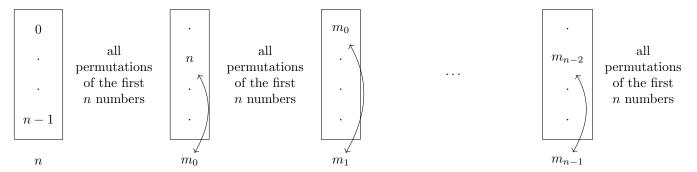
NOTES ON CRYPTARITHM SOLVER AND PERMUTATIONS

ERIC MARTIN

1. Heap's algorithm

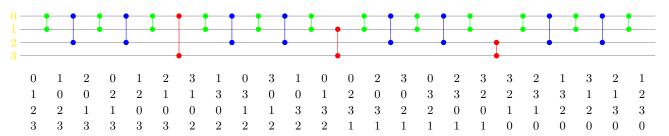
Let a nonzero natural number n be given. Heap's algorithm generates all permutations of a set S with n+1 elements, in such a way that any permutation, the first one excepted, is obtained from the previous one by exchanging two of S's elements. Without loss of generality, take for S the set $\{0, 1, \ldots, n\}$. The recursive version of Heap's algorithm can be illustrated as follows.



So the algorithm generates all permutations of the form $L \star n$, then all permutations of the form $L \star m_0$, then all permutations of the form $L \star m_1, \ldots$, and eventually all permutations of the form $L \star m_{n-1}$. The scheme is correct if $\{m_0, m_1, \ldots, m_{n-1}\} = \{0, \ldots, n-1\}$: at every stage, the algorithm has to select a new number from the first n ones (and exchange it with the current (n+1)st number). Heap's algorithm uses the following strategy:

- In case n is odd, select the first number, then the second number, then the third number...
- In case n is even, always select the first number.

The following diagram illustrates with n = 3.



Note that starting with (0,1,2,3), Heap's algorithm produces (1,2,3,0) as last permutation. We will see that starting with (0,1,2,3,4,5), it would produce (3,4,1,2,5,0) as last permutation; starting with (0,1,2,3,4,5,6,7), it would produce (5,6,1,2,3,4,7,0) as last permutation. More generally, starting with $(0,1,2,\ldots,2n+1)$, Heap's algorithm will produce as last permutation $(2n-1,2n,1,2,\ldots,2n-2,2n+1,0)$.

Note that starting with (0,1,2), Heap's algorithm produces (2,1,0) as last permutation. We will see that starting with (0,1,2,3,4), it would produce (4,1,2,3,0) as last permutation; starting with (0,1,2,3,4,5,6), it would produce (6,1,2,3,4,5,0) as last permutation. More generally, starting with $(0,1,2,\ldots,2n)$, Heap's algorithm will produce as last permutation $(2n,1,2,\ldots,2n-1,0)$.

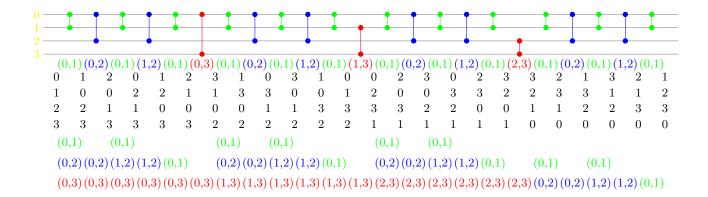
The previous formulas can be used to generalise Heap's algorithm and generate all sequences of k numbers chosen from $\{0,1,\ldots,n\}$: when the last number has been selected, and the penultimate number has been selected, ..., and the

2 ERIC MARTIN

(n-k+1)st number has been selected, it suffices to stop the recursion and "simulate" all permutations of the remaining n+1-k numbers by applying those formulas.

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The following is a possible recursive implementation of Heap's algorithm.
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The following is a possible iterative implementation of Heap's algorithm.



2. Proof of correctness

We prove that Heap's algorithm is correct and that moreover, the following holds for all $n \geq 1$:

(1) starting with $(0,1,2,\ldots,2n)$, all permutations of $(0,1,2,\ldots,2n)$ are generated, ending in

$$(2n, 1, 2, \ldots, 2n - 1, 0)$$

(2) starting with $(0,1,2,\ldots,2n+1)$, all permutations of $(0,1,2,\ldots,2n+1)$ are generated, ending in

$$(2n-1,2n,1,2,\ldots,2n-2,2n+1,0)$$

Proof is by induction. The base case n = 1 is straightforward, so let $n \ge 1$ be given, and assume that (1) holds. We show that (2) holds too.

- Starting from $(0, 1, 2, ..., 2n 1, 2n) \star 2n + 1$, Heap's algorithm generates all permutations of the form $L \star 2n + 1$, ending with $(2n, 1, 2, ..., 2n 1, 0) \star 2n + 1$.
- Permuting first and last elements, $(2n, 1, 2, ..., 2n 1, 0) \star 2n + 1$ is changed to $(2n + 1, 1, 2, ..., 2n 1, 0) \star 2n$. Starting from $(2n + 1, 1, 2, ..., 2n - 1, 0) \star 2n$, the algorithm then generates all permutations of the form $L \star 2n$, ending with $(0, 1, 2, ..., 2n - 1, 2n + 1) \star 2n$.
- Permuting second and last elements, $(0, 1, 2, ..., 2n 1, 2n + 1) \star 2n$ is changed to $(0, 2n, 2, ..., 2n 1, 2n + 1) \star 1$. Starting from $(0, 2n, 2, ..., 2n - 1, 2n + 1) \star 1$, the algorithm then generates all permutations of the form $L \star 1$, ending with $(2n + 1, 2n, 2, ..., 2n - 1, 0) \star 1$.
- Permuting third and last elements, $(2n+1,2n,2,\ldots,2n-1,0) \star 1$ is changed to $(2n+1,2n,1,3\ldots,2n-1,0) \star 2$. Starting from $(2n+1,2n,1,3\ldots,2n-1,0) \star 2$, the algorithm then generates all permutations of the form $L \star 2$, ending with $(0,2n,1,3\ldots,2n-1,2n+1) \star 2$.
- Permuting fourth and last elements, $(0, 2n, 1, 3..., 2n-1, 2n+1) \star 2$ is changed to $(2n+1, 2n, 1, 2..., 2n-1, 0) \star 3$. Starting from $(2n+1, 2n, 1, 2..., 2n-1, 0) \star 3$, the algorithm then generates all permutations of the form $L \star 3...$ till all permutations of the form $L \star 2n-1$, ending in $(2n+1, 2n, 1, 2, ..., 2n-2, 0) \star 2n-1$.
- Permuting last two elements, $(2n+1,2n,1,2,\ldots,2n-2,0)\star 2n-1$ is changed to $(2n+1,2n,1,2,\ldots,2n-2,2n-1)\star 0$. Starting from $(2n+1,2n,1,2,\ldots,2n-2,2n-1)\star 0$, the algorithm then generates all permutations of the form $L\star 0$, ending with $(2n-1,2n,1,2,\ldots,2n-2,2n+1)\star 0$.

So we have established that (2) holds.

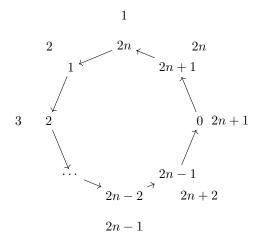
Now assume that (2) holds. We show that (1) with n replaced by n+1 holds too. The inner circle of the following diagram shows how elements move from one position to another one after all permutations of a list consisting of the 2n+2 numbers $0, \ldots, 2n+1$ have been performed by Heap's algorithm. For instance, the first element ends up as the last element, moving from position (index) 0 and eventually ending up at position (index) 2n+1. After all permutations of $(0,1,2,\ldots,2n,2n+1)$ have been generated, ending in $(2n-1,2n,1,2,\ldots,2n-2,2n+1,0)$, so after all permutations of $(0,1,2,\ldots,2n,2n+1) \star 2n+2$ have been generated, ending in $(2n-1,2n,1,2,\ldots,2n-2,2n+1,0) \star 2n+2$, Heap's algorithm replaces the element now at position 0, that is, 2n-1 (originally at position 2n-1), with 2n+2. This is depicted in the following diagram with 2n+2 on the outer circle facing 2n-1 on the inner circle. At the end of each of the following stages, the algorithm permutes the element currently at position 0 with the element currently at position 2n+2, that is, the element at position 0 at the end of previous stage. Hence as illustrated in the diagram, move to position 2n+2; first 2n-1 replaced by 2n+2, then 2n-2 replaced by 2n-1, ..., then 2 replaced by 3, then 1 replaced by 2, then 2n replaced by 1, then 2n+1 replaced by 2n, and eventually 0 replaced by 2n+1. Finally, all permutations of the numbers then at position 0, ..., 2n+1 (those numbers being 1, 2, ..., 2n+2) are generated, corresponding to a last, (2n+2)nd rotation in the following diagram, hence a rotation following a "full circle". This means that:

- ends up at position 0 the element which at the beginning of this last round of permutations, is a position 2n-1, that is, 2n+2,
- ends up at position 1 the element which at the beginning of this last round of permutations, is a position 2n, that is, 1,

• ...

resulting in the final list $(2n+2,1,2,3,\ldots,2n-1,2n,2n+1) \star 0$. So we have established that (1) with n replaced by n+1 holds.

4 ERIC MARTIN



3. Notes on the implementation of the cryptarithm solver

The first version is a minor adaptation of code written by Raymond Hettinger as part of ActiveState Code Recipes. It uses the permutations() function, imported from itertools (and also the findall() function, imported from re). The second version does not import anything.

The first version filters out the permutations that assign 0 to one of the letters that start a word; the second version does not produce those permutations. The first version creates a string where the letters starting a word come first, followed by the letters not starting a word; it is the other way around for the second version. For instance, with the cryptarithm SEND + MORE == MONEY, the first version creates a string which could be SMENDORY, whereas the second version creates a string which could be ENDORYSM. Let us still use that example to explain how we proceed in the second version.

- Starting with the list L = (0, 1, 2, 3, 4, 5, 6, 7, 8, 9), we use the generalisation of Heap's algorithm to generate lists of the form L_1L_2 , one for each possible list L_2 of two nonzero digits (so we ignore 0, as if L started with the element of index 1). This determines a possible assignment to SM.
- For each list of the form L_1L_2 generated as described, we make a copy of L_1L_2 and we use again the generalisation of Heap's algorithm to generate from the copy lists of the form $L_{11}L_{12}L_2$, one for each possible list L_{12} of six digits, amongst those in L_1 (so 0 is now allowed but we ignore L_2 , as if we were working with a list that ended at index 7). We return the last 8 digits of $L_{11}L_{12}L_2$, that is, $L_{12}L_2$, allocating the digits in L_{12} to ENDORY and the digits in L_2 to SM.

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