

# Period Finding

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Mainly referenced from [NC00], [GHH08]

## 1 Problem

Given  $f : \mathbb{Z} \rightarrow \mathbb{N} \cup \{0\}$  such that

$$\begin{aligned} \forall x \ f(x) &= f(x+r) \quad \text{for some } r \in \mathbb{N} \\ \forall x, y \in \{0, \dots, r-1\} \ x \neq y &\rightarrow f(x) \neq f(y) \end{aligned}$$

find  $r$ .

## 2 Solution

### 2.1 The circuit

Build a quantum gate of  $(m+n)$  qubits,  $U : |x\rangle |y\rangle \rightarrow |x\rangle |y \oplus f(x)\rangle$  such that

$$2r^2 \leq N \leq M \tag{1}$$

where

$$M \stackrel{\text{def}}{=} 2^m, N \stackrel{\text{def}}{=} 2^n, \oplus \text{ is XOR.}$$

Using the above quantum gate, Hadamard gates  $H$ , and Quantum Fourier Transform gate  $QFT$  [Ros03], build the following circuit.

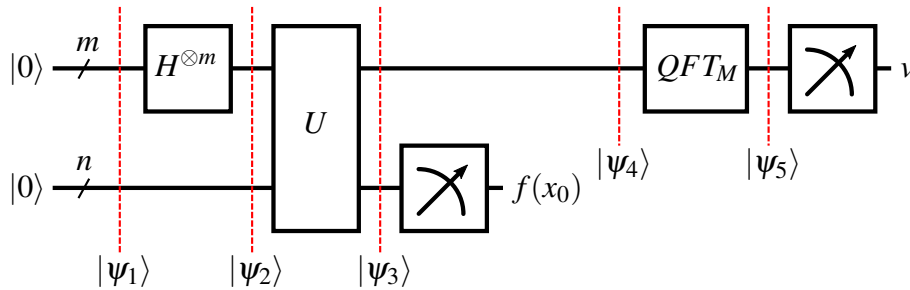


Figure 1: Circuit for period finding

Let's take a look at each state  $|\psi\rangle$  one by one. The first initial state is

$$|\psi_1\rangle = |0\rangle |0\rangle$$

Generating a superposition on the first register,

$$|\psi_2\rangle = \frac{1}{\sqrt{M}} \sum_{x=0}^{M-1} |x\rangle |0\rangle$$

Next,  $U$  produces

$$|\psi_3\rangle = \frac{1}{\sqrt{M}} \sum_{x=0}^{M-1} |x\rangle |f(x)\rangle$$

We measure the second register, which results  $f(x_0)$ . From now, we consider only the state of first register. The first register is composed of states that produce  $f(x_0)$ . Assuming the number of states of that kind is  $\mu$ .

$$|\psi_4\rangle = \frac{1}{\sqrt{\mu}} \sum_{y=0}^{\mu-1} |x_0 + yr\rangle$$

where

$$\mu = \lfloor M/r \rfloor \text{ or } \lfloor M/r \rfloor + 1 \quad (2)$$

according to  $x_0, r, M$ . To elicit  $r$ , do the phase analysis on it by applying  $QFT$ .

$$\begin{aligned} |\psi_5\rangle &= \frac{1}{\sqrt{\mu}} \sum_{y=0}^{\mu-1} QFT_M |x_0 + yr\rangle \\ &= \frac{1}{\sqrt{\mu}} \sum_{y=0}^{\mu-1} \frac{1}{\sqrt{M}} \sum_{x=0}^{M-1} e^{i2\pi x(x_0 + yr)/M} |x\rangle \\ &= \frac{1}{\sqrt{M\mu}} \sum_{x=0}^{M-1} e^{i2\pi x x_0/M} \sum_{y=0}^{\mu-1} e^{i2\pi x y r/M} |x\rangle \end{aligned}$$

where  $i = \sqrt{-1}$ . Let  $p(v)$  be the probability of getting  $v$  after measuring the first register. Then,

$$p(v) = \frac{1}{M\mu} |c(v)|^2$$

where

$$c(v) = \sum_{y=0}^{\mu-1} e^{i2\pi v y r/M}$$

## 2.2 Characteristic of outcome

Let  $E(v)$  be the distance between  $v$  and the nearest  $\frac{M}{r}k$  where  $k \in \mathbb{Z}$ . Then,

$$\exists k \in \mathbb{Z} \quad E(v) = \left| v - \frac{M}{r}k \right| < 1 \quad (3)$$

with high probability.

*Proof.* We consider two cases,  $r|M$  and  $r \nmid M$ .

i)  $r|M$

The number of states that produces specific result  $f(x_0)$  is a constant.

$$\mu = \frac{M}{r}$$

Again, considering two cases of  $\frac{M}{r}|v$  and  $\frac{M}{r} \nmid v$ . If  $\frac{M}{r}|v$ ,

$$e^{i2\pi vr/M} = 1 \Rightarrow c(v) = \frac{M}{r} \Rightarrow p(v) = \frac{1}{r}$$

If  $\frac{M}{r} \nmid v$ ,

$$e^{i2\pi vr/M} \neq 1 \Rightarrow c(v) = \sum_{y=0}^{M/r-1} \left( e^{i2\pi vr/M} \right)^y = \frac{1 - e^{i2\pi v}}{1 - e^{i2\pi vr/M}} = 0 \Rightarrow p(v) = 0$$

$\therefore$  (3) is held.

ii)  $r \nmid M$

We prove (3) by showing that  $\forall v_0 E(v_0) \leq 1$ ,  $\forall v' E(v') > 1$ ,

$$\frac{p(v_0)}{p(v')} \geq 9 \quad (4)$$

To prove it, we obtain a lower bound of  $p(v_0)$  and an upper bound of  $p(v')$ .

Lower bound first. If  $v = 0$ ,  $E(v) = 0$ . Getting  $p(v)$ ,

$$p(0) = \frac{1}{M\mu} \mu^2 = \frac{\mu}{M} \quad (5)$$

Keeping it for later, now assume  $v \neq 0$ . Since  $r \nmid M$ , or  $r \nmid 2^m$ ,  $r$  has at least one prime factor other than 2. Hence,

$$M \nmid vr \Rightarrow e^{i2\pi vr/M} \neq 1$$

This follows

$$c(v) = \sum_{y=0}^{\mu-1} \left( e^{i2\pi vr/M} \right)^y = \frac{1 - e^{i2\pi v\mu r/M}}{1 - e^{i2\pi vr/M}}$$

Let  $\alpha \stackrel{\text{def}}{=} 2\pi vr/M$ . Then,

$$\begin{aligned} |c(v)|^2 &= \left| \frac{1 - e^{i\alpha\mu}}{1 - e^{i\alpha}} \right|^2 = \left| \frac{1 - \cos \alpha\mu - i \sin \alpha\mu}{1 - \cos \alpha - i \sin \alpha} \right|^2 \\ &= \frac{(1 - \cos \alpha\mu)^2 + \sin^2 \alpha\mu}{(1 - \cos \alpha)^2 + \sin^2 \alpha} = \frac{2 - 2\cos \alpha\mu}{2 - 2\cos \alpha} = \frac{\sin^2 \frac{\alpha\mu}{2}}{\sin^2 \frac{\alpha}{2}} \end{aligned}$$

Replacing  $\alpha$  back,

$$|c(v)|^2 = \frac{\sin^2(\pi v\mu r/M)}{\sin^2(\pi vr/M)} \quad (6)$$

Let

$$\delta \stackrel{\text{def}}{=} v - \frac{M}{r}k \quad \text{for some } k \in \mathbb{Z} \quad (7)$$

Then,

$$\pi v \frac{r}{M} = \pi \left( \frac{M}{r}k + \delta \right) \frac{r}{M} = \pi k + \pi \delta \frac{r}{M}$$

Applying it to (6),

$$|c(v)|^2 = \frac{\sin^2(\pi \delta \mu r / M)}{\sin^2(\pi \delta r / M)} \stackrel{\text{def}}{=} h(\delta) \quad (8)$$

Given  $k$ , if  $-\frac{1}{2} \leq \delta < \frac{1}{2}$ , then there is exactly one  $v$  that satisfies (7). So

$$v_0 \stackrel{\text{def}}{=} \frac{M}{r}k + \delta \quad \left( -\frac{1}{2} \leq \delta < \frac{1}{2} \right)$$

Note that

$$|\sin x| \geq \left| \frac{x}{\pi/2} \right| \quad \text{for } x \in \left[ -\frac{\pi}{2}, \frac{\pi}{2} \right]$$

This follows

$$\sin^2 x \geq \frac{4x^2}{\pi^2} \quad (9)$$

Also, since  $\pi \delta r / M \ll 1$  by (1),

$$\sin \frac{\pi \delta r}{M} \approx \frac{\pi \delta r}{M} \quad (10)$$

Using (9) and (10),

$$\begin{aligned} h(\delta) &\geq \frac{\frac{4}{\pi^2} (\pi \delta \mu r / M)^2}{(\pi \delta r / M)^2} = \frac{4\mu^2}{\pi^2} \\ \therefore p(v_0) &= \frac{1}{M\mu} |c(v_0)|^2 = \frac{1}{M\mu} h(\delta) \geq \frac{4\mu}{\pi^2 M} \end{aligned} \quad (11)$$

This bound is persistent even if we remember (5). Note that this lower bound covers  $\forall v_0 \ E(v_0) \leq 1$  because  $p(v) \geq 0$  and  $-\frac{1}{2} \leq \delta < \frac{1}{2}$  implies  $|\delta| \leq 1$  and  $E(v_0) \leq 1$ .

Now look at the case of  $1 < |\delta| \leq \frac{M}{2r}$ . We'll get an upper bound of  $f(\delta)$ . To begin with,

$$g(x) \stackrel{\text{def}}{=} \frac{\sin x \mu \beta}{\sin x \beta} \quad \text{for } 1 \leq x \leq \frac{M}{2r}$$

where  $\beta = \pi r / M$ . We'll obtain the maximum value of  $g(x)$  and utilize it to get an upper bound of the case  $\forall v' \ E(v') > 1$ . To achieve it, get the values at critical points and compare them with  $g(1)$  and  $g(\frac{M}{2r})$ . Differentiating  $g(x)$ ,

$$g'(x) = \frac{\mu \beta \cos x \mu \beta \cdot \sin x \beta - \sin x \mu \beta \cdot \beta \cos x \beta}{\sin^2 x \beta}$$

Let  $x_0 \in (1, \frac{M}{2r})$  such that  $g'(x_0) = 0$ . Then,

$$\begin{aligned}\mu \cos x_0 \mu \beta \cdot \sin x_0 \beta &= \sin x_0 \mu \beta \cdot \cos x_0 \beta \\ \Rightarrow \mu \tan x_0 \beta &= \tan x_0 \mu \beta\end{aligned}$$

Putting back  $\beta$ ,

$$\mu \tan \frac{\pi x_0 r}{M} = \tan \frac{\pi x_0 \mu r}{M} \quad (12)$$

Meanwhile, from (2),

$$\mu \leq \frac{M}{r} \leq \mu + 2$$

Giving a boundary of  $r\mu/M$ ,

$$1 - \frac{2r}{M} < \frac{r\mu}{M} \leq 1 \quad (13)$$

From (1) and (13),

$$\mu r/M \approx 1 \quad (14)$$

From (12) and (14),

$$\mu \tan \frac{\pi x_0 r}{M} > \pi x_0$$

This states that the left hand side of (12) is large enough to approximate its solution as

$$\frac{\pi x_0 \pi r}{M} \approx \frac{\pi}{2} + \pi n \quad (15)$$

where  $n \in \mathbb{N} : x_0 > 1$  and (14). Now enumerating major values,

$$\begin{aligned}g(1) &= \frac{\sin(\pi \mu r/M)}{\sin(\pi r/M)} = \frac{\sin(\pi - \pi \mu r/M)}{\sin(\pi r/M)} \approx \frac{\pi - \pi \mu r/M}{\pi r/M} \\ &= \frac{M}{r} - \mu \leq \frac{M}{r} - \left\lfloor \frac{M}{r} \right\rfloor < 1\end{aligned}$$

$$g\left(\frac{M}{2r}\right) = \sin \frac{\pi}{2} \mu \leq 1$$

$$g(x_0) = \frac{1}{\sin\left[\frac{\pi}{\mu}\left(\frac{1}{2} + n\right)\right]} \leq \frac{1}{\sin\left[\frac{\pi}{\mu}\left(\frac{1}{2} + 1\right)\right]} = \frac{1}{\sin \frac{3\pi}{2\mu}} \approx \frac{2\mu}{3\pi}$$

Among the values,  $\frac{2\mu}{3\pi}$  is the largest one. Hence,

$$f(\delta) \leq \{g(\delta)\}^2 \leq \frac{4\mu^2}{9\pi^2} \quad \text{for } 1 < |\delta| \leq \frac{M}{2r}$$

If  $v' = \frac{M}{r}k + \delta$  such that  $1 < |\delta| \leq \frac{M}{2r}$ , then

$$p(v') = \frac{1}{M\mu} |c(v')|^2 = \frac{f(\delta)}{M\mu} \leq \frac{4\mu}{9\pi^2 M} \quad (16)$$

Combining two bounds, (11) and (16),

$$\frac{p(v_0)}{p(v')} \geq \frac{4\mu}{\pi^2 M} \frac{9\pi^2 M}{4\mu} = 9$$

Thus, when we measure  $|\psi_5\rangle$ , the integer  $v$  is at least 9 times likely to satisfy  $E(v) \leq 1$  than  $E(v) > 1$ . □

Here is a sample case.

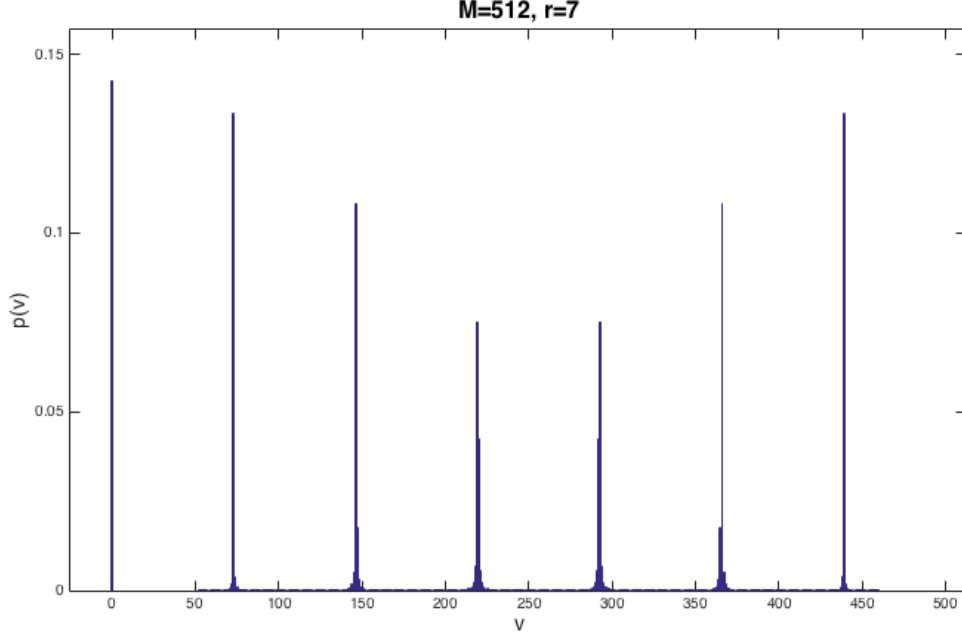


Figure 2: Sample distribution of  $p(v)$

### 2.3 Drawing out the period

Since it is highly probable  $E(v) \leq 1$ , from now, we assume that  $v$  an integer such that  $E(v) \leq 1$ . It follows

$$\left| v - \frac{M}{r}k \right| \leq 1 \quad \Rightarrow \quad \left| \frac{v}{M} - \frac{k}{r} \right| \leq \frac{1}{M} \leq \frac{1}{2r^2}$$

Using continued fractions technique, Theorem A4.16 at [NC00],  $\frac{k}{r}$  appears among convergents of  $\frac{v}{M}$ , giving  $r$  unless  $\gcd\{k, r\} \neq 1$ . A couple of ways are introduced to overcome the case  $\gcd\{k, r\} \neq 1$  in p.229 at [NC00]. Among them, one simple way is to repeat this procedure until  $k$  is a prime so that  $\gcd\{k, r\} = 1$ .

Problem 4.1 on p.638 at [NC00] states that there are at least  $\frac{r/2}{\lg r}$  primes in  $\{1, \dots, r\}$ . Therefore, considering that  $k$ 's are uniformly distributed over  $\{1, \dots, r\}$ , the probability that  $k$  is a prime is at least  $\frac{1}{2\lg r}$ . Furthermore, we know

$$\frac{1}{2\lg r} > \frac{1}{2\lg N}$$

Let  $p \stackrel{\text{def}}{=} \frac{1}{2\lg N}$  and assume  $p$  is sufficiently small (this happens frequently since this algorithm is usually applied when  $N$  is large). The probability of obtaining a prime  $k$  within  $s$  tries is

$$1 - (1 - p)^s \approx 1 - (1 - sp) = sp$$

Hence,  $s = 2\lg N$  tries pretty ensure that we earn prime  $k$  and  $r$ .

## 2.4 Summary

In sum,

1. Run the circuit.
2. Obtain an approximation of  $\frac{M}{r}k$ .
3. Elicit  $k'$  and  $r'$  using continued fractions technique.
4. Check if  $r' = r$  by checking  $f(0) = f(r')$ .
5. If  $r' = r$ , done. Otherwise, go to 1.

The repetition will end within  $2\lg N$  with high probability.

## References

- [GHH08] Andrew Wiles G. H. Hardy, Edward M. Wright. *An Introduction to the Theory of Numbers*. Oxford University Press, 6th edition, 2008.
- [NC00] Michael A Nielson and Isaac L Chuang. *Quantum computation and quantum information*. Cambridge University Press, 2000.
- [Ros03] Burton Rosenberg. Quantum fourier transforms. 2003.