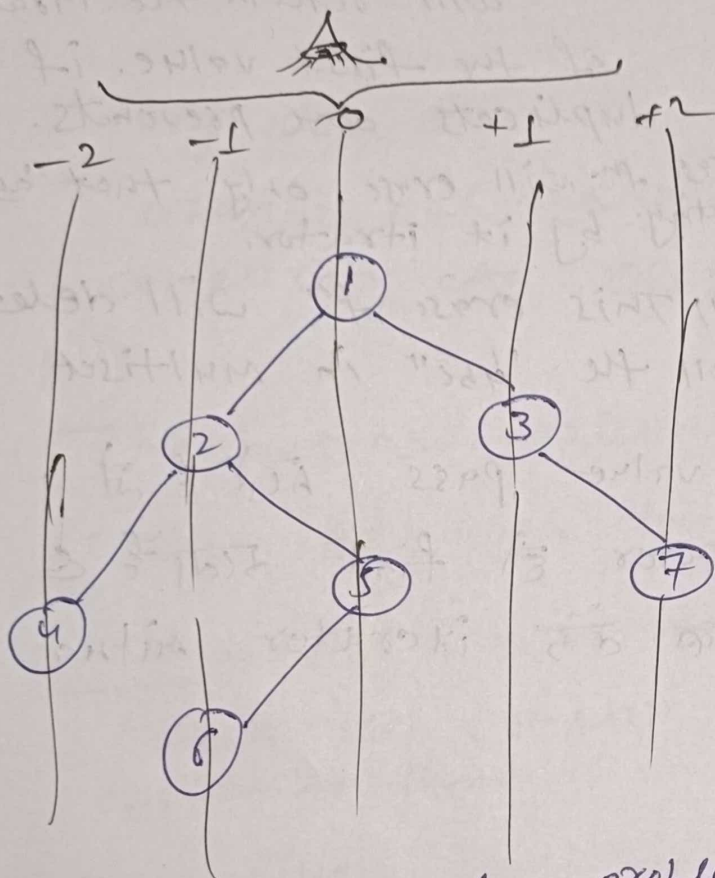


Top View of Binary Tree



Top View $\rightarrow 4\ 2\ 1\ 3\ 7$

$$T.C. = O(N)$$

$$S.C. = O(N)$$

In order to solve problem on BT we have to traverse the tree for which we have traversal

BFS
(level order)

DFS
(inorder, preorder, postorder)

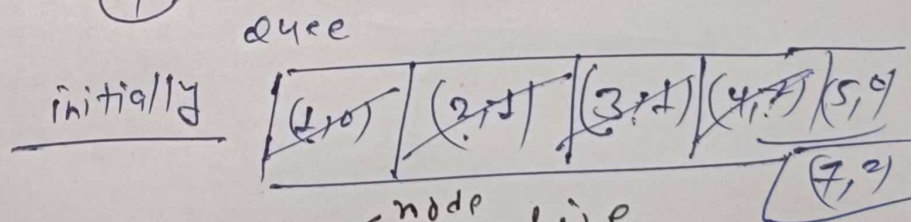
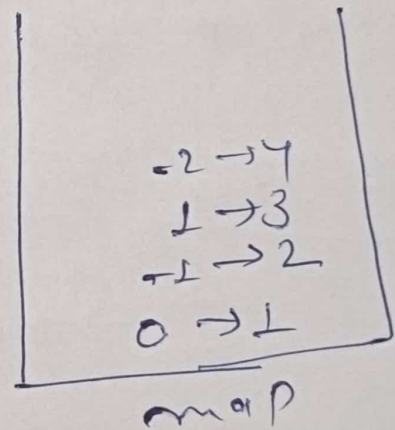
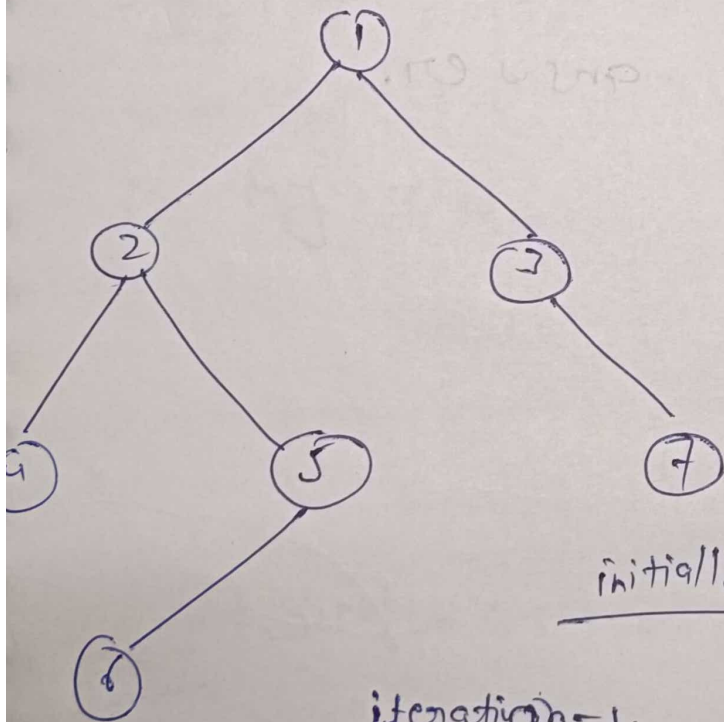
we use level order traversal here.

why not recursive (DFS) traversal this will be discuss at end.

Intuition :- In vertical order traversal we use line concept where we assign every node to a coordinate here we are using similar kind of concept where which we use only line concept like that the first node in every line will be my top view.

Data structure

- (i) queue $\langle \text{Node}^*, \text{line} \rangle$
- (ii) map $\langle \text{line}, \text{Node}^* \text{val} \rangle$



iteration-1. node = (1,0) ← line

if (map.find(0) == map.end())

map.insert(0, 1)

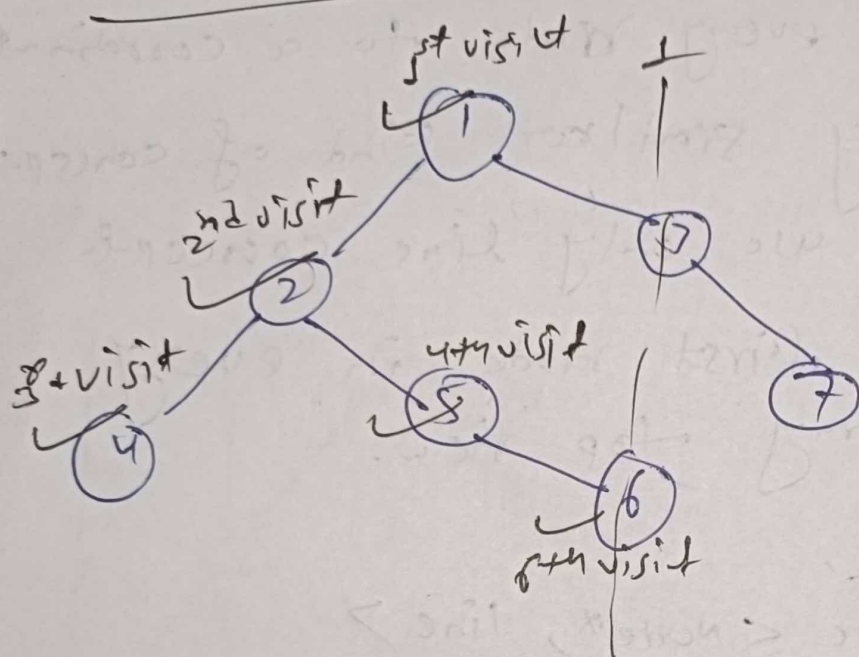
insert left & right child in que

2nd → node = (2,1)

map.find(-1) == map.end() →

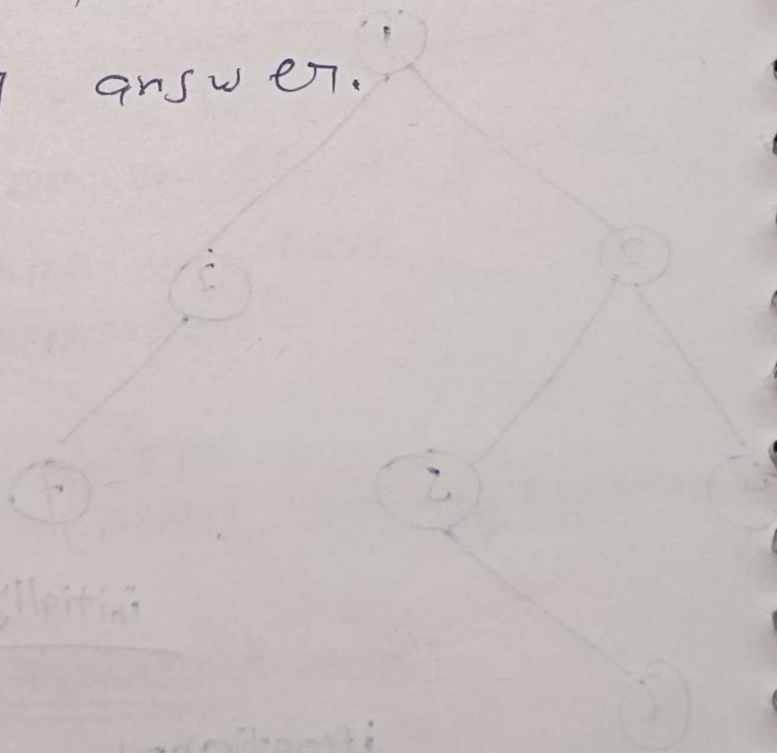
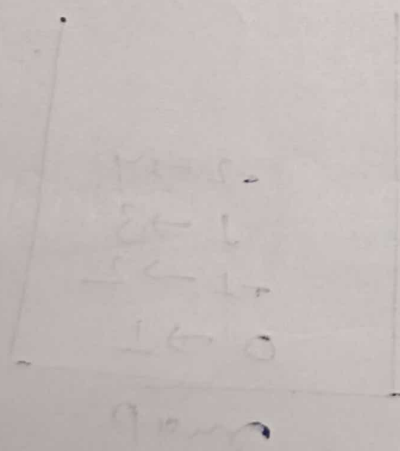
3rd → node = (5,0) but map.find(0) != map.end()

Why we not use recursive traversal.



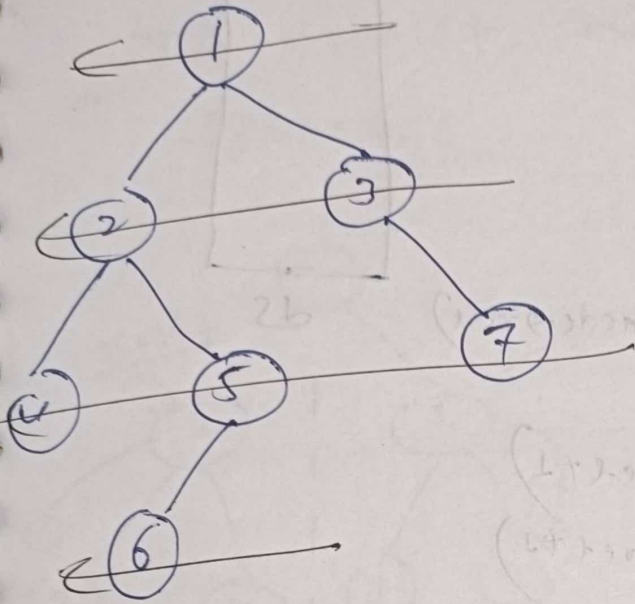
Suppose you in order traversal.

for line 1 we visit 6 before
 5 that's why will give
 wrong answer.



Right/left View of Binary Tree

Right view = 1 3 7 6



The last node of every level is indeed in my right side view

If I reverse the logic then the first node of every level is ~~may~~ indeed in my right side view.

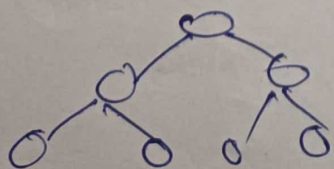
what kind of traversal we can use here

Recursive		Iterative (level order)
$T.C. = O(N)$ $S.C. = O(H)$		
$O(N)$ case worst \leftarrow height		

why we are not going with level order traversal.

b2. In the worst case

$T.C. = O(N)$
 $S.C. = O(N)$



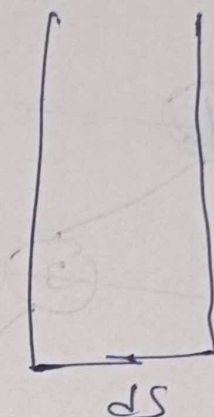
last level is filled so queue contain at max $n/2$ nodes.

$f(\text{node}, \text{level})$

2

if (node == NULL)
return;

if (level == ds.size)
ds.push_back(node->val)



$f(\text{node} \rightarrow \text{right}, \text{level} + 1)$
 $f(\text{node} \rightarrow \text{left}, \text{level} + 1)$

4

→ Root left right

Reverse Preorder

→ Root right left

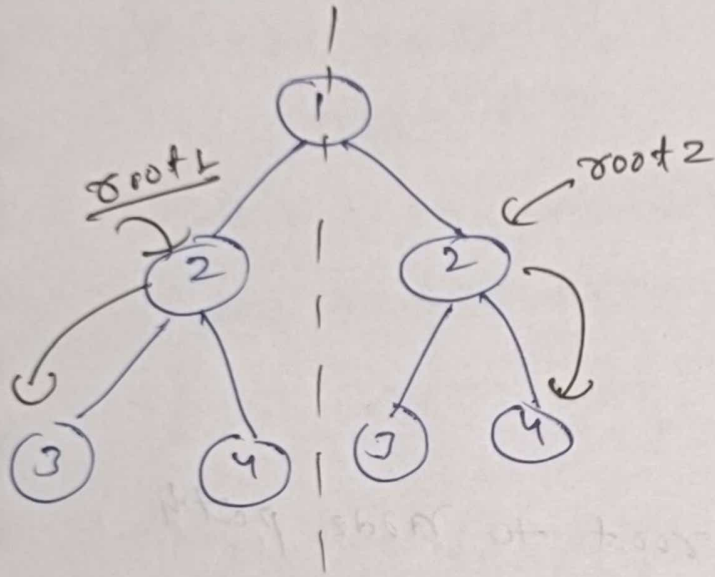
left side view

~~change the~~ apply only pre order

$f(\text{node} \rightarrow \text{left}, \text{---})$
 $f(\text{node} \rightarrow \text{right}, \text{---})$

Symmetrical Binary Tree

If tree is mirror of itself then symmetric,



if (root 1 == root 2)
return false

~~if~~
int x = f(root->left,
root->right);
int y = f(root->right,
root->left);

return (x & y);