

# Morris Traversal

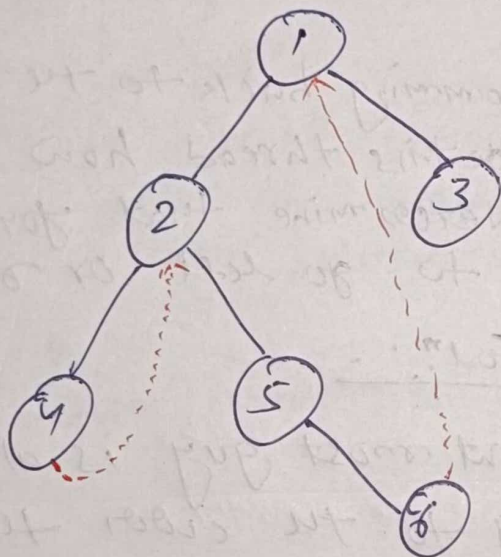
$$T.C = O(N)$$

$$S.C = O(1)$$

Morris Traversal uses the concept of threaded BT

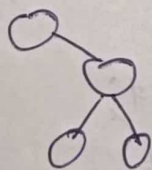
Inorder (left root right)

inorder - 4 2 5 6 1 3



pattern  $\rightarrow$  from the last node of any <sup>left</sup> subtree you go back to the root and after that you back to the right.

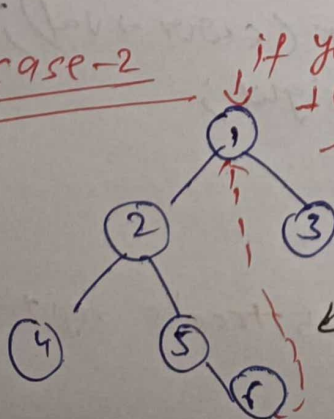
Case-1



if (left == NULL)

inorder.push\_back(root->val);  
root = root->right;

Case-2

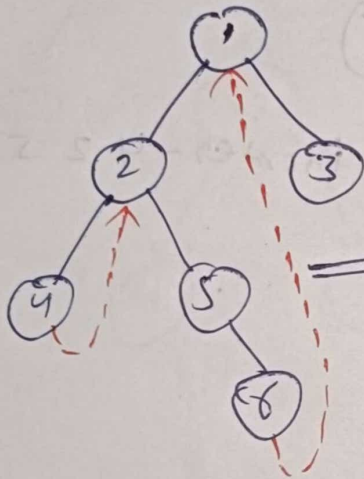


if you are standing at  
+ (cursor) before moving to  
left you should make the  
thread/pointer

BCZ if you once  
move to the 2  
you can't come back  
at 1 without the  
thread.

Before going to the left subtree create a thread by the right most guy of the left subtree to the root.

after that  $cwr = cwr \rightarrow left$ .



after coming back to the 1 via this thread how you can determine that you have to go left or right.

Soln:

if right most guy is already pointing to the cwr then go to the right of the cwr & remove the thread.

Code

```
vector<int> inorder (TreeNode* root)
{
    vector<int> inorder;
    TreeNode* cwr = root;
    while (cwr != NULL)
    {
        // case-1
        if (cwr->right == NULL)
        {
            inorder.push_back(cwr->val);
            cwr = cwr->right;
        }
        else
        {
            // go to the left subtree's right most guy for thread.

```



TreeNode\* prev = curr->left;

while (prev->right && prev->right != curr)

{  
    prev = prev->right;

if (prev->right == NULL)

{  
    prev->right = curr; ← create a thread to point to the root;

for preorder  
remove the  
line & put  
the arrow  
location.

curr = curr->left;

else

if thread already exist.  
prev->right = NULL;

inorder push-back (curr->val);

curr = curr->right;

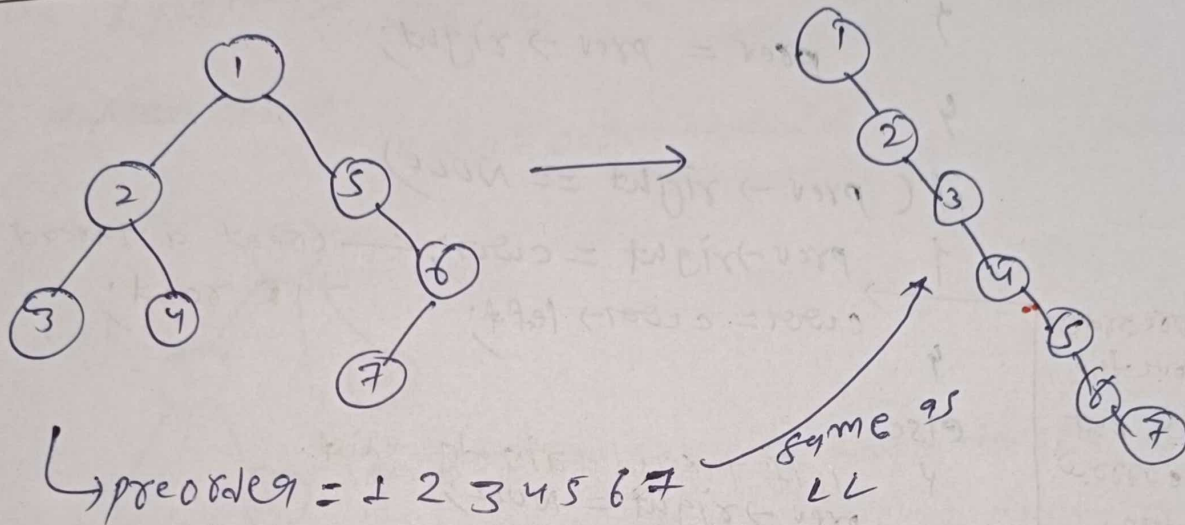
return inorder;

T.C. =  $O(N)$

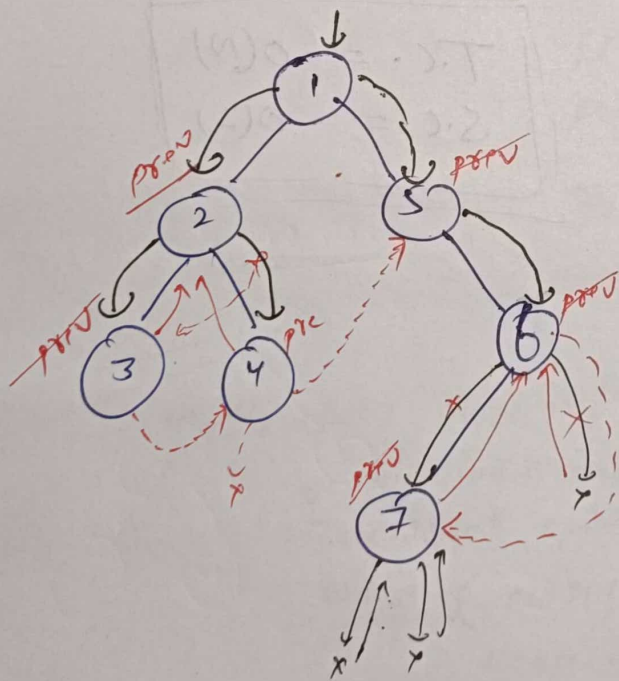
S.C. =  $O(1)$

The inner while loop will  
run completely  $O(N)$  at  
most.

# Flatten a Binary Tree to a Linked List



~~pre = NULL~~ 7 6 5 4 3 2



Code

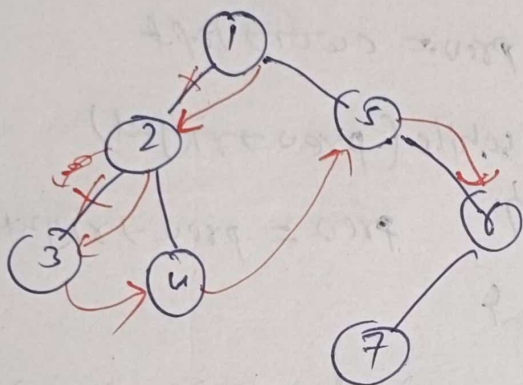
```
prev = NULL;  
flatten(node)  
{  
    if (node == NULL)  
        return;  
  
    flatten(node -> right);  
    flatten(node -> left);  
  
    node -> right = prev;  
    node -> left = NULL;  
    prev = node;  
}
```

T.C. =  $O(N)$   
S.C. =  $O(N)$



## Approach - 2

curr = ~~1~~ ~~2~~ ~~3~~ ~~4~~ ~~5~~ ~~6~~ ~~7~~



~~7~~  
~~6~~  
~~5~~  
~~4~~  
~~3~~  
~~2~~  
~~1~~

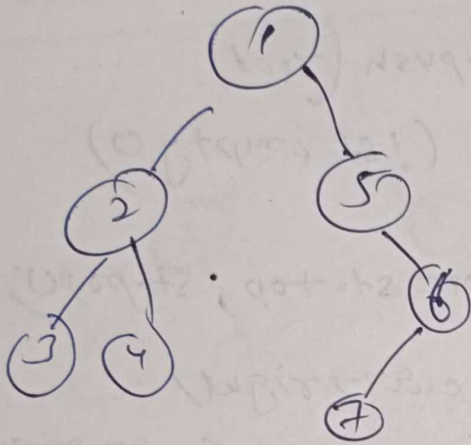
```
st.push(root)
while (!st.empty())
{
    curr = st.top(); st.pop();

    if (curr->right)
        st.push(curr->right);
    if (curr->left)
        st.push(curr->left);

    if (!st.empty())
        curr->right = st.top();
    curr->left = NULL;
}
```

T.C. =  $O(N)$   
S.C. =  $O(10)$

Now we use the morris traversal



$T.C. = O(N)$   
 $S.C. = O(1)$

```

curr = root,
while (curr != NULL)
{
  if (curr->left != NULL)
  {
    prev = curr->left
    while (prev->right)
    {
      prev = prev->right
    }
  }

```

```

  prev->right = curr->right;
  curr->right = curr->left;

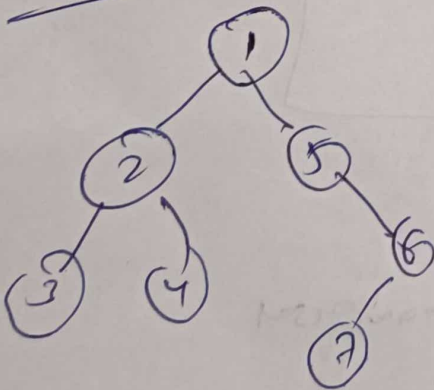
```

```

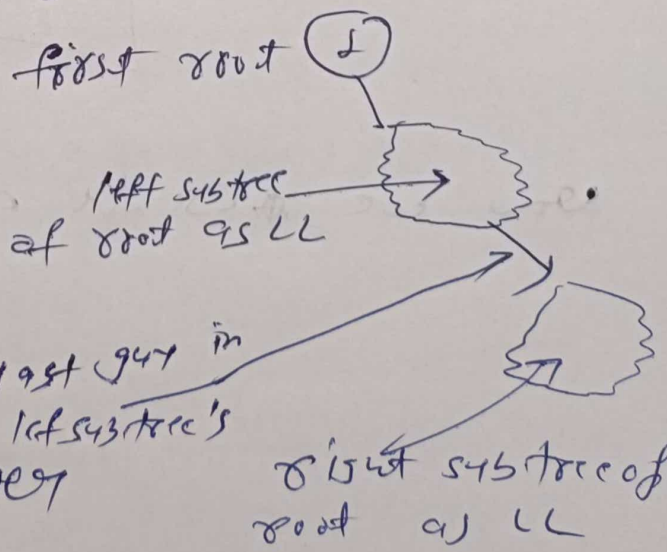
  curr = curr->right;
}

```

Thought process



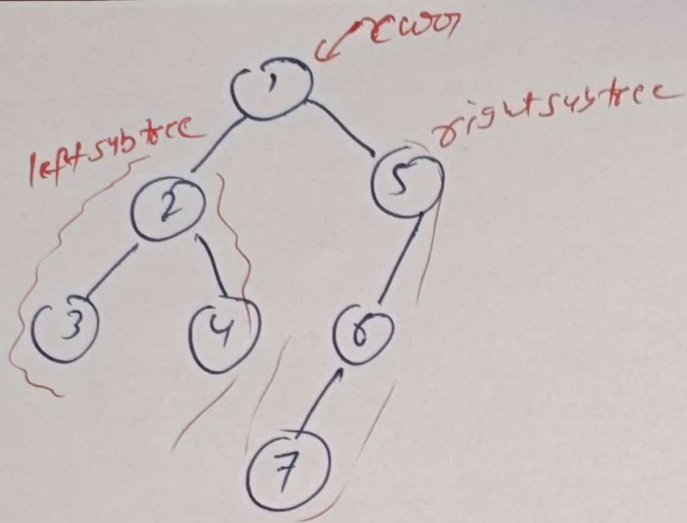
BT  $\xrightarrow{\text{convert to LL}}$  LL  
 ↳ preorder same as



the last node in  
 preorder of left subtree's  
 right pointer

right subtree of  
 root as LL





left subtree में record का last guy find करो और उसे right की curr के right में connect कर दो & curr को right में curr का left कर दो

if (curr->left != NULL)

↑ prev = curr->left

while (prev->right)

prev = prev->right

prev->right = curr->right

curr->right = curr->left

{