

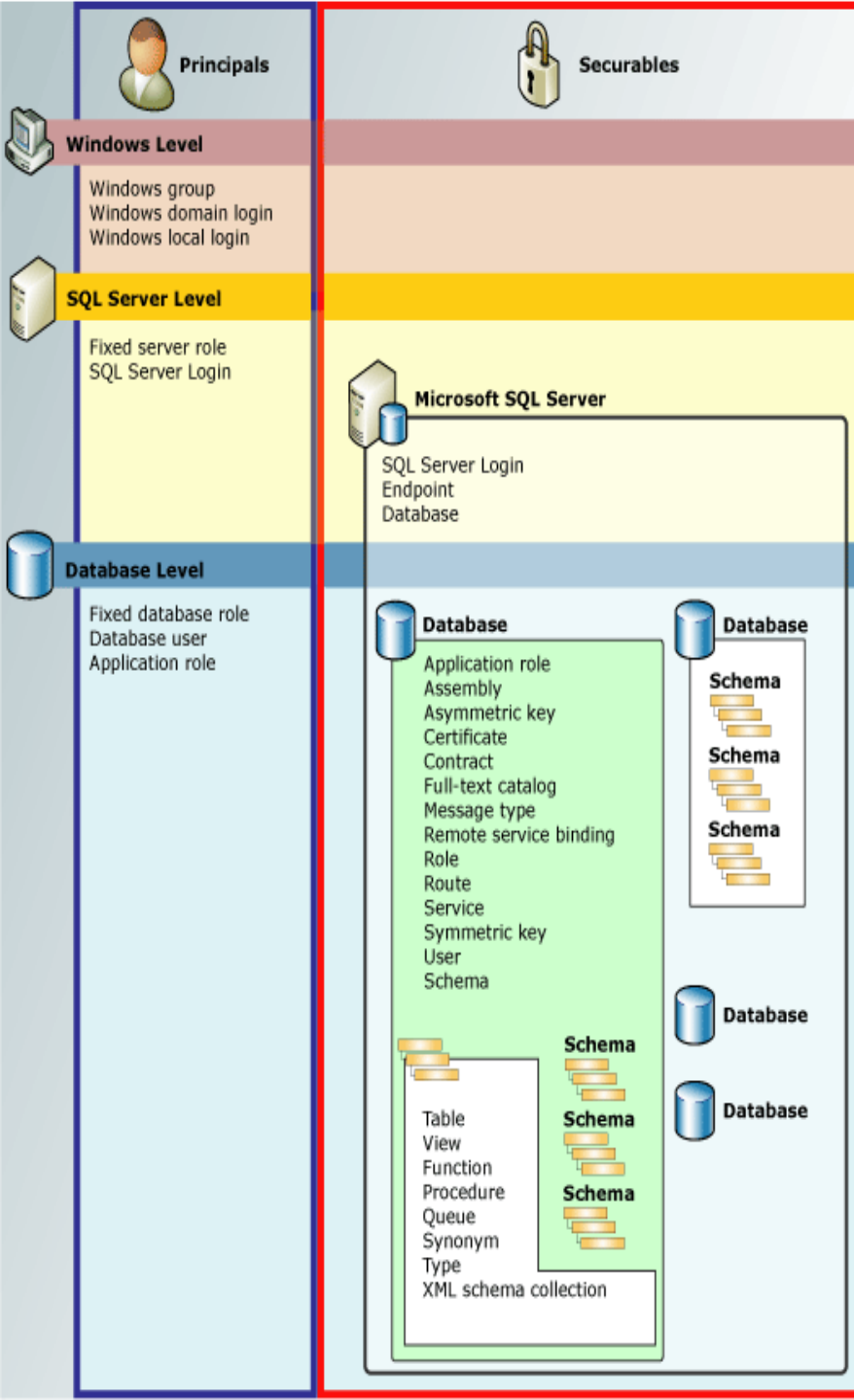
Chapter 2

RELATIONAL DATABASE



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OUTLINE

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Part 1

Relational Model Concepts



Relational Model Concepts

- Represent **data as a collection of relations**
- **Table** of values
 - Each **row** (*tuple*)
 - Represents a **record of related data values**
 - Facts that typically correspond to a real-world entity or relationship
 - Each **column** (*attribute*)
 - Holds a corresponding value for each row
 - Slot for a specific interpretation for a row

Relational Model Concepts

- Relational database: a set of relations.
- **Relation: consists of two parts:-**
 1. **Schema describes table**
 - Table name, attribute names and types
 - Specifies the **relation's name**, as well as the **name** and type of each **column**.
E.g. **Students(sid: string, name: string, login: string, age: integer, gpa: real)**
 2. **The current contents of the table are indicated by instance.**
 - Table of columns and rows
 - #rows = cardinality**
 - #fields = degree**

Example

The diagram illustrates the components of a database table. At the top, 'Relation Name' points to 'STUDENT'. 'Attributes' points to the column headers: 'Name', 'Ssn', 'Home_phone', 'Address', 'Office_phone', 'Age', and 'Gpa'. 'Tuples' points to the rows of data in the table.

| Name | Ssn | Home_phone | Address | Office_phone | Age | Gpa |
|----------------|-------------|---------------|----------------------|---------------|-----|------|
| Benjamin Bayer | 305-61-2435 | (817)373-1616 | 2918 Bluebonnet Lane | NULL | 19 | 3.21 |
| Chung-cha Kim | 381-62-1245 | (817)375-4409 | 125 Kirby Road | NULL | 18 | 2.89 |
| Dick Davidson | 422-11-2320 | NULL | 3452 Elgin Road | (817)749-1253 | 25 | 3.53 |
| Rohan Panchal | 489-22-1100 | (817)376-9821 | 265 Lark Lane | (817)749-6492 | 28 | 3.93 |
| Barbara Benson | 533-69-1238 | (817)839-8461 | 7384 Fontana Lane | NULL | 19 | 3.25 |

#rows = cardinality = 5 (Number of Tuples)

#fields = degree = 7 (Number of Attributes)

Part 2

Basic Structure



Characteristics Of Relations

- A Relation initially was a mathematical concept based on the ideas of **sets**.

product_name = {keyboard, mouse, pendrive, CD, ...} /* Set of all product names */

product_origin = {KL, Perak, Kedah, ...} /* Set of all product's origin*/

- The relation above also have **keys**, such as:
 - Primary Key, Foreign Key

CONSTRAINTS

- **Constraints are conditions** that must **hold on all valid relation states**.
- There are three main types of (**explicit schema based**) constraints that can be expressed in the relational model:

Key
constraints

Entity
integrity
constraints

Referential
integrity
constraints

- Another schema based constraint is the **domain constraint**.
 - Any value in a tuple must be from the attribute's domain (or NULL if that option is available).

Key Constraints

- Constraints are the **rules to follow** when entering data into database table columns.
- Constraints ensure that the user's data in columns **meets the condition's requirement**.

NOT NULL:

- ensures that the specified *column doesn't contain a NULL value*.

UNIQUE :

- *provides a unique/distinct values* to specified columns.

DEFAULT:

- *provides a default value to a column* if none is specified.

CHECK :

- *checks for the predefined conditions before inserting* the data inside the table.

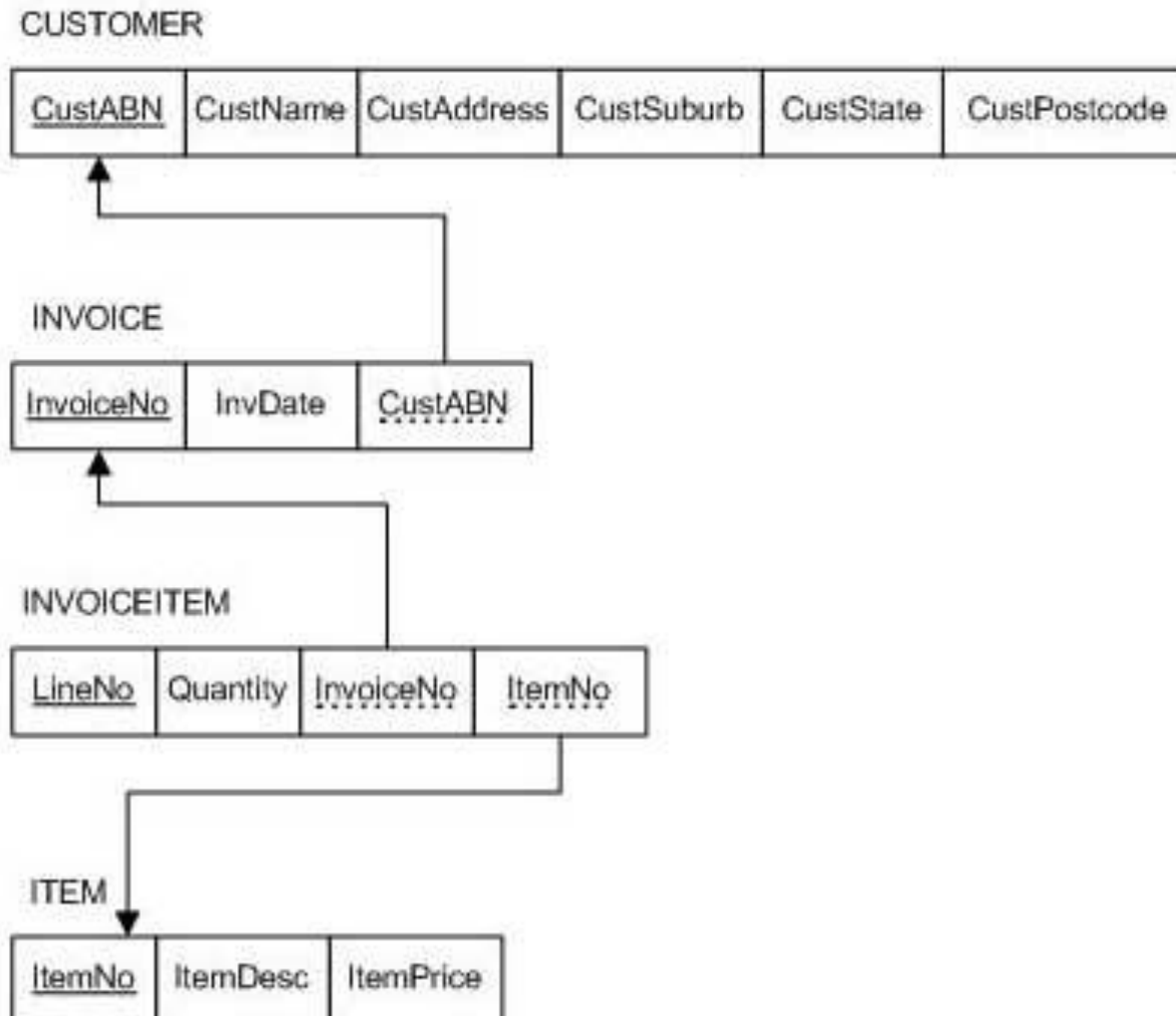
PRIMARY KEY:

- it *uniquely identifies a row* in a table.

FOREIGN KEY:

- ensures *referential integrity* of the relationship

Relational Database Schema



Update Operation in Relation

- **INSERT** a tuple.
- **DELETE** a tuple.
- **MODIFY** a tuple.

Class Exercise

Consider the following relations for a database that keeps track of student enrollment in courses and the books adopted for each course:

STUDENT (SID, Name, Major, Bdate)

COURSE (Course#, Cname, Dept, TID)

ENROLL (Enroll#, Course#, Quarter, Grade, SID, Semester)

TEACHER (TID, Qualification, Position)

Draw a **relational database schema** diagram specifying the foreign keys for this schema. Do underline the primary key.

**Part 3
(Week 3)**

Relational Algebra

Query Language

- A Language which is used to **store and retrieve data** from database is known as query language.
 - For example – **SQL**

Procedural Query language



- ☐ **User give direction** to the system to perform a series of operations to produce the desired output.
- ☐ **User tells what data to be retrieved** from database and **how to** retrieve it.

Non-procedural query language



- ☐ **User tell the system** to come up with output **without telling how to** do it.
- ☐ Users **tells what data** to be retrieved from database but **doesn't tell how to** retrieve it.

Relational Algebra & Relational Calculus



- **Relational Algebra:**

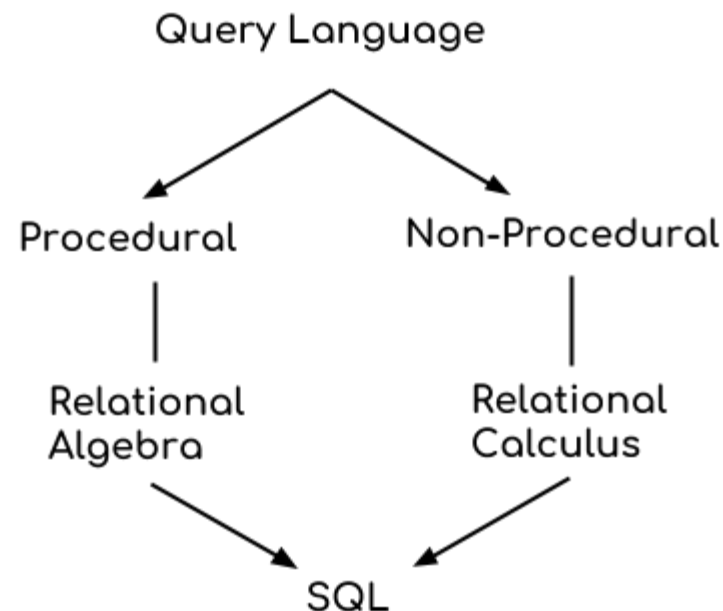
- **Relational algebra** is a conceptual **procedural query language** used on relational model.

- **Relational Calculus:**

- **Relational calculus** is a conceptual **non-procedural query language** used on relational model.

Difference between Relational Algebra, Calculus, RDBMS & SQL:

- **Relational algebra** and **calculus** are the **theoretical concepts** used on relational model.
- **RDBMS** is a **practical implementation of relational model**.
- **SQL** is a **practical implementation of relational algebra and calculus**.



Relational Algebra in Database

- Its based on the **procedural query language**.
 - to retrieve data from database or perform various operations such as **insert, update, delete** on the data.
- There are TWO categories of procedural language operation :

Basics

- Select (σ)
- Project (π)
- Union (\cup)
- Set Difference ($-$)
- Cartesian product (\times)
- Rename (ρ)

Derived

- Natural Join (\bowtie)
- Left, Right, Full outer join (\ltimes , \rtimes , $\ltimes\rtimes$)
- Intersection (\cap)
- Division (\div)



Basics

- Select (σ)
- Project (Π)
- Union (\cup)
- Set Difference ($-$)
- Cartesian product (\times)
- Rename (ρ)

Select (σ)

Represented by sigma (σ) and it is used to **find the rows** in a **table** which satisfy the **given condition** (**WHERE** statement in SQL).

σ Condition (Table name)

STUDENT

| Stu_Id | Stu_Name | City | Stu_Age |
|--------|----------|--------------|---------|
| 101 | Hamzah | Kuala Lumpur | 21 |
| 102 | Ajeet | Selangor | 19 |
| 103 | Liyau | Shenzen | 21 |
| 104 | Jason | Shenzen | 20 |

σ City= "Shenzen" (STUDENT)



| Stu_Id | Stu_Name | City | Stu_Age |
|--------|----------|---------|---------|
| 103 | Liyau | Shenzen | 21 |
| 104 | Jason | Shenzen | 20 |

Select (σ)

Display the information about the Student from Shenzen and age his/her above 20.

σ City = "Shenzen" AND Stu_Age > 20 (STUDENT)

| Stu_Id | Stu_Name | City | Stu_Age |
|--------|----------|--------------|---------|
| 101 | Hamzah | Kuala Lumpur | 21 |
| 102 | Ajeet | Selangor | 19 |
| 103 | Liyau | Shenzen | 21 |
| 104 | Jason | Shenzen | 20 |



| Stu_Id | Stu_Name | City | Stu_Age |
|--------|----------|---------|---------|
| 103 | Liyau | Shenzen | 21 |

Project (Π)

- Used to **select specific columns** (or attributes) from a table (or relation).
- Similar to **SELECT** statement in SQL.

Π column_name1, column_name2,
(table_name)

Lets say we want to **select only two column** form Student table

Π City, Stu_Age (STUDENT)



STUDENT

| Stu_Id | Stu_Name | City | Stu_Age |
|--------|----------|--------------|---------|
| 101 | Hamzah | Kuala Lumpur | 21 |
| 102 | Ajeet | Selangor | 19 |
| 103 | Liyau | Shenzen | 21 |
| 104 | Jason | Shenzen | 20 |

| City | Stu_Age |
|--------------|---------|
| Kuala Lumpur | 21 |
| Selangor | 19 |
| Shenzen | 21 |
| Shenzen | 20 |

Union (U)

Used to **select all the rows** (tuples) **from two tables** (relations).

```
table_name1 U table_name2
```

STUDENT

| Stu_Id | Stu_Name | City | Stu_Age |
|--------|----------|--------------|---------|
| 101 | Hamzah | Kuala Lumpur | 21 |
| 102 | Ajeet | Selangor | 19 |
| 103 | Liyau | Shenzen | 21 |
| 104 | Jason | Shenzen | 20 |

COURSE

| Stu_Id | Stu_Name |
|--------|----------|
| 104 | Jason |
| 105 | Maria |
| 106 | Xuan |

Π Stu_Name (STUDENT)
 \cup Π Stu_Name (COURSE)



| Stu_Name |
|----------|
| Hamzah |
| Ajeet |
| Liyau |
| Jason |
| Maria |
| Xuan |

Intersection (\cap)

- Used to **select common rows** (tuples) **from two tables** (relations).

`table_name1 \cap table_name2`

STUDENT

| Stu_Id | Stu_Name | City | Stu_Age |
|--------|----------|--------------|---------|
| 101 | Hamzah | Kuala Lumpur | 21 |
| 102 | Ajeet | Selangor | 19 |
| 103 | Liyau | Shenzen | 21 |
| 104 | Jason | Shenzen | 20 |

COURSE

| Stu_Id | Stu_Name |
|--------|----------|
| 104 | Jason |
| 105 | Maria |
| 106 | Xuan |

| Stu_Name |
|----------|
| Jason |



Π Stu_Name (STUDENT) \cap Π Stu_Name (COURSE)

Difference -

Used to **select all tuples(rows)** that are **present in Relation R1 but not present in Relation R2**

`table_name1 - table_name2`

STUDENT

| Stu_Id | Stu_Name | City | Stu_Age |
|--------|----------|--------------|---------|
| 101 | Hamzah | Kuala Lumpur | 21 |
| 102 | Ajeet | Selangor | 19 |
| 103 | Liyau | Shenzen | 21 |
| 104 | Jason | Shenzen | 20 |

COURSE

| Stu_Id | Stu_Name |
|--------|----------|
| 104 | Jason |
| 105 | Maria |
| 106 | Xuan |

Π Stu_Name (STUDENT) - Π Stu_Name (COURSE)



| Stu_Name |
|----------|
| Hamzah |
| Ajeet |
| Liyau |

Cartesian Product (X)

If we have two relations R1 and R2 then the Cartesian product of these two relations ($R1 \times R2$) would **combine each tuple of first relation R1** with the **each tuple of second relation R2**.

R1 X R2

STUDENT : B

| Stu_Id | Stu_Name |
|--------|----------|
| 101 | Hamzah |
| 102 | Ajeet |

COURSE : C

| Course_Id | Course_Name |
|-----------|-------------|
| CST403 | Database |
| CST104 | Intro |

B X C



| Stu_Id | Stu_Name | Course_Id | Course_Name |
|--------|----------|-----------|-------------|
| 101 | Hamzah | CST403 | Database |
| 101 | Hamzah | CST104 | Intro |
| 102 | Ajeet | CST403 | Database |
| 102 | Ajeet | CST104 | Intro |

Matrix 2 x 2

Rename (ρ)

Used to change or **rename the table** (relation name) or **attributes** of the table.

$\rho(\text{new_relation_name}, \text{old_relation_name})$

STUDENT

| Stu_Id | Stu_Name | City | Stu_Age |
|--------|----------|--------------|---------|
| 101 | Hamzah | Kuala Lumpur | 21 |
| 102 | Ajeet | Selangor | 19 |
| 103 | Liyau | Shenzen | 21 |
| 104 | Jason | Shenzen | 20 |

Rename the attributes (Stu_Age) from table above to "Age"

$\rho(\text{Age}, \Pi(\text{Stu_Age})(\text{STUDENT}))$

| Age |
|-----|
| 21 |
| 19 |
| 21 |
| 20 |

Links to study



<https://www.youtube.com/watch?v=KaIRmVD-v3U> → Basic operator of relational algebra

<https://www.youtube.com/watch?v=xZELQc11ACY> → Intersection and Union Set Concepts

<https://www.mathsisfun.com/sets/venn-diagrams.html> → Venn Diagram explanation