## Monima Atif Don

	Ti: Inserts "N"
	T: luserls 'a'
	Precingtion - cuz not Thread safe
	T2: Inserts G?
	Premption
(E)	T, : Inserts (z'
(7)	
	Premption
	T2: huserts (u)
(io)	Ti : Inserts (2)
(1)	Tz; luserto (a)

Best suited: Semaphores

2 Semaphores is best suited as it doesn't have

spin tock problem so the cycles are not wasted.

It's implemented in hardware and instructions

are atomic which is another plus point.

Q3 My second choice would be Hardware

Test and Set instructions became They are

atomic. So nutual exclusion will be taken care of.

And They are in hardware so They would be

fest.

Q4 Because semaphores don't have the problem of 'busy-waiting'. The Thread that has to wait is sent to blocked slate until it can be given trun. Whereas 'busy-waiting' so it wastes cycles of CPU. are running on same core so schedular has to give tuens to both of them concurrently. For example, schedular gives tuen to Tisead 1 and Their offices there puts Thread I in waiting state and schedules Thread 2. If The premiption does not occur or lets say Thread I is being executed on one core and Threed 2 is being executed on a separate core Then my amsevers to the questions above will be incorrect.