

Buy and Sell Stocks III

Part II \rightarrow Infinite Transaction

part III \rightarrow 2 Transaction

Where we are putting the bound & this we can relate to the knapsack where we are bounding of our bag capacity

$f(idx, buy, cap)$ \rightarrow 2 (initially)

{
if (cap == 0) return 0;
if (idx == n) return 0;

if (buy)
{
return max {
- prices[idx] + f(idx+1, 0, cap);
0 + f(idx+1, 1, cap)}}

\rightarrow money goes into market

else
{
return max {
prices[idx] + f(idx+1, 1, cap-1);
0 + f(idx+1, 0, cap);
}}

\rightarrow money get from market

Memoization

<u>States</u>	idx	o/l buy	cap
	↓	↓	↓
	0 to N	2	0/1/2 = 3

$$dp[N][2][3]$$

$$T.C. = O(N * 2 * 3)$$

$$S.C. = O(N * 2 * 3) + O(N)$$

Tabulation

① Base cases

② Changing parameters

;

buy

cap

③ copy paste recurrence

④ Base cases

(i) cap = 0 that means idx & buy can be anything

for (idx = 0 to n-1)

for (buy = 0 to 1)

$$dp[idx][buy][0] = 0$$

(ii) $idx = n$ then buy & cap can be anything

for (buy = 0 to 1)

for (cap = 0 to 2)

$dp[n][buy][cap] = 0$

(2) Changing parameters

$i \rightarrow n-1$ to 0

buy $\rightarrow 0$ to 1

cap $\rightarrow 0$ to 2

(3) copy past recurrence

(4) return $dp[0][1][2]$

other sol

$dp[n][4]$ 2 transaction

$f(idx, transaction)$

B	S	B	S
0	1	2	3

if (idx == n || transaction == 4)

return 0

return 0

Buy

|| = odd for sell

if (transaction % 2 == 0) // buy

else // sell

max

$f(idx+1, transaction)$

$-prices[idx] + f(idx+1, transaction)$

else //sell

$\max \{ \text{prices}[\text{idx}] + f(\text{idx}+1, \text{trans}+1);$
 $f(\text{idx}+1, \text{trans}); \}$

$dp[N][4]$