### **Operations ... SQL Specifics**



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- What is difference between following?
  - SELECT dno FROM EMPLOYEE
  - SELECT DISTINCT dno FROM EMPLOYEE
- What is meaning of
  - COUNT(\*)
  - COUNT(DISTINT dno)

## Issues with Null Values

- An attribute having NULL could mean either of following-
  - Value is unknown or not available right now
  - Value is not application for the tuple: a employee not having supervisor will have null in this attribute
- Consider following two SQL statements-

```
SELECT e.salary*1.1 from employee as e;
Select * from employee as e where e.salary > 50000;
```

 Interpret e here as tuple variable that ranges over all tuples of employee relations. Try finding result of expressions in blue for tuples where salary is NULL?

## Issues with Null Values

- Arithmetic expressions (+,-,\*,/) involving null values result null value for result
- When NULL values appears for attributes used in WHERE clause then boolean expression like this t.a < 10 then interpretation of attribute reference is UNKOWN.</li>
- When we compare a NULL value with another value including NULL, result is UNKNOWN.
- UNKOWN is treated as third truth (in addition to TRUE and FALSE) value in SQL where clause evaluation

# Truth values for UNKOWN

- NOT
  - NOT UNKOWN -> UNKWON
- AND
  - TRUE AND UNKOWN -> UNKOWN
  - FALSE AND UNKOWN -> FALSE
  - UNKWON AND UNKOWN -> UNKOWN
- OR
  - TRUE OR UNKOWN -> TRUE
  - FALSE OR UNKOWN -> UNKOWN
  - UNKWON OR UNKOWNUNKOWN

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### **Null Values and Comparisons**

- While evaluating WHERE clause tuples with UNKOWN or FALSE truth values are not included in result
- Following query will not include any tuple where either of value in NULL irrespective value in other attribute

```
SELECT * FROM EMPLOYEE WHERE bdate < DATE '2001-01-01' AND salary > 30000
```

 Following query will not include a tuple only when both are NULL, if one of attribute meets the condition then it will get included in result

```
SELECT * FROM EMPLOYEE WHERE bdate < DATE '2001-01-01' OR salary > 30000
```

# Null Values and Comparisons - IS NULL

Following will not give desired result. Why? -

```
SELECT * FROM employee

WHERE superssn = NULL;
```

- This is so because Null = Null is also UNKOWN. For checking an attribute for having NULL value, SQL provides IS NULL (and IS NOT NULL)
- We write as following for such situations –

```
SELECT * FROM employee
WHERE superssn IS NULL ;
```



### **Sub-queries in SQL**

# Subquery in SQL

 A Query that is part of another query is subquery. A subquery may also have subquery, and so forth upto any level

 A subquery in SQL is written as a query expression enclosed in parentheses, and is in following form-

```
"(SELECT ... FROM ...)" as a part of some existing query
```

Result of sub-query is again a relation;

### **Subquery in WHERE clause**

We have seen sub-query in IN, as
 WHERE ... IN (SELECT ... )

For example,
 SELECT \* FROM employee WHERE ssn IN
 ( SELECT essn FROM works on );

How such a query is evaluated?

# **Execution of Subquery**

- SUB-Query may not execute for every tuple of outer query
- Consider another query-

 Typically, after execution of inner query, outer query may be translated to:

```
SELECT * FROM student WHERE
progid IN (BEC, BEE);
```

 However this optimization may not be possible when you have "correlated sub-query"

## Correlated Sub-Queries

 When inner query makes a reference to tuple of outer query then it is correlated sub-query. Consider following query -

 List employees, whose salary is more than department average:

```
SELECT ssn, fname FROM employee as e
WHERE salary > (SELECT AVG(salary) FROM
employee WHERE dno = e.dno)
```



#### **Execution of Correlated Sub-Queries**

Consider same query

```
SELECT ssn, fname FROM employee as e WHERE salary > (SELECT AVG(salary) FROM employee WHERE dno = e.dno)
```

- Logically, it is as following: For each tuple of outer query, execute inner query.
- Note that it can not be executed once for all tuples of outer query, as the case be with un-related inner query, and we have to execute SUB-Query for every tuple of outer query
- This is identified problem with correlated sub-queries.



### Correlated Sub-Queries could be expensive to execute – therefore should be avoided

- Correlated queries are expensive to execute, and can be avoided; for example the previous example
- SELECT ssn, fname FROM employee as e WHERE salary > (SELECT AVG(salary) FROM employee WHERE dno = e.dno)
- can be re-written as-SELECT ssn, fname, salary FROM employee as e NATURAL JOIN (SELECT dno, AVG(salary) as avg\_sal FROM employee GROUP BY dno) as av WHERE salary > av.avg\_sal;



### more Correlated Sub-queries

List down employees having salary greater than their immediate supervisors.

```
select * from employee as e1 where e1.salary >
  (select salary from employee as e2 where e2.ssn =
e1.superssn);
```

- Select employees having dependents older than 18 years:
   SELECT \* FROM employee AS e WHERE ssn IN (SELECT essn FROM dependent AS d WHERE essn = e.ssn AND age(d.bdate) > interval '18 years');
- Attempt re-writting them without correlated query.

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## Compare a values with a bag of values (SQL)

 For example consider following two queries [Find out employee who have salary greater some or all employees of dno = 4]

SELECT ssn, fname FROM employee WHERE salary

> **SOME** (SELECT salary FROM employee WHERE dno = 4);

SELECT ssn, fname FROM employee WHERE salary

> ALL (SELECT salary FROM employee WHERE dno = 4);

# Compare a values with a bag of values (SQL)

- Note the equivalences:
  - SELECT ssn, fname FROM employee WHERE salary
  - > SOME (SELECT salary FROM employee WHERE dno = 4); and SELECT ssn, fname FROM employee WHERE salary
    - > (SELECT min(salary) FROM employee WHERE dno = 4);

SELECT ssn, fname FROM employee WHERE salary

- > ALL (SELECT salary FROM employee WHERE dno = 4); and SELECT ssn, fname FROM employee WHERE salary
  - > (SELECT max(salary) FROM employee WHERE dno = 4);

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### Compare a values with a bag of values (SQL)

Comparative operators could be, one of following-

```
>SOME, >=SOME, <=SOME, =SOME,
<>SOME
>ALL, >=ALL, <=ALL, <ALL, =ALL, <>ALL
```

- Note: it can be easily proved that
  - **=SOME** is identical to **IN**, and
  - <>SOME is not identical to NOT IN
  - = ALL is not identical to IN
  - **ALL** (mean = NONE) and is same as **NOT IN**, and Earlier versions of SQL used **ANY** for **SOME**; today both keywords are used as synonymous.



### **Sub-queries in Update statements**

```
• UPDATE employee
   SET salary = salary * 1.1
   WHERE ssn IN ( ... );
```

```
• DELETE employee WHERE ssn = ( ... );
```

## **EXISTS and NOT EXISTS in SQL**

- Checks for emptiness of a relation and returns true or false.
- EXISTS (r) can be interpreted as "is there some tuple exists in relation r"
- EXISTS (r) returning true says that argument relation r is not empty
- Similarly, **NOT EXISTS (r)** returning true says that argument relation r empty

# Example EXISTS

List employees who have dependents older than 18 years

```
    SELECT * FROM employee AS e WHERE EXISTS
        (SELECT * FROM dependent AS d WHERE d.essn
        = e.ssn AND age(d.bdate) > interval '18
        years');
```

# SQL- EXISTS and IN

- While they might appear to be serving similar purposes, semantically are different.
- Both appear as part of predicate in WHERE clause of SELECT
- IN:
  - Syntax: x IN ( r )
  - Meaning: checks existence of tuple x in relation r, if found returns true, other wise false. Normally x is a scalar value and r is a single column relation.
- EXISTS:
  - Syntax: EXISTS ( r )
  - Meaning: checks if r is a non empty relation. Returns true if the relation has at least one tuple, otherwise false.
- In both above cases  ${f r}$  is a *relational expression* resulting a relation.



#### **Bags and Relational Operations**

# Relational Operations and multiset (or bag)

- By Definition, relations are set; but implementations may permit duplicate tuples and such relations are called *bags*
- Normally stored relations (base) relations should still be sets, because most relations have Primary Key
- However SQL SELECT results are often bags, possibly because duplicate removal is expensive.
- To get set you use DISTINCT keyword

## SQL and Multiset (or Bag)

- SET operations, that are UNION, INTESECT, and EXCEPT in SQL yield their result as SET, that means duplicates are removed
- SQL however provides options by which you can have bag results by adding ALL keyword to operation name, i.e. UNION ALL, EXCEPT ALL or so.
- Let us see an example-

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## UNION/INTERSECT/EXCEPT ALL in SQL

Compare result of following queries:

```
SELECT superssn FROM employee; --Q1
SELECT mgrssn FROM department; --Q2
SELECT superssn FROM employee
UNION
SELECT mgrssn FROM department; --Q3
SELECT superssn FROM employee
UNION ALL
SELECT mgrssn FROM department; --Q4
```



R	superssn
	102
	101
	101
	102
	101
	(Null)
	108
	108

S	mgrssn	
		101
		102
		108

R UNION S	superssn
	101
	102
	108
	(Null)

R UNION ALL S	superssn
	102
	101
	101
	102
	101
	(Null)
	108
	108
	101
	102
pecifics	108



R	superssn
	102
	101
	101
	102
	101
	(Null)
	108
	108

S	mgrssn	
		101
		102
		108

R EXCEPT S	superssn
	(Null)

R EXCEPT ALL S	superssn
	101
	101
	102
	108
	(Null)

# UNION/INTESECT/EXCEPT ALL in SQL

- UNION ALL
  - count of an element e in result is sum of count in R and S
- INTERSECT ALL
  - min(count-r, count-s) of an element in R and S, is taken as result
- EXCEPT ALL:
  - Every occurrence of an element e in S decreases its count in R by one.



### **Functions and Operators in SQL**

## Functions and Operators

 SQL provides various functions and operators that can be used to create a new attribute in resultant relations

 There are typically, type conversion, arithmetic operators, mathematical, and string manipulation operators and functions. For example: substring, upper, lower, sqrt, ln, etc.

 Details for PostgreSQL functions can be seen at: <a href="http://intranet.daiict.ac.in/~pm\_jat/postgres/html/functions.html">http://intranet.daiict.ac.in/~pm\_jat/postgres/html/functions.html</a>.



```
SELECT ssn,
  fname || ' ' || minit || '. ' || lname AS name,
  current date - bdate AS age FROM employee;
SELECT essn, hours*50 AS amount FROM works on;
SELECT upper (fname) AS name, ln(salary) AS x FROM
  employee;
SELECT * FROM employee
      WHERE upper(fname) = 'FRANKLIN';
SELECT essn FROM dependent WHERE age(d.bdate) > interval
  '18 years');
```

- **BETWEEN**, **LIKE** are used in predicate:
  - SELECT .... WHERE A BETWEEN 10 TO 20;
  - SELECT ... WHERE A1 LIKE '%IX%' OR A2 LIKE 'ABC%' OR A3 LIKE '%XYZ';
  - SELECT ... WHERE A1 LIKE '\_X\_%';
- Also: NOT BETWEEN and NOT LIKE.

# Regular Expression Matching in PostgreSQL

 PostgreSQL also allows regular expression matching in string match using IS SIMILAR TO <reg-ex>

# ORDER BY

- ORDER BY CLAUSE is used for ordering the resultant tuples of a SQL query.
- Following statements returns row-set from employee table, and rows are sorted based on salary. To order in descending order, we add DESC keyword after attribute name.

```
select * from employee order by salary;
select * from employee order by salary desc;
```

 Following statement returns row-set from employee table, and rows are sorted in ascending order of dno, and within dno all rows are sorted on salary in descending order-

select \* from employee order by dno, salary desc;

# LIMIT and OFFSET

- Examples below should be self explanatory
- Gives top three earners
  - select \* from employee order by salary desc limit 3
- Gives next two earners after top 3
  - select \* from employee order by salary desc offset 3 limit 2