  digits.

## ③ Wilson's characterization of primes:

### Wilson's theorem:

A natural number  $n > 1$  is a prime number if and only if the product of all positive integers less than  $n$  is one less than a multiple of  $n$ .

that is

~~That is~~ (using the notations of modular Arithmetic), the factorial  $(n-1)! = 1 \times 2 \times 3 \times \dots \times (n-1)$  satisfies:

$$(n-1)! \equiv -1 \pmod{n}$$

$$\text{OR } \frac{(n-1)! + 1}{n} = 0$$

implies  $n$  is prime

Ex:

For  $n = 2$

~~$$(n-1)! = 2$$~~

~~correct calculation~~

~~$$-1 \pmod{2} = 2 - 1 = 1$$~~

$$(n-1)! = (2-1)! = 1 \quad \text{LHS}$$

$$-1 \pmod{2} = 2 - 1 = 1 \quad \text{RHS}$$

LHS = RHS, Hence  $n = 2$  prime

$n = 13$

$$(n-1)! = (13-1)! = 12!$$

$$= 479001600$$

$$(n-1)! + 1 = 479001600 + 1 = 479001601$$

$$\text{since } 479001601 \pmod{13} = 0$$

$\Rightarrow 13$  prime



## Complexity Analysis:

From the Wilson's characterization of primes we can determine the primality of an integer  $n$  by calculating  $(n-1)! \pmod{n}$ .

But the computation requires  $(n-1)$  multiplications, making it very time consuming.

## ④ Euler test:

Euler test is based on simple lemma:

Lemma:  $n$  is prime iff  $\phi(n) = n-1$

Proof: If  $n$  is prime number, then every integer less than  $n$  is relatively prime to it, hence by definition  $\phi(n) = n-1$

Conversely, if  $n > 1 = \text{composite}$ , then  $n$  has a divisor  $d$  such that  $1 < d < n$ . It follows that there are at least 2 integers among  $1, 2, \dots, n$  which are not relatively prime to  $n$ , namely  $d$  &  $n$  itself. As a result,  $\phi(n) < n-1$ . This proves the lemma



## Algorithm (Euler Test):

Input:  $n \in \mathbb{N}$

check  $\phi(n)$ ;

if  $\phi(n) = n - 1$

then output "prime"

else

output "composite"

## Complexity Analysis:

From the Euler test, we can determine the primality of integer  $n$  by calculating  $\phi(n)$ .

But for calculating  $\phi(n)$ , we require the factors of  $n$  and factorization is more difficult problem than primality test.

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