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## Homework

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**Exercise: More on hash functions** Let  $p$  be a prime number, let  $\mathbb{Z}_p$  denote the field of arithmetic mod  $p$ . For each  $a, b, c \in \mathbb{Z}_p$ , let  $P_{a,b,c} : \mathbb{Z}_p \rightarrow \mathbb{Z}_p$  be the quadratic polynomial defined by  $P_{a,b,c}(x) = ax^2 + bx + c \bmod p$ . Show that the family of functions  $\{P_{a,b,c} : a, b, c \in \mathbb{Z}_p\}$  satisfies the following 3-universal property: for every three *distinct*  $x_1, x_2, x_3 \in \mathbb{Z}_p$  and every  $y_1, y_2, y_3 \in \mathbb{Z}_p$  we have

$$\Pr_{a,b,c \in_r \mathbb{Z}_p} [P_{a,b,c}(x_i) = y_i \text{ for } i = 1, 2, 3] = \frac{1}{p^3}.$$

The notation “ $a, b, c \in_r \mathbb{Z}_p$ ” under “Pr” means that  $a, b$  and  $c$  are chosen uniformly and independently at random in  $\mathbb{Z}_p$ .

**Exercise: More on hashing for estimating sizes of sets** Let  $m$  and  $k$  be positive integers and let  $U = U_{m,k}$  be a 2-universal family of hash functions from  $m$  bits to  $k$  bits. For any fixed set  $S \subseteq \{0, 1\}^m$  and a randomly chosen  $h \in U$ , let  $I(S, h)$  be the indicator random variable for the event that  $h$  has collisions on  $S$ : “there exist two distinct  $x, y \in S$  with  $h(x) = h(y)$ ”. Prove the following:

1. If  $|S| > 2^k$ , then  $\Pr_{h \in_r U} [I(S, h) = 1] = 1$
2. If  $|S| \leq 2^{k/2}$ , then  $\Pr_{h \in_r U} [I(S, h) = 1] \leq 1/2$ .

Recall that the notation “ $h \in_r U$ ” under “Pr” means that  $h$  is chosen uniformly at random in  $U$ .

**Exercise: Approximate counting up to a square** Use the result of the previous exercise to show that for any #P function  $f$  there exists a *deterministic* polynomial time algorithm that, given an  $n$ -bit input  $x$  and access to an  $\text{NP}^{\text{NP}}$ -oracle, outputs a number  $t$  satisfying  $t \leq f(x) \leq t^2$ .