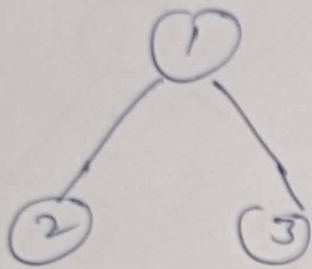
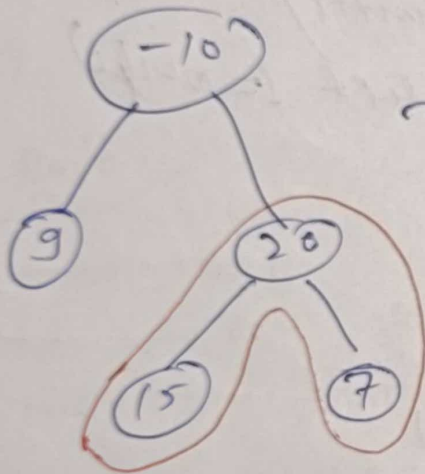


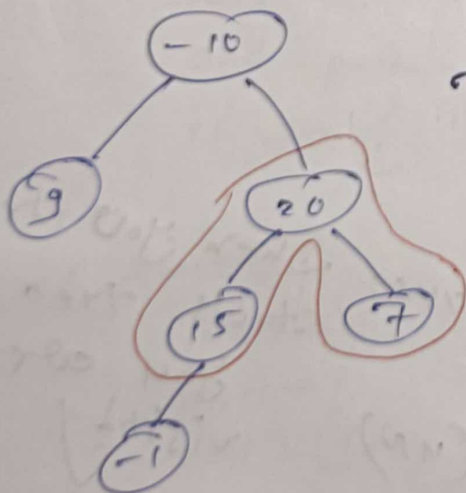
## 124. Binary Tree Maximum Path Sum



maximum path sum =  $2 \rightarrow 1 \rightarrow 3$   
 $2 + 1 + 3 = 6$



maximum path sum  $\Rightarrow 15 \rightarrow 20 \rightarrow 7$   
 $= 42$



maximum path sum  $\Rightarrow 15 \rightarrow 20 \rightarrow 7$   
 $= 42$

### Approach - 1 Brute force approach.

Go for every possible path & whichever path is giving me maximum path sum we pick up that path.

Now if we have  $N$  nodes  
nodes  $\rightarrow N$

total paths possible for every node  $\rightarrow N^2$

that means for  $N$  nodes  $\rightarrow$

$$\boxed{\text{T.C.} = O(N \times N^2) = O(N^3)}$$

This is not an optimal sol<sup>n</sup>

Approach - 2:

```
int solve (node, max_path_sum)
{
    if (node == NULL)
        return 0;
```

```
    int left_sum = max(0, solve(node->left, max_path_sum));
```

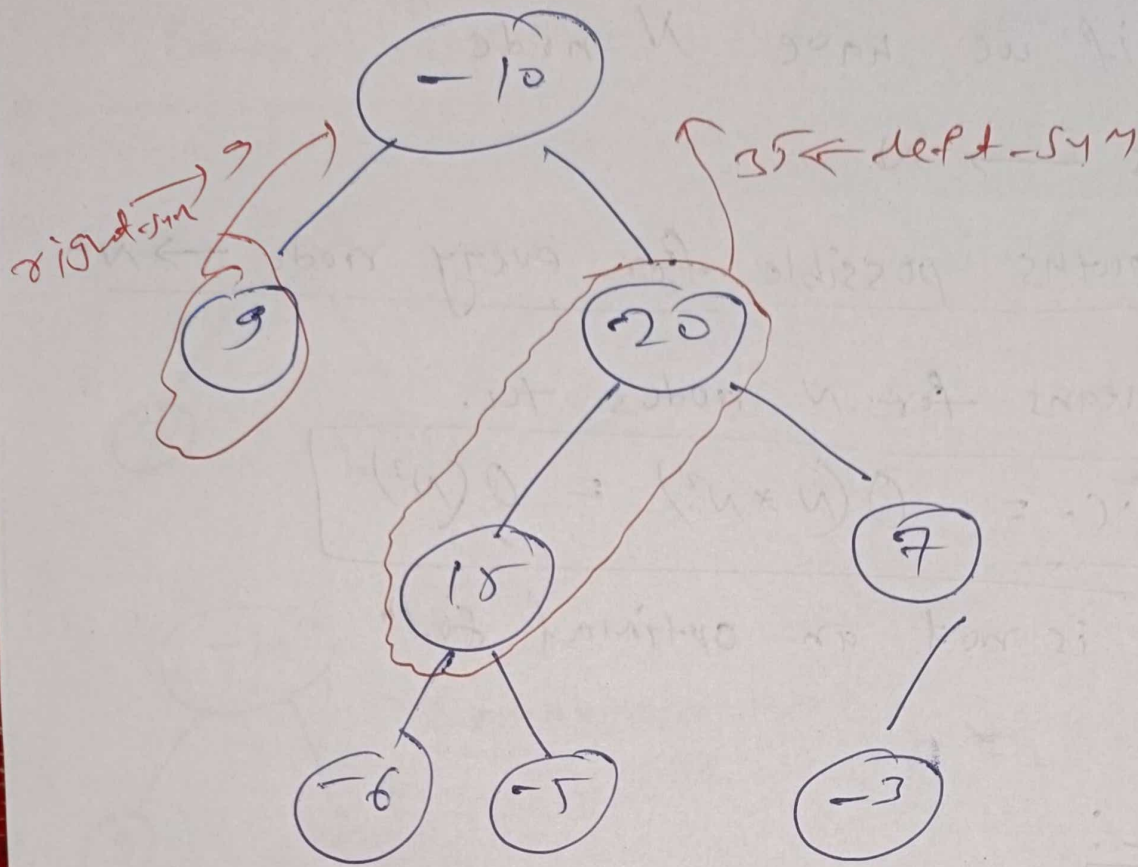
```
    int right_sum = max(0, solve(node->right, max_path_sum));
```

```
    max_path_sum = max(max_path_sum, (node->val + left_sum + right_sum));
```

```
    return node->val + max(left_sum, right_sum);
```

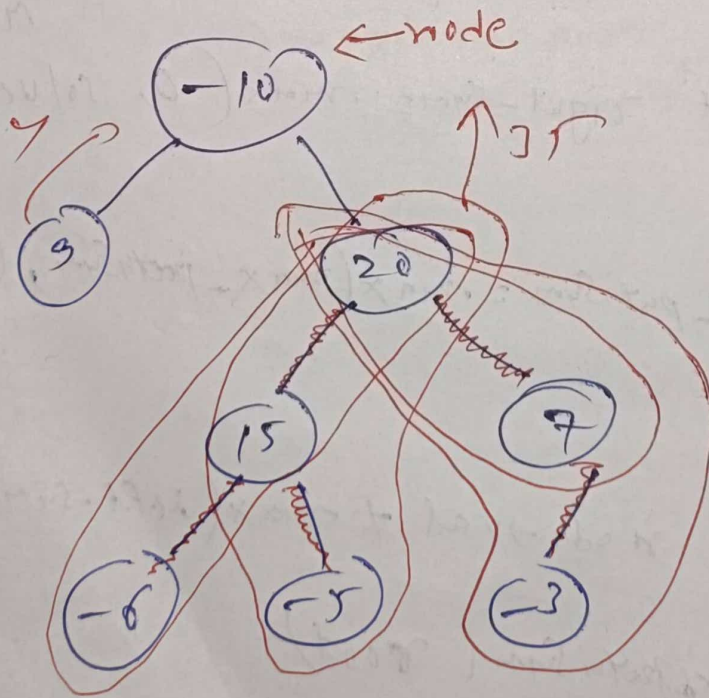
```
int max_path_sum (root)
```

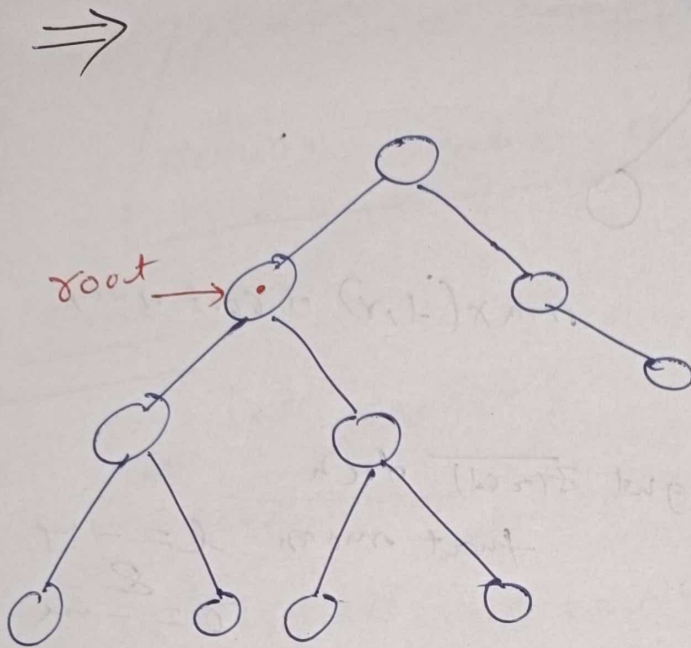




$\text{left-sum} = \text{node \& left sub tree \& possible paths}$   
 $\text{h} \uparrow \text{max put sum store in set}$

$\text{right-sum} = \text{sum of left sum}$





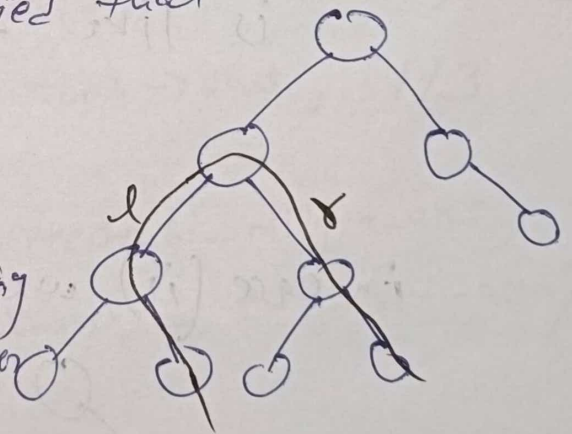
write down all possibilities that which path gives you best answer.

(i) left side, right side  $\Rightarrow$   $\overline{\text{actual sum}}$   $\Rightarrow$   $\text{root}$   $\Rightarrow$   $\text{val}$

That means you are satisfied that this is your path sum.

$$l + r + \text{root} \rightarrow \text{val}$$

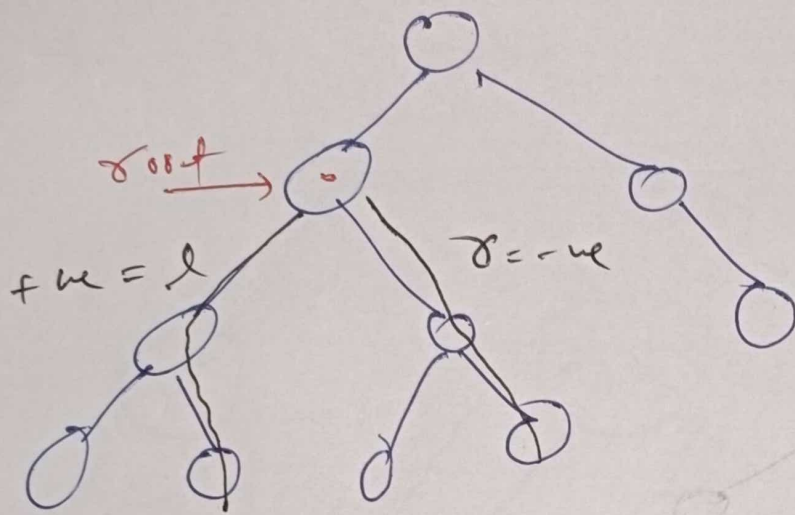
By this we are deciding that  $\Rightarrow$   $\text{actual answer}$



(ii) left  $\Rightarrow$  right  $\Rightarrow$   $\overline{\text{actual sum}}$   $\Rightarrow$   $\text{val}$   $\Rightarrow$   $\text{root}$   $\Rightarrow$   $\text{val}$  means if get  $-ve$  in left or right  $\Rightarrow$  right and we add this in  $\Rightarrow$   $\text{root} \rightarrow \text{val}$  this will decrease  $\Rightarrow$   $\text{own sum}$

$$\max(l, r) + \text{root} \rightarrow \text{val}$$





$$\max(l, r) + \text{root} + \text{val}$$

(iii) if left, if right ~~tree~~ then  
that means  $l = -\infty$   
 $\text{root} + \text{val}$   $r = \infty$

(iv) in (i) case path is like .

and we are not returning this because  
if we return this that means ~~the~~ the path  
is like that

but this is not  
be true

in case (ii) we can return  $\max(l, r) + \text{root} + \text{val}$

because we are hopping  
from upper node to  
this node and not  
skip it,

same as in case (iii).

```

    maxSum = max(maxSum, 1, 2, 3)
    return max(2, 3);

```

code

```

int maxSum;

int solve (TreeNode* root)
{
    if (root == NULL)
        return 0;

    int l = solve (root -> left);
    int r = solve (root -> right);

    int neeche-hi-milgya-ans = l + r + root -> val; // 1
    int koi-ek-acha = max(l, r) + root -> val; // 2
    int only-root-acha = root -> val; // 3

    maxSum = max({maxSum, neeche-hi-milgya-ans,
                  koi-ek-acha, only-root-acha});

    return max(koi-ek-acha, only-root-ans);
}

int maxPathSum (TreeNode* root);
{
    maxSum = INT_MIN;
    solve (root);
    return maxSum;
}

```