

Surabhi Narayan

Department of Computer Science & Engineering



DECREASE AND CONQUER

Surabhi Narayan

Department of Computer Science & Engineering

Decrease and Conquer

Combinatorial Objects

- > Permutations
- **≻**Combinations
- ➤ Subsets of a given set



Decrease and Conquer

Generating Permutations

- ➤ Underlying set elements are to be permuted
- ➤ Decrease and conquer approach
- > Satisfies the minimal change requirement
- ➤ Example: Johnson- Trotter algorithm



Decrease and Conquer

Generating Permutations

ALGORITHM JohnsonTrotter(n)

```
//Implements Johnson-Trotter algorithm for generating permutations
//Input: A positive integer n
//Output: A list of all permutations of \{1, \ldots, n\}
initialize the first permutation with 1 \ 2 \ldots n

while the last permutation has a mobile element do
find its largest mobile element k
swap k with the adjacent element k's arrow points to
reverse the direction of all the elements that are larger than k
add the new permutation to the list
```



Decrease and Conquer



| 1234 | 1243 | 1423 | 4 1 2 3 |
|---|--|--|---|
| $\frac{1}{4}$ $\frac{2}{1}$ $\frac{3}{3}$ $\frac{4}{2}$ | $\begin{array}{c} 1 & 2 & 4 & 5 \\ \hline 1 & 4 & 3 & 2 \end{array}$ | $\begin{array}{c} 1 & 4 & 2 & 3 \\ \hline 1 & 3 & 4 & 2 \end{array}$ | $\begin{array}{c} \stackrel{4}{\cancel{}} \stackrel{1}{\cancel{}} \stackrel{2}{\cancel{}} \stackrel{3}{\cancel{}} \\ 1 & 3 & 2 & 4 \end{array}$ |
| $\frac{2}{3}$ $\frac{2}{1}$ $\frac{2}{4}$ | $\frac{2}{3}$ $\frac{2}{4}$ $\frac{2}{2}$ | $\frac{2}{3}$ | $\frac{\cancel{}$ |
| $\overrightarrow{4}\overrightarrow{3}\overrightarrow{2}\overrightarrow{1}$ | $\overrightarrow{3}\overrightarrow{4}\overset{\leftarrow}{2}\overset{\leftarrow}{1}$ | $\overrightarrow{3} \overset{\leftarrow}{2} \overset{\leftarrow}{4} \overset{\leftarrow}{1}$ | $\overrightarrow{3} \overset{\leftarrow}{2} \overset{\leftarrow}{1} \overset{\rightarrow}{4}$ |
| $2\overrightarrow{3}\overrightarrow{1}\overrightarrow{4}$ | 2341 | 2431 | 4231 |
| $\overrightarrow{4} \overrightarrow{2} \overrightarrow{1} \overrightarrow{3}$ | $24\overline{13}$ | $2\overrightarrow{1}\overrightarrow{4}\overrightarrow{3}$ | $2\overrightarrow{1}\overrightarrow{3}\overrightarrow{4}$ |

Decrease and Conquer



Generating Subsets:

Knapsack problem needed to find the most valuable subset of items that fits a knapsack of a given capacity.

Powerset: set of all subsets of a set. Set $A=\{1, 2, ..., n\}$ has 2^n subsets.

Generate all subsets of the set $A=\{1, 2, ..., n\}$.

```
Any decrease-by-one idea? # of subsets of \{\} = 2^0 = 1, which is \{\} itself Suppose, we know how to generate all subsets of \{1,2,...,n-1\} Now, how can we generate all subsets of \{1,2,...,n\}?
```

Decrease and Conquer



Generating Subsets:

```
All subsets of \{1,2,...,n-1\}: 2^{n-1} such subsets
```

```
All subsets of \{1,2,...,n\}:
 2^{n-1} subsets of \{1,2,...,n-1\} and
 another 2^{n-1} subsets of \{1,2,...,n-1\} having 'n' with them.
```

That adds up to all 2^n subsets of $\{1,2,...,n\}$

Decrease and Conquer



Alternate way of Generating Subsets:

Knowing the binary nature of either having **n**th element or not, any idea involving binary numbers itself?

One-to-one correspondence between all 2^n bit strings $b_1b_2...b_n$ and 2^n subsets of $\{a_1, a_2, ..., a_n\}$.

Each bit string $b_1b_2...b_n$ could correspond to a subset.

In a bit string $b_1b_2...b_n$, depending on whether b_i is 1 or 0, a_i is in the subset or not in the subset.

| 000 | 001 | 010 | 011 | 100 | 101 | 110 | 111 |
|-----|-----------|-----------|----------------|-----------|----------------|----------------|---------------------|
| Ø | $\{a_3\}$ | $\{a_2\}$ | $\{a_2, a_3\}$ | $\{a_1\}$ | $\{a_1, a_3\}$ | $\{a_1, a_2\}$ | $\{a_1, a_2, a_3\}$ |

Decrease and Conquer



Generating Subsets in Squashed order:

Squashed order: any subset involving a_j can be listed only after all the subsets involving $a_1, a_2, ..., a_{j-1}$

Both of the previous methods does generate subsets in squashed order.

000 001 010 011 100 101 110 111
$$\varnothing$$
 { a_3 } { a_2 } { a_2 , a_3 } { a_1 } { a_1 , a_3 } { a_1 , a_2 } { a_1 , a_2 , a_3 }

Decrease and Conquer



Generating Subsets in Squashed order:

Squashed order: any subset involving a_j can be listed only after all the subsets involving $a_1, a_2, ..., a_{i-1}$

Can we do it with minimal change in bit-string (actually, just one-bit change to get the next bit string)? This would mean, to get a new subset, just change one item (remove one item or add one item).

Binary reflected gray code:

000 001 011 010 110 111 101 100



THANK YOU

Surabhi Narayan

Department of Computer Science & Engineering

surabhinarayan@pes.edu