1043. Partition Array for maximum Sum

arr = [1, 15, 7, 9, 2, 5, 10] , K=3

ans = 84

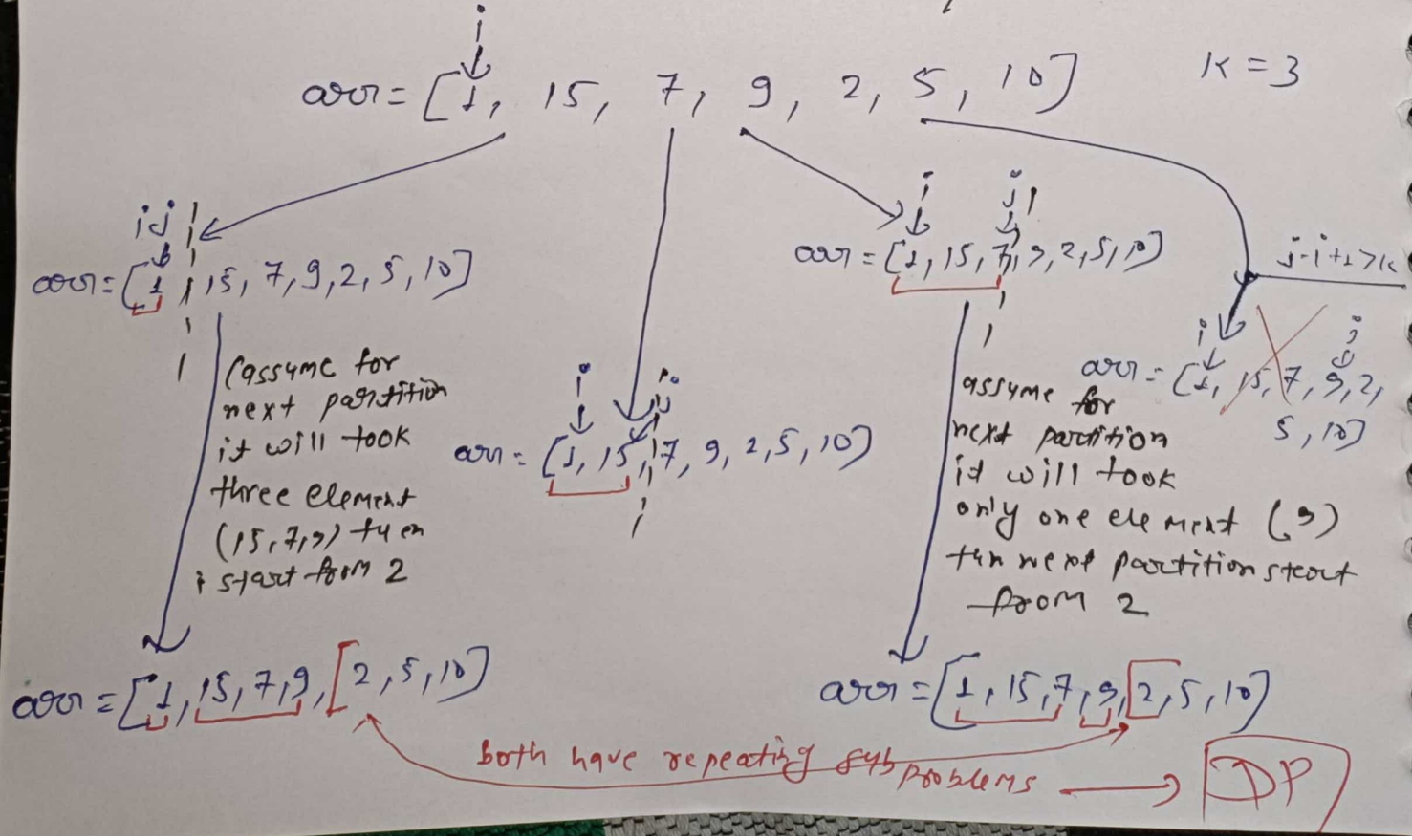
explaintion

arr = [1, 15 | 7, 9, 2 | 5, 10] K=3
15 15 | 9 9 9 | 10 10

arr = [1, 15, 7 | 9, | 2, 5, 10]
15 15 15 | 9 | 10 10 10

soln

We try all possibility to place the partition & store + Bars [
↳ T.C. = exponential / factorial



Imp

Recursion :-

```
j = i ;   j < min(n, i+k) ;   j++
```

why we wrote this

Sub array size can be at most K so $j < i+K$

if our i is at

[1, 15, 7, 9, 2, 5, 10]
 ↑
 i

$i+K = 6+3 = 9$.

but $n = 7$

so $j < n$

T.C. = exponential / factorial

~~memoization~~

solve (i, arr) → This will return the largest sum of the given array after partition

or.

[1, 15, 7, 9, 2, 5, 10] , K = 3

solve (2, arr) → from 2nd index to n the largest sum after partition this fn will return

memoization

dp → 1D array

dp[i] = solve (i, arr);

Bottom-UP

Top-Down
start idx = 0 & go till (n-1)

in BOttom-UP
start idx = n-1 & go till 0

T.C. = O(n*k)
S.C. = O(n)

Space Optimization

we can optimize space from O(n) to O(k)
use dp[k+1] ← initially all values = 0.

Imp.

dp[i%k] = max(dp[i], (maxVal * (j-i+1) + dp[(j+1)%k]));

↳ this will give you wrong ans bcz.

6%2 = 0 & 8%2 = 0

↳ for two different values % can gives same answer.

In top-Down approach

$dp[i] = max(dp[i], (maxVal * (j - i + 1) + dp[j + 1]))$

These both are different but in MCM we while space optimization these can be same at some point.

- So we use ans variable