## MOHET SHARMA (2019 CS 10372)

- A1.3.1: n operations performed on DLL.
  - 1) Worst lase time complexity of <u>Insert</u>: O(1)
    In insert, a new node of the given data is made which takes constant time and then 4 pointer values are initialised which is ob constant time.
    So, obserall time complexity is O(1).
  - In worst lase time complexity of the get First: O(n) In worst lase, DII will have n nodes after n operations. If get First is called from tail node, then the while look nums exactly n times. So, T(n) = n = O(n) other helper functions used (is Empty (), is Head () are of constant time complexity).
  - In delete, get First function is called which may takes n unite of operation in worst case.

    Then the while loop would run over all the nodes in DII (n), in worst case when the last node (before tail node) matches the node to be deleted, so loop runs n times.

    The deletion then takes \$8 hointer initialisation in worst case which is 0(1)

    get Next function show used is taking 0(1) time (shown below).

    So, in worst case, T(n) = n + n + o(1) = O(n).
  - 4) Worst lase time complexity of get Next: O(1)Helper functions used (is Empty, is tail) do constant number of matches so they have O(1) time complexity.

    In get Next, these constant time taking functions are used 3 times in worst case.

    So, T(n) = O(1)

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- 5) Worst lase time complexity of Find: O(n).

  The get First function is used, it takes on unit of time operation in worst lase.

  Then the while loop is run n times over all DII nodes in worst case when no match is found.

  So, T(n) = n + n + o(1) = O(n).
- 6) Worst lose time Nomplexity of Sanity: O(n).

  In sanity function, the cycletrists function takes O(n) time. Standard slow and fast pointer technique is used in which slow pointer moves a max of 2n times.

  Other than this, the DII is traversed in forward and lackward direction, taking a total of n+3 iterations. (head and tail nodes are included).

  So, T(n)= 2n + n+3 + O(1)= O(n).

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- A.1.3.2: n operations performed of Allocate / Free
  - In worst case free llk can have  $[n/2] \perp n/2 \rfloor + 1$  nodes. So number of nodes is O(n) in feweblk. DLL. In worst case, in Allocate function, a find 2 delete function is used on besellk and two inserts one mallocable k one on freeBlk. The salls are made from fead so get First in find k so,  $T(n) = k \lfloor n/2 \rfloor + 1 + \lfloor n/2 \rfloor + 1 + O(1)$ .

Find Delete Insertions.

T(n) = O(n).

2) Worst lase time tomplexity of Free: O(n).

In worst case, allocally can have n nodes.

So find and delete on allocally will take no operations.

As call is made from head node, so it into find and delete, get First used takes O(1) time.

So, t(n) = n + n + o(1)Find Delete Insertion t(n) = o(n).