

IT308: Operating Systems

Condition Variables

Review: Concurrency Objectives

- Mutual Exclusion – if one thread is in critical section then no other thread is
 - Solved using locks
- Ordering – B runs after A does something
 - Solved using condition variables

Example 1: Thread Join

```
pthread_t p1, p2;
```

```
// create child threads
```

```
pthread_create(&p1, NULL, mythread, "A");
```

```
pthread_create(&p2, NULL, mythread, "B");
```

```
...
```

```
// join waits for the child threads to finish
```

```
thr_join(p1, NULL);
```

```
thr_join(p2, NULL);
```

how to implement `thr_join()`?

```
return 0;
```

Waiting for an Event

- Parent thread has to wait until child terminates
- Option 1: spin until that happens
 - Waste of CPU time

Waiting for an Event

- Parent thread has to wait until child terminates
- Option 1: spin until that happens
 - Waste of CPU time
- Option 2: wait (sleep) in a queue until that happens
 - Better use of CPU time
 - Child thread will signal the parent to wake up before its termination

Generalizing Option 2

- **Condition Variable:** queue of waiting threads with two basic operations

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- Thread B waits for a signal on cv before running
 - `cond_wait(cv, ...)`

Generalizing Option 2

- **Condition Variable**: queue of waiting threads with two basic operations
- Thread B waits for a signal on cv before running
 - `cond_wait(cv, ...)`
- Thread A sends signal to cv to wake-up one waiting thread
 - `cond_signal(cv, ...)`

Thread join: Attempt 1

Parent

```
void thr_join() {  
    cond_wait(&c);  
}
```

Child

```
void thr_exit() {  
    cond_signal(&c);  
}
```

- Does this work? If not, what's the problem?

Thread join: Attempt 1

Parent

```
void thr_join() {  
    cond_wait(&c);  
}
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Child

```
void thr_exit() {  
    cond_signal(&c);  
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- Does this work? If not, what's the problem?
- Child may run and call `cond_signal()` before parent called `cond_wait()`

Thread join: Attempt 1

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void thr_join() {  
    cond_wait(&c);  
}
```

Child

```
void thr_exit() {  
    cond_signal(&c);  
}
```

- Does this work? If not, what's the problem?
- Child may run and call `cond_signal()` before parent called `cond_wait()`
 - `cond_signal()` signals nobody (parent is not there yet!)

Thread join: Attempt 1

Parent

```
void thr_join() {  
    cond_wait(&c);  
}
```

Child

```
void thr_exit() {  
    cond_signal(&c);  
}
```

- Does this work? If not, what's the problem?
- Child may run and call `cond_signal()` before parent called `cond_wait()`
 - `cond_signal()` signals nobody (parent is not there yet!)
 - parent goes to sleep indefinitely

Thread join: Attempt 2

Parent

```
void thr_join() {  
    if (done == 0) {  
        cond_wait(&c);  
    }  
}
```

Child

```
void thr_exit() {  
    done = 1;  
    cond_signal(&c);  
}
```

- Let's keep some state then

Thread join: Attempt 2

Parent

```
void thr_join() {  
    if (done == 0) {  
        cond_wait(&c);  
    }  
}
```

Child

```
void thr_exit() {  
    done = 1;  
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}
```

- Let's keep some state then
- Is there a problem here?

Thread join: Attempt 2

Parent

```
void thr_join() {  
    if (done == 0) {  
        cond_wait(&c);  
    }  
}
```

Child

```
void thr_exit() {  
    done = 1;  
    cond_signal(&c);  
}
```

- Trace

Parent: a b

Child: x y

Thread join: Attempt 2

Parent

```
void thr_join() {  
    if (done == 0) { //a  
        cond_wait(&c); //b  
    }  
}
```

Child

```
void thr_exit() {  
    done = 1; //x  
    cond_signal(&c); //y  
}
```

- Trace

Parent: a b

Child: x y

- Again, parent may sleep indefinitely
- Solution?

Using Locks to Achieve Atomicity

- `wait(cond_t *cv, mutex_t *lock)`
 - assumes the lock is held when `wait()` is called
 - puts caller to sleep + releases the lock (atomically)
 - when awoken, reacquires lock before returning
- `signal(cond_t *cv)`
 - wake a single waiting thread (if ≥ 1 thread is waiting)
 - if there is no waiting thread, just return, doing nothing

Thread join: Attempt 3

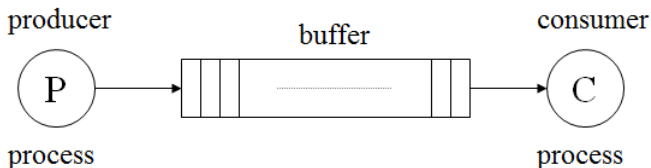
Parent:

```
void thread_join() {  
    Mutex_lock(&m);           // w  
    if (done == 0)           // x  
        Cond_wait(&c, &m); // y  
    Mutex_unlock(&m);        // z  
}
```

Child:

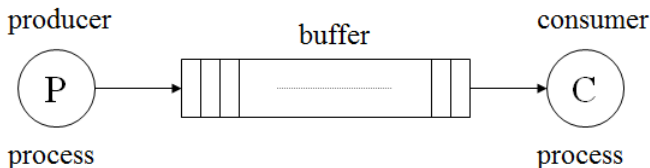
```
void thread_exit() {  
    Mutex_lock(&m);           // a  
    done = 1;                // b  
    Cond_signal(&c);         // c  
    Mutex_unlock(&m);        // d  
}
```

Producer/Consumer (bounded buffer) Problem



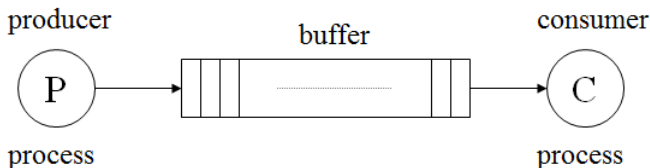
- Assume shared, finite size buffer
- from time to time, the producer adds items to buffer
- the consumer removes items from buffer

Producer/Consumer (bounded buffer) Problem



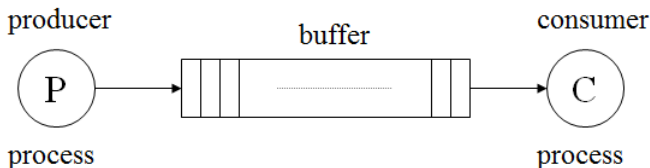
- Assume shared, finite size buffer
- from time to time, the producer adds items to buffer
- the consumer removes items from buffer
- careful synchronization required

Producer/Consumer (bounded buffer) Problem



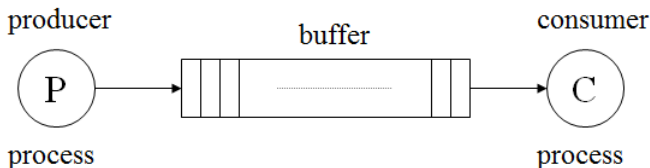
- Assume shared, finite size buffer
- from time to time, the producer adds items to buffer
- the consumer removes items from buffer
- careful synchronization required
 - the consumer must wait when the buffer is empty

Producer/Consumer (bounded buffer) Problem



- Assume shared, finite size buffer
- from time to time, the producer adds items to buffer
- the consumer removes items from buffer
- careful synchronization required
 - the consumer must wait when the buffer is empty
 - the producer must wait when the buffer is full

Producer/Consumer (bounded buffer) Problem



- Assume shared, finite size buffer
- from time to time, the producer adds items to buffer
- the consumer removes items from buffer
- careful synchronization required
 - the consumer must wait when the buffer is empty
 - the producer must wait when the buffer is full
- typical solution would use a shared variable count

The Put and Get Routines (Version 1)

- Assume buffer can hold only one item

```
1      int buffer;
2      int count = 0;    // initially, empty
3
4      void put(int value) {
5          assert(count == 0);
6          count = 1;
7          buffer = value;
8      }
9
10     int get() {
11         assert(count == 1);
12         count = 0;
13         return buffer;
14     }
```


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```

- Only put data into the buffer when count is zero (buffer is empty)

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10     int get() {
11         assert(count == 1);
12         count = 0;
13         return buffer;
14     }
```

- Only put data into the buffer when count is zero (buffer is empty)
- Only get data from the buffer when count is one (buffer is full)

Producer/Consumer Threads (Version 1)

- Using a condition variable (and mutex) to synchronise producer and consumer

```
1      cond_t cond;
2      mutex_t mutex;
3
4      void *producer(void *arg) {
5          int i, loops = (int) arg;
6          for (i = 0; i < loops; i++) {
7              Pthread_mutex_lock(&mutex);           // p1
8              if (count == 1)                       // p2
9                  Pthread_cond_wait(&cond, &mutex); // p3
10             put(i);                                // p4
11             Pthread_cond_signal(&cond);           // p5
12             Pthread_mutex_unlock(&mutex);         // p6
13         }
14     }
15
16     void *consumer(void *arg) {
17         int i, loops = (int) arg;
18         for (i = 0; i < loops; i++) {
19             Pthread_mutex_lock(&mutex);           // c1
```

Producer/Consumer Threads (Version 1)

```
20         if (count == 0)                                // c2
21             Pthread_cond_wait(&cond, &mutex);          // c3
22         int tmp = get();                                  // c4
23         Pthread_cond_signal(&cond);                     // c5
24         Pthread_mutex_unlock(&mutex);                   // c6
25         printf("%d\n", tmp);
26     }
27 }
```

- p1–p3: A producer waits for the buffer to be empty
- c1–c3: A consumer waits for the buffer to be full

Producer/Consumer Threads (Version 1)

```
20         if (count == 0)                                // c2
21             Pthread_cond_wait(&cond, &mutex);           // c3
22         int tmp = get();                                  // c4
23         Pthread_cond_signal(&cond);                      // c5
24         Pthread_mutex_unlock(&mutex);                    // c6
25         printf("%d\n", tmp);
26     }
27 }
```

- p1–p3: A producer waits for the buffer to be empty
- c1–c3: A consumer waits for the buffer to be full
- The above code works for 1P and 1C. Can you find a problematic timeline with 2 consumers (still 1 producer)?

Thread Trace: Broken Solution (Version 1)

```
void *consumer(void *arg) {
    int i;
    for (i = 0; i < loops; i++) {
        Pthread_mutex_lock(&mutex);           // c1
        if (count == 0)                       // c2
            Pthread_cond_wait(&cond, &mutex); // c3
        int tmp = get();                      // c4
        Pthread_cond_signal(&cond);           // c5
        Pthread_mutex_unlock(&mutex);         // c6
        printf("%d\n", tmp);
    }
}

void *producer(void *arg) {
    int i;
    for (i = 0; i < loops; i++) {
        Pthread_mutex_lock(&mutex);           // p1
        if (count == 1)                       // p2
            Pthread_cond_wait(&cond, &mutex); // p3
        put(i);                               // p4
        Pthread_cond_signal(&cond);           // p5
        Pthread_mutex_unlock(&mutex);         // p6
    }
}
```

T_{c1}	State	T_{c2}	State	T_p	State	Count	Comment
c1	Running		Ready		Ready	0	Nothing to get
c2	Running		Ready		Ready	0	
c3	Sleep		Ready		Ready	0	

Thread Trace: Broken Solution (Version 1)

```

void *consumer(void *arg) {
    int i;
    for (i = 0; i < loops; i++) {
        Pthread_mutex_lock(&mutex);          // c1
        if (count == 0)                      // c2
            Pthread_cond_wait(&cond, &mutex); // c3
        int tmp = get();                     // c4
        Pthread_cond_signal(&cond);          // c5
        Pthread_mutex_unlock(&mutex);        // c6
        printf("%d\n", tmp);
    }
}

void *producer(void *arg) {
    int i;
    for (i = 0; i < loops; i++) {
        Pthread_mutex_lock(&mutex);          // p1
        if (count == 1)                      // p2
            Pthread_cond_wait(&cond, &mutex); // p3
        put(i);                              // p4
        Pthread_cond_signal(&cond);          // p5
        Pthread_mutex_unlock(&mutex);        // p6
    }
}

```

T_{c1}	State	T_{c2}	State	T_p	State	Count	Comment
c1	Running		Ready		Ready	0	Nothing to get
c2	Running		Ready		Ready	0	
c3	Sleep		Ready		Ready	0	
	Sleep		Ready	p1	Running	0	Buffer now full T_{c1} awoken
	Sleep		Ready	p2	Running	0	
	Sleep		Ready	p4	Running	1	
	Ready		Ready	p5	Running	1	
	Ready		Ready	p6	Running	1	
	Ready		Ready	p1	Running	1	
	Ready		Ready	p2	Running	1	

Thread Trace: Broken Solution (Version 1)

```

void *consumer(void *arg) {
    int i;
    for (i = 0; i < loops; i++) {
        Pthread_mutex_lock(&mutex);          // c1
        if (count == 0)                      // c2
            Pthread_cond_wait(&cond, &mutex); // c3
        int tmp = get();                     // c4
        Pthread_cond_signal(&cond);          // c5
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        printf("%d\n", tmp);
    }
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void *producer(void *arg) {
    int i;
    for (i = 0; i < loops; i++) {
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        if (count == 1)                      // p2
            Pthread_cond_wait(&cond, &mutex); // p3
        put(i);                              // p4
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    }
}

```

T _{c1}	State	T _{c2}	State	T _p	State	Count	Comment
c1	Running		Ready		Ready	0	
c2	Running		Ready		Ready	0	
c3	Sleep		Ready		Ready	0	Nothing to get
	Sleep		Ready	p1	Running	0	
	Sleep		Ready	p2	Running	0	
	Sleep		Ready	p4	Running	1	Buffer now full
	Ready		Ready	p5	Running	1	T _{c1} awoken
	Ready		Ready	p6	Running	1	
	Ready		Ready	p1	Running	1	
	Ready		Ready	p2	Running	1	
	Ready		Ready	p3	Sleep	1	Buffer full; sleep
	Ready	c1	Running		Sleep	1	T _{c2} sneaks in ...
	Ready	c2	Running		Sleep	1	
	Ready	c4	Running		Sleep	0	... and grabs data
	Ready	c5	Running		Ready	0	T _p awoken
	Ready	c6	Running		Ready	0	

Thread Trace: Broken Solution (Version 1)

```

void *consumer(void *arg) {
    int i;
    for (i = 0; i < loops; i++) {
        Pthread_mutex_lock(&mutex);          // c1
        if (count == 0)                      // c2
            Pthread_cond_wait(&cond, &mutex); // c3
        int tmp = get();                     // c4
        Pthread_cond_signal(&cond);          // c5
        Pthread_mutex_unlock(&mutex);        // c6
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    }
}

```

T_{c1}	State	T_{c2}	State	T_p	State	Count	Comment
c1	Running		Ready		Ready	0	
c2	Running		Ready		Ready	0	
c3	Sleep		Ready		Ready	0	Nothing to get
	Sleep		Ready	p1	Running	0	
	Sleep		Ready	p2	Running	0	
	Sleep		Ready	p4	Running	1	Buffer now full
	Ready		Ready	p5	Running	1	T_{c1} awoken
	Ready		Ready	p6	Running	1	
	Ready		Ready	p1	Running	1	
	Ready		Ready	p2	Running	1	
	Ready		Ready	p3	Sleep	1	Buffer full; sleep
	Ready	c1	Running		Sleep	1	T_{c2} sneaks in ...
	Ready	c2	Running		Sleep	1	
	Ready	c4	Running		Sleep	0	... and grabs data
	Ready	c5	Running		Ready	0	T_p awoken
	Ready	c6	Running		Ready	0	
c4	Running		Ready		Ready	0	Oh oh! No data

Thread Trace: Broken Solution (Version 1)

- After the producer woke T_{c1} , but before T_{c1} ever ran, the state of the bounded buffer changed by T_{c2}

Thread Trace: Broken Solution (Version 1)

- After the producer woke T_{c1} , but before T_{c1} ever ran, the state of the bounded buffer changed by T_{c2}
- There is no guarantee that when the woken thread runs, the state will still be as desired \rightarrow Mesa semantics
 - Virtually every system ever built employs Mesa semantics

Thread Trace: Broken Solution (Version 1)

- After the producer woke T_{c1} , but before T_{c1} ever ran, the state of the bounded buffer changed by T_{c2}
- There is no guarantee that when the woken thread runs, the state will still be as desired \rightarrow Mesa semantics
 - Virtually every system ever built employs Mesa semantics
- Hoare semantics provides a stronger guarantee that the woken thread will run immediately upon being woken

Producer/Consumer (Version 2)

- Consumer T_{c1} wakes up and re-checks the state of the shared variable
- If the buffer is empty, the consumer simply goes back to sleep

```
1      cond_t cond;
2      mutex_t mutex;
3
4      void *producer(void *arg) {
5          int i, loops = (int) arg;
6          for (i = 0; i < loops; i++) {
7              Pthread_mutex_lock(&mutex);           // p1
8              while (count == 1)                    // p2
9                  Pthread_cond_wait(&cond, &mutex); // p3
10             put(i);                                // p4
11             Pthread_cond_signal(&cond);            // p5
12             Pthread_mutex_unlock(&mutex);          // p6
13         }
14     }
15
```

Producer/Consumer (Version 2)

```
(Cont.)
16  void *consumer(void *arg) {
17      int i, loops = (int) arg;
18      for (i = 0; i < loops; i++) {
19          Pthread_mutex_lock(&mutex);           // c1
20          while (count == 0)                    // c2
21              Pthread_cond_wait(&cond, &mutex); // c3
22          int tmp = get();                       // c4
23          Pthread_cond_signal(&cond);           // c5
24          Pthread_mutex_unlock(&mutex);         // c6
25          printf("%d\n", tmp);
26      }
27  }
```

- A simple rule to remember with condition variables is to always use while loops
- However, this code still has a bug!

Thread Trace: Broken Solution (Version 2)

```
void *consumer(void *arg) {
    int i;
    for (i = 0; i < loops; i++) {
        Pthread_mutex_lock(&mutex);          // c1
        while (count == 0)                   // c2
            Pthread_cond_wait(&cond, &mutex); // c3
        int tmp = get();                     // c4
        Pthread_cond_signal(&cond);          // c5
        Pthread_mutex_unlock(&mutex);        // c6
        printf("%d\n", tmp);
    }
}

void *producer(void *arg) {
    int i;
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    }
}
```

T_{c1}	State	T_{c2}	State	T_p	State	Count	Comment
c1	Running		Ready		Ready	0	Nothing to get
c2	Running		Ready		Ready	0	
c3	Sleep		Ready		Ready	0	

Thread Trace: Broken Solution (Version 2)

```

void *consumer(void *arg) {
    int i;
    for (i = 0; i < loops; i++) {
        Pthread_mutex_lock(&mutex);          // c1
        while (count == 0)                   // c2
            Pthread_cond_wait(&cond, &mutex); // c3
        int tmp = get();                     // c4
        Pthread_cond_signal(&cond);          // c5
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        printf("%d\n", tmp);
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void *producer(void *arg) {
    int i;
    for (i = 0; i < loops; i++) {
        Pthread_mutex_lock(&mutex);          // p1
        while (count == 1)                   // p2
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        put(i);                             // p4
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    }
}

```

T_{c1}	State	T_{c2}	State	T_p	State	Count	Comment
c1	Running		Ready		Ready	0	
c2	Running		Ready		Ready	0	
c3	Sleep		Ready		Ready	0	Nothing to get
	Sleep	c1	Running		Ready	0	
	Sleep	c2	Running		Ready	0	
	Sleep	c3	Sleep		Ready	0	Nothing to get

Thread Trace: Broken Solution (Version 2)

```

void *consumer(void *arg) {
    int i;
    for (i = 0; i < loops; i++) {
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        while (count == 0)                   // c2
            Pthread_cond_wait(&cond, &mutex); // c3
        int tmp = get();                     // c4
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    }
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void *producer(void *arg) {
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    for (i = 0; i < loops; i++) {
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    }
}

```

T_{c1}	State	T_{c2}	State	T_p	State	Count	Comment
c1	Running		Ready		Ready	0	
c2	Running		Ready		Ready	0	
c3	Sleep		Ready		Ready	0	Nothing to get
	Sleep	c1	Running		Ready	0	
	Sleep	c2	Running		Ready	0	
	Sleep	c3	Sleep		Ready	0	Nothing to get
	Sleep		Sleep	p1	Running	0	
	Sleep		Sleep	p2	Running	0	
	Sleep		Sleep	p4	Running	1	Buffer now full
	Ready		Sleep	p5	Running	1	T_{c1} awoken
	Ready		Sleep	p6	Running	1	
	Ready		Sleep	p1	Running	1	
	Ready		Sleep	p2	Running	1	
	Ready		Sleep	p3	Sleep	1	Must sleep (full)

Thread Trace: Broken Solution (Version 2)

```

void *consumer(void *arg) {
    int i;
    for (i = 0; i < loops; i++) {
        Pthread_mutex_lock(&mutex);           // c1
        while (count == 0)                    // c2
            Pthread_cond_wait(&cond, &mutex); // c3
        int tmp = get();                      // c4
        Pthread_cond_signal(&cond);           // c5
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        printf("%d\n", tmp);
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```

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	Sleep	c2	Running		Ready	0	
	Sleep	c3	Sleep		Ready	0	Nothing to get
	Sleep		Sleep	p1	Running	0	
	Sleep		Sleep	p2	Running	0	
	Sleep		Sleep	p4	Running	1	Buffer now full
	Ready		Sleep	p5	Running	1	T _{c1} awoken
	Ready		Sleep	p6	Running	1	
	Ready		Sleep	p1	Running	1	
	Ready		Sleep	p2	Running	1	
	Ready		Sleep	p3	Sleep	1	Must sleep (full)
c2	Running		Sleep		Sleep	1	Recheck condition
c4	Running		Sleep		Sleep	0	T _{c1} grabs data
c5	Running		Ready		Sleep	0	Oops! Woke T _{c2}
c6	Running		Ready		Sleep	0	
c1	Running		Ready		Sleep	0	
c2	Running		Ready		Sleep	0	
c3	Sleep		Ready		Sleep	0	Nothing to get

Thread Trace: Broken Solution (Version 2)

```

void *consumer(void *arg) {
    int i;
    for (i = 0; i < loops; i++) {
        Pthread_mutex_lock(&mutex);          // c1
        while (count == 0)                    // c2
            Pthread_cond_wait(&cond, &mutex); // c3
        int tmp = get();                      // c4
        Pthread_cond_signal(&cond);           // c5
        Pthread_mutex_unlock(&mutex);         // c6
        printf("%d\n", tmp);
    }
}

void *producer(void *arg) {
    int i;
    for (i = 0; i < loops; i++) {
        Pthread_mutex_lock(&mutex);          // p1
        while (count == 1)                    // p2
            Pthread_cond_wait(&cond, &mutex); // p3
        put(i);                              // p4
        Pthread_cond_signal(&cond);           // p5
        Pthread_mutex_unlock(&mutex);         // p6
    }
}

```

T _{c1}	State	T _{c2}	State	T _p	State	Count	Comment
c1	Running		Ready		Ready	0	
c2	Running		Ready		Ready	0	
c3	Sleep		Ready		Ready	0	Nothing to get
	Sleep	c1	Running		Ready	0	
	Sleep	c2	Running		Ready	0	
	Sleep	c3	Sleep		Ready	0	Nothing to get
	Sleep		Sleep	p1	Running	0	
	Sleep		Sleep	p2	Running	0	
	Sleep		Sleep	p4	Running	1	Buffer now full
	Ready		Sleep	p5	Running	1	T _{c1} awoken
	Ready		Sleep	p6	Running	1	
	Ready		Sleep	p1	Running	1	
	Ready		Sleep	p2	Running	1	
	Ready		Sleep	p3	Sleep	1	Must sleep (full)
c2	Running		Sleep		Sleep	1	Recheck condition
c4	Running		Sleep		Sleep	0	T _{c1} grabs data
c5	Running		Ready		Sleep	0	Oops! Woke T _{c2}
c6	Running		Ready		Sleep	0	
c1	Running		Ready		Sleep	0	
c2	Running		Ready		Sleep	0	
c3	Sleep		Ready		Sleep	0	Nothing to get
	Sleep	c2	Running		Sleep	0	
	Sleep	c3	Sleep		Sleep	0	Everyone asleep...

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Producer/Consumer (Version 3)

- A consumer should not wake other consumers, only producers, and vice-versa
- Can't use the same cond var for two things (signaling the buffer is empty and the buffer is full)