

# Analysis of Loops

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We have already discussed *Asymptotic Analysis*, *Worst, Average and Best Cases* and *Asymptotic Notations*. In this , analysis of iterative programs with simple examples is discussed.

1.

**O(1):** Time complexity of a function (or set of statements) is considered as O(1) if it doesn't contain loop, recursion and call to any other non-constant time function.

```
// set of non-recursive and non-loop statements
```

For example swap( ) function has O(1) time complexity. A loop or recursion that runs a constant number of times is also considered as O(1). For example the following loop is O(1).

```
// Here c is a constant
for (int i = 1; i <= c; i++) {
    // some O(1) expressions
}
```

2.

**O(n):** Time Complexity of a loop is considered as O(n) if the loop variables is incremented / decremented by a constant amount. For example following functions have O(n) time complexity.

```
// Here c is a positive integer constant
for (int i = 1; i <= n; i += c) {
    // some O(1) expressions
}
```

```
for (int i = n; i > 0; i -= c) {
    // some O(1) expressions
}
```

3.

**O(n^c):** Time complexity of nested loops is equal to the number of times the innermost statement is executed. For example the following sample loops have O(n<sup>2</sup>) time complexity

```

for (int i = 1; i <=n; i += c) {
    for (int j = 1; j <=n; j += c) {
        // some O(1) expressions
    }
}

```

```

for (int i = n; i > 0; i -= c) {
    for (int j = i+1; j <=n; j += c) {
        // some O(1) expressions
    }
}

```

4.

**O(Logn):** Time Complexity of a loop is considered as  $O(\text{Log}n)$  if the loop variables is divided / multiplied by a constant amount.

```

for (int i = 1; i <=n; i *= c) {
    // some O(1) expressions
}

```

```

for (int i = n; i > 0; i /= c) {
    // some O(1) expressions
}

```

For example Binary Search(refer iterative implementation) has  $O(\text{Log}n)$  time complexity. Let us see mathematically how it is  $O(\text{Log} n)$ . The series that we get in first loop is 1, c,  $c^2$ ,  $c^3$ , ...  $c^k$ . If we put k equals to  $\text{Log}c^n$ , we get  $c^{(\text{Log}c^n)}$  which is n.

5.

**O(LogLogn):** Time Complexity of a loop is considered as  $O(\text{LogLog}n)$  if the loop variables is reduced / increased exponentially by a constant amount.

```

// Here c is a constant greater than 1
for (int i = 2; i <=n; i = pow(i, c)) {
    // some O(1) expressions
}

```

```

// Here fun() is function to find square root
// or cuberoot or any other constant root
for (int i = n; i > 1; i = fun(i)) {
    // some O(1) expressions
}

```

**How to combine time complexities of consecutive loops?** When there are consecutive loops, we calculate time complexity as sum of time complexities of individual loops.

```
for (int i = 1; i <= m; i += c) {  
    // some O(1) expressions  
}
```

```
for (int i = 1; i <= n; i += c) {  
    // some O(1) expressions  
}
```

```
Time complexity of above code is  $O(m) + O(n)$  which is  $O(m+n)$   
If  $m == n$ , the time complexity becomes  $O(2n)$  which is  $O(n)$ .
```

**How to calculate time complexity when there are many if, else statements inside loops?**

As discussed earlier, the worst-case time complexity is the most useful among best, average and worst. Therefore we need to consider the worst case. We evaluate the situation when values in if-else conditions cause a maximum number of statements to be executed. For example, consider the linear search function where we considered the case when an element is present at the end or not present at all. When the code is too complex to consider all if-else cases, we can get an upper bound by ignoring if else and other complex control statements.