· Turing machine (algorithm) for checking if a string (of zonoes) is of the form 02" (i.e, 2" 2010es, n7,0) Language A : {02": n 70}

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Algorithm outline: (idea) - if total number of zeroes is odd, we can reject right away! conty 2 states needed) But that is not enough. can we extend this idea? . If only one zono, accept.

What else?

-when it is not a single zero, (in a loop) cross off every alternative zero. so half the no of zeroes remain. if this number is odd (and not 1), reject! if it is 1 (single 0), accept! -we have an algorithm!

36)

U, L (hit right (first 2010) end, even) (ends in odd Unu, R no. of zeroes)

(37)

## -understand how it works:

- If in any round, after crossing of alternative zeroes, we hit the right end at an odd length, it rejects.
- when it is a valid string, what happens? say, 0000.

21, tape: 40000 (U × 92 00 U) 92, tape: UX00U UXOOU (jumps over 0) 92, Now, we start going backwards.

(39)

Goes to 24: UXOXU GOUX024XU) It jumps over all zeroes and x's. After 3 more moves: 24 4 10 X 1 Goes right: 21, 4000 q, uxoxu Jumps over: 0-x, gog to 22: - XXX

- . No more zeroes remaining.
- . on blank it goes to 24, then goes to 21.
- · Jumps over X's, goes right,
  - . hits blank (right end), accepts.

. Multitape turing machines: -Makes things simple. - Think of the first example, on each tape, move should be clear. I.E, S: QXTK-> QXTKX{L,R3K or Qxrk-- Qx([x(L,R])k

(42)

what it means is that -it has only one state at a time, - K tapes (and k tape heads), -so k "current" symbols. eg:  $\delta(q, a_1, a_2) = (r_1(b_1, L), (b_2, R))$ or  $(\gamma, b_1, b_2, b, R)$ tape 1 tape 2