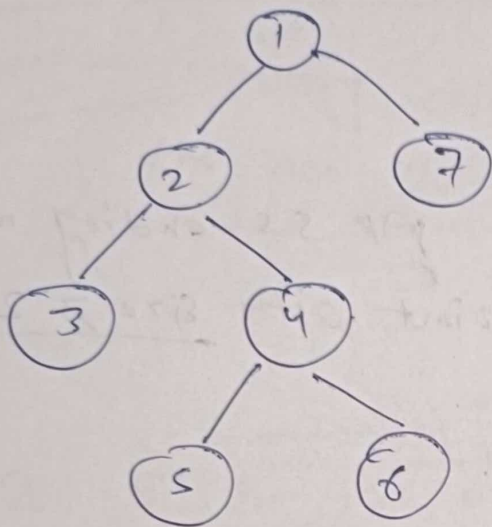


# Iterative Preorder Traversal in Binary Tree



Preorder  $\rightarrow 1\ 2\ 3\ 4\ 5\ 6\ 7$

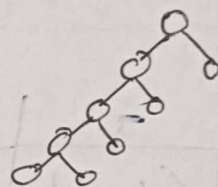
$$T.C. = O(n)$$

$$S.C. = O(n) \approx O(h)$$

$\downarrow$  when tree like one right & more left

data structure use  $\rightarrow$  stack

initially put root onto stack

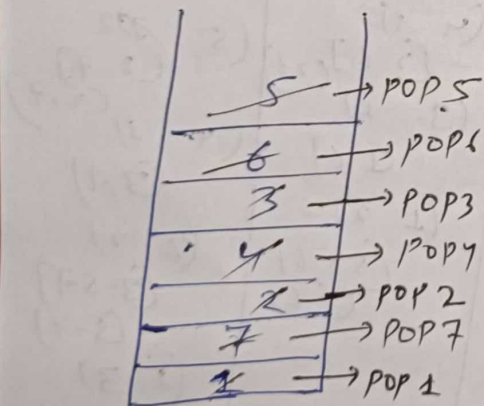


Now in loop while stack is not empty do following steps.

(i) pop the Node from stack top & print it

(ii) put right child onto stack

(iii) put left child onto stack.



1 2 3 4 5 6 7

preorder.

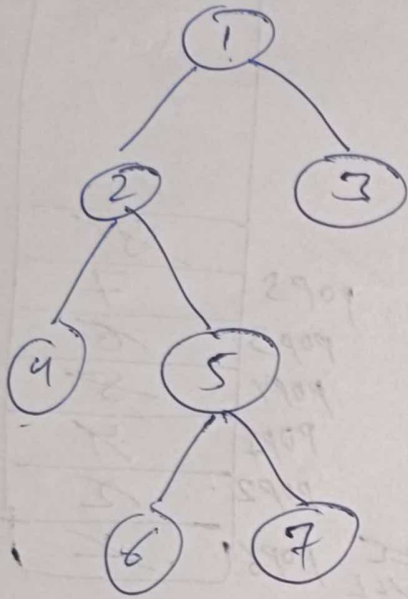
why we put first right & then left it is because in preorder

we go for root left right that means after the root left part traversed

stack is LIFO Data structure & if we put

left in last that means it stay on top of the stack.

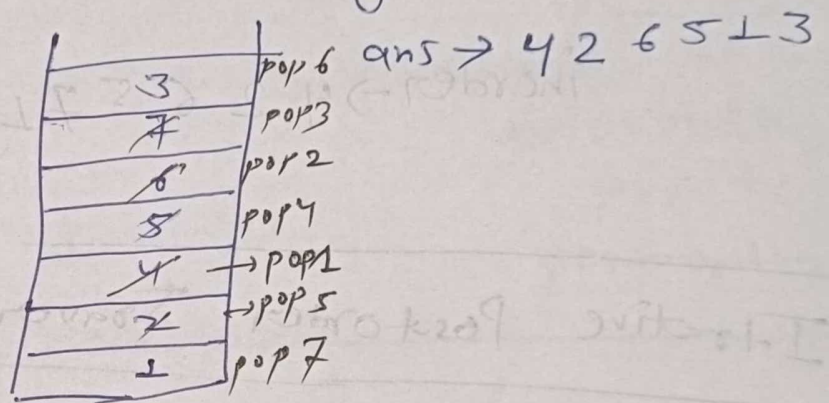
# Iterative Inorder Traversal for Binary Tree



Inorder  $\rightarrow$  left Root right

$\hookrightarrow$  4 2 6 5 7 1 3

In recursion how stack is working

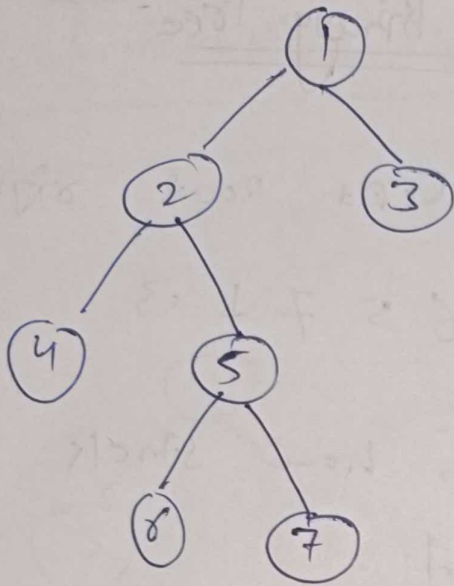


In iterative ~~and~~ inorder we maintain a stack data structure similar to recursion that means

we take a node = root variable ~~and while no.~~ and in every iteration we follow these steps.

```
if (node != NULL)
{
    st.push(node);
    node = node->left;
}
else
{
    if (st.empty()) break;
    node = st.top();
    st.pop();
    print (node);
    node = node->right;
}
```





node = ~~1~~  
2  
~~4~~

NULL  
NULL  
~~8~~

~~6~~  
NULL  
NULL

~~7~~  
NULL  
NULL

inorder → 4 2 6 5 7 3

pop5

pop3

pop4

pop1

pop2

pop6

<del>8</del>
7
<del>6</del>
<del>5</del>
<del>4</del>
2
<del>1</del>

## Iterative Postorder Traversal

### Approach 1 - using 2 stack

take 2 stack st1 & st2

st1 initially store the root.

apply while loop & do following steps.

while ( !st1.empty() )

{

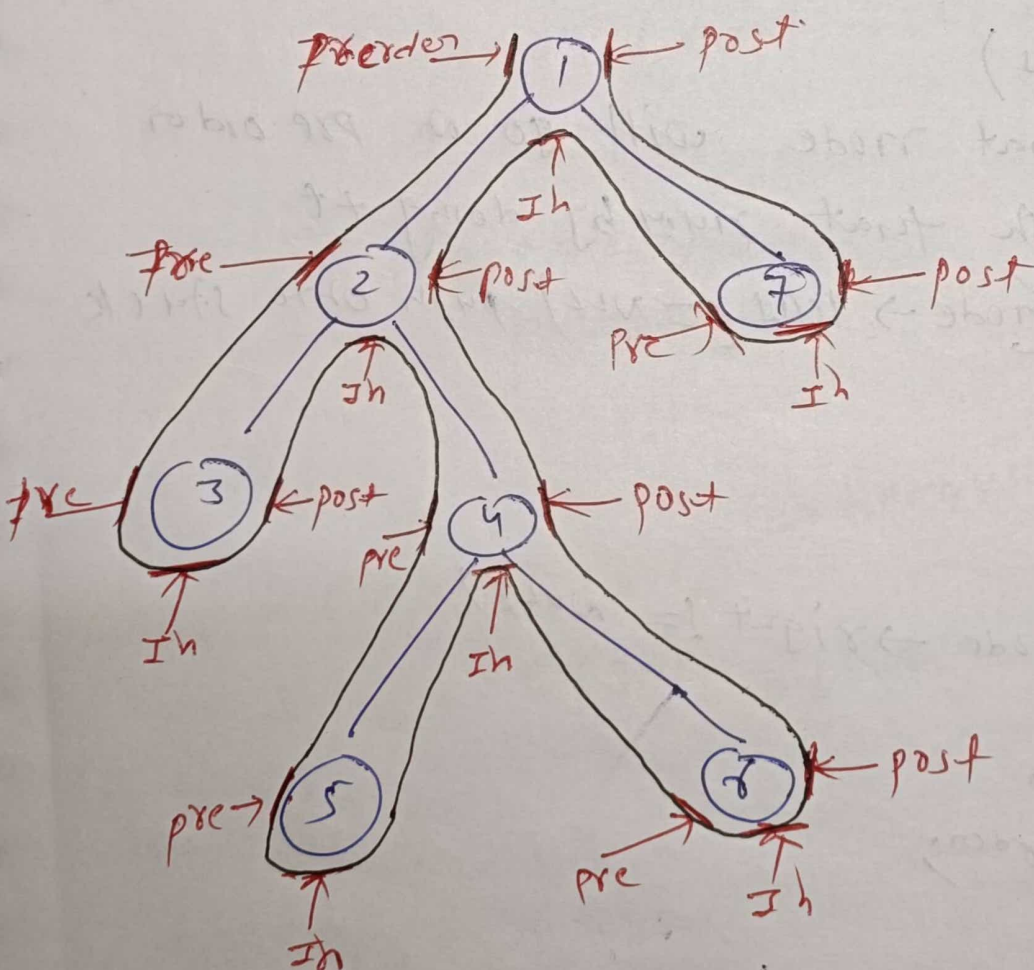
step 1 → pop the top element from st1.

and push this element into st2,

step 2 → push the left child of the pop node if not NULL

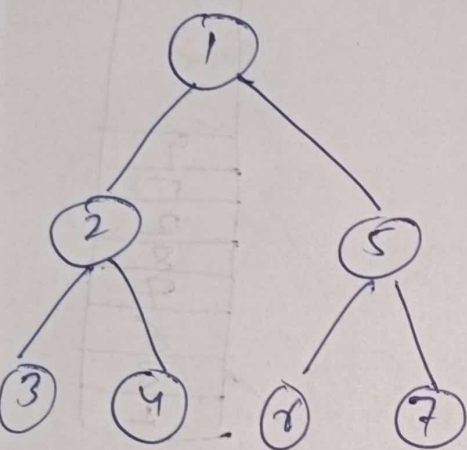
step 3 → push the right child of the pop node if not NULL

}





# Inorder / Preorder / Postorder



Initially take a stack  
 & insert (root, val) that  
 means stack store a  
 pair  $\langle \text{Node}^*, \text{int} \rangle$   
 ↑                      ↑  
 pointing              this  
 to                      will tell  
 node of              that inorder  
 tree                      pre order  
                              &  
                              post order

Rule:-

~~while you~~  
 if whenever you taking out an  
 element from the stack & the num of the element  
 is like

if (num == 1)

{ then that node will go on pre order  
 and push that num by doing ++

if ~~the~~ (node → left != NULL) push onto stack

↓

if num == 2

inorder

num ++

push (node → right != NULL)

if

num == 3

post order

Rule for every iteration

```
pair <Node*, int> = st.top()
st.pop();
```

```
int num = pair.second;
```

```
if (num == 1)
```

```
{
```

```
    pre order
```

```
    push (pair.first, num+1)
```

```
    push (pair.first->left, -1)
```

```
}
```

```
if (num == 2)
```

```
{ in order
```

```
    push (pair.first, num)
```

```
    push (pair.first->right, -1)
```

```
}
```

```
if (num == 3)
```

```
    post order.
```

T.C. =  $O(3N)$

S.C. =  $O(3N)$  or  $O(4N)$

because

when you draw  
stack then  
you are  
visiting  
every node  
3 times.