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P18-0030

*Rough working*

|   |   |   |   |   |   |   |   |
|---|---|---|---|---|---|---|---|
|   | 1 | 2 | 3 | 4 | 5 | 6 | 7 |
| A | 3 | 9 | 5 | 5 | 6 | 3 | 9 |

|   |   |   |   |   |   |   |   |   |   |
|---|---|---|---|---|---|---|---|---|---|
|   | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 |
| C | 0 | 0 | 2 | 0 | 2 | 1 | 0 | 0 | 2 |

|   |   |   |   |   |   |   |   |   |   |
|---|---|---|---|---|---|---|---|---|---|
|   | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 |
| C | 0 | 0 | 2 | 0 | 2 | 1 | 0 | 0 | 2 |

|   |   |   |   |   |   |   |   |
|---|---|---|---|---|---|---|---|
|   | 1 | 2 | 3 | 4 | 5 | 6 | 7 |
| B | 3 | 3 | 5 | 5 | 6 | 6 | 9 |

Counting Sort is optimal than Quick Sort as in Counting Sort only simple iterations are involved.

for  $i = 1$  to  $k$  }  $k$  times  
 $C[i] = 0$

for  $j = 1$  to  $\text{len}(A)$

for  $j = 1$  to  $\text{len}(A)$  }  $n$ -elements  
do put count of numbers in  $C$  time

for  $f = 2$  to  $k$  }  $k-1$  times  
do accumulate  $C$

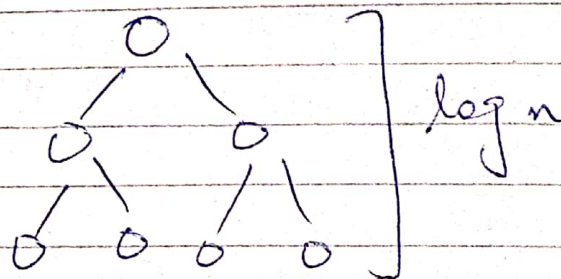
for  $m = \text{len}(A)$  to  $1$  (reverse) }  $n$  times  
do put values in  $B$

$$\begin{aligned}
 &= k + n + (k-1) + n \\
 &= k + n + k - 1 + n \\
 &= 2k + 2n - 1 \\
 &= O(k+n)
 \end{aligned}$$

if  $k = n$  then  $O(n)$ .

Whereas in Quick Sort we do partition and BST is formed at backend in best case.

Best case in Quick Sort:



Partition algo takes  $O(n)$  running time

So total running time =  ~~$O(n)$~~   $O(n \log n)$

In worst case:

(edge case) where pivot is at index 0 or index  $n-1$

$$\begin{aligned}
 &n + (n-1) + \cancel{n-2} + (n-2) + \dots + 1 \\
 &= \frac{n(n-1)}{2} \\
 &= O(n^2)
 \end{aligned}$$

So Counting Sort is better as its running time is more efficient and less than Quick Sort.