# **Detecting Infected Hosts and Domains**

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https://github.com/moneydance/591project

#### **Abstract**

Advanced persistent threats have become a major concern for IT professionals around the world. Their stealthy distributed nature makes them difficult to identify and remove. However because the connections between backdoors and command and control centers leave telltale traces in DNS data, certain graph theory techniques can be used to identify malicious domains, and infected hosts.

# 1 Introduction

Cyber security is an ever-changing field. As malware becomes more advanced so do methods of detection. Recently a sophisticated attack called an Advanced Persistent Threat (APT) has emerged. This stealthy form of malware attempts to blend in to the normal operations of the target organization. Fortunately due to the communication and infection vectors employed by APTs, tell tale patterns are left in an organization's DNS logs. In our implementation we leverage these communication patterns to detect suspicious activity that potentially indicates an APT infection.

### 2 APTs

# 2.1 Infection

APT's enter an organizations network through various means. USB's carrying malicious code, emails, and compromised websites are the most common vectors.

# 2.2 Covertly Spread

After infection the APT attempts to move through the network by infecting additional hosts. It does this covertly taking advantage of unpatched vulnerabilities and hijacked credentials.

#### 2.3 Exfiltrate Data

After collecting the correct credentials and moving to the target hosts. The APT will begin to silently pass sensitive data out of the organization's network.

### 2.4 Call Home

In order to spread through the network and steal information the backdoor established by the APT must communicate and take commands from a hacker outside of the organization's network. Because organizations block inbound traffic, the communication must be initialized within the organization. To do this http/https connections are made from backdoors to command and control servers (C&C). These malicious domains are contacted at certain automated time intervals, asking for additional instructions from the server. Because these command and control domains are only contacted by infected machines their traffic is low. By finding rarely frequented domains and looking for patterns in the time intervals between a host contacting them we can label potentially malicious domains.

# 3 Methodology

# Algorithm

- 1. Parse the logs into a graph. We build a bipartite graph where nodes consist of two sets, domains and hosts, edges represent a connection between a domain and host.
- 2. Run degree centrality on the graph to find domains with a low number of connections.
- 3. From these rare domains look for suspicious behavior. If the domain seems to be interacting with hosts in a suspicious manner mark it as a potential C&C domain.
- 4. Use these potential C&C domains and hosts connecting to these domains as a seed for our belief propagation algorithm.

### 3.1 Rare Domains

Malicious domains are contacted by a small subset of infected hosts. Because of this the degree of these nodes tends to be low. We can take advantage of this by using a degree centrality algorithm to discover rarely contacted domains. With this list of rarely contacted domains we can then look at the properties of the edges between the domains and the hosts that contacted them. We use two properties to define suspicious edges. The first is the time intervals of an individual host contacting the domain. If the standard deviation of the time intervals is low this indicates scheduled behavior (i.e a host was contacting the server for instructions on a regular basis). The second property is the number of CNAMES used by the DNS server. CNAMES map one domain name to another. A malicious domain could use multiple CNAMES to obfuscate the intended destination domain and make the original destination appear more legitimate. Using these two properties we are able to distinguish between legitimate low traffic websites and potential C&C domains.

### 3.2 Belief Propagation

We can use a belief propagation algorithm on a seed of malicious hosts and domains to iteratively grow our subgraph of infected hosts and know C&C domains. We do this by first looking at all the rare domains contacted by our infected seed hosts. If these domains exhibit malicious C&C like behavior we add them to a list of new malicious domains. If we have not detected any new malicious domains we begin belief propagation. We assign rare domains a score based on their properties being similar to known C&C domains. If this score is above a certain threshold we add the domain to our list of new malicious domains. After these two steps we update our known malicious domains with the newly found ones. We then update our malicious hosts with the hosts that contacted these new malicious domains, and update the rare domains we were looking at with the rare domains contacted by these new malicious hosts. We continue to do this until a certain number of C&C domains and malicious hosts are detected.

# 4 Results

Due to the sheer size of the data itself, our experiments were focused on just a single day's worth of DNS logs, from which we could potentially iterate on and refine our methods.

This could potentially form a basis for creating a more in-depth, focused, and localized area of interest within the possible infected sub-network(s) of the entire network, which could be more useful than a broad analysis involving an entire month of data.

We run our code with the following settings, and output the sets of domains and hosts we find, as well as edge information between them, to a log file:

```
infile = '2013-03-17'
num_edges = 500000
degelist = parseToGraph.parse(infile, num_edges=num_edges)
G = construct_graph(edgelist)
hosts, doms = belief_propagation(G, set(), set(), threshold=0.7)
```

```
74.92.39.47
74.92.32.18
```

<sup>74.92.74.110</sup> 

```
74.92.47.73
```

- 74.92.67.20
- 74.92.174.204
- 74.92.74.157
- 74.92.38.152
- 74.92.210.24
- 74.92.169.178
- 74.92.14.160
- 74.92.111.62
- 74.92.62.205
- 74.92.241.121
- 74.92.36.107
- 74.92.56.80
- 74.92.39.83
- 74.92.69.169
- 74.92.10.60
- 74.92.169.56
- 74.92.42.46
- 74.92.155.178
- 184.202.111.41
- 74.92.96.151
- 74.92.148.15
- 184.202.20.220
- 74.92.163.47
- 74.92.171.191
- 74.92.39.79
- 74.92.30.112
- 74.92.220.19
- 74.92.30.116
- 74.92.140.160
- 74.92.136.59
- 74.92.39.53
- 92.160.212.105
- 74.92.23.86
- 74.92.185.4
- 74.92.4.27
- 74.92.107.121
- 74.92.245.101
- 252.90.80.26
- 74.92.180.150
- 74.92.8.144
- 74.92.243.44
- 74.92.36.115
- 74.92.74.11
- 74.92.46.106
- 74.92.118.27
- 74.92.94.183
- 74.92.123.9
- 74.92.138.187
- 74.92.111.222
- 74.92.176.102
- 58.229.128.1
- 74.92.77.104 74.92.169.10
- 74.00.14.07
- 74.92.14.27
- 74.92.208.220
- 74.92.231.233 74.92.172.7
- 74.92.100.219

```
74.92.226.78
```

74.92.169.110

74.92.4.74

74.92.215.80

74.92.80.56

74.92.179.46

74.92.147.238

74.92.114.49

74.92.208.178

74.92.50.119

74.92.175.32

58.208.125.7

58.208.125.6

74.92.224.26

74.92.12.8

58.229.45.32

74.92.190.162

74.92.125.38

74.92.240.231

74.92.49.12

74.92.54.53

74.92.26.44

74.92.80.130

74.92.100.40

74.92.74.94

74.92.4.214

74.92.185.33

74.92.50.52

74.92.182.142

184.202.159.108

74.92.43.108

74.92.157.35

74.92.38.3

74.92.139.55

74.92.94.190

184.202.84.131

74.92.77.153

74.92.50.31

74.92.77.115

74.92.132.75

14.92.132.13

184.202.138.101

74.92.248.83

74.92.83.155 74.92.250.98

74.92.240.78

74.92.255.26

74.92.151.144

74.92.12.52

74.92.89.91

74.92.65.85

74.92.81.93

74.92.25.140

74.92.81.90

74.92.64.45

184.202.152.7

74.92.38.238 74.92.79.62

74.92.195.204

74.92.150.213

74.92.125.134 74.92.156.31 184.202.58.166 74.92.213.113 ump.thumb.dimly.wad fulfil.johannes.wad hastening.nullify.wad suites.dusted.wad cot.ledger.wad rattiest.add.wad lam.preponderances.wad peddler.vet.wad requisition.vanishing.wad fa.fop.plot.hated.wad gluey.jeans.tad.wad shrimp.ab-z7g6r.noe fa.ad-.plot.hated.wad u.refurnished.wad cot.fireproof.wad oberon.aacire9v9zf.wad pestilence.jocasta.noe ob.enhanced.wad step.hated.wad braved.racier.wad rev.quileless.wad pit.clashed.rimbaud.wad blantyre.superstitions.wad bacteriologists.sapped.console.val i.refurnished.wad pours.comb.co.rd dick.ably.ox.wad ump.frill.dimly.wad pert.la.agt.slyly.wad sn.hated.wad cot.preponderances.wad na.blantyre.superstitions.wad cot.ki.wad ming.foam.inching.hated.wad blantyre.rev.console.noe occlusion.rimbaud.wad rev.faint.wad ripple.nails.wad stoppering.conversationalists.wad aaaa8y5807h1ayfufc0u7.tamra.b.tridents.noe cot.jeans.tad.wad rho.sedation.relentless.wad kempis.jeans.tad.wad pm.ohio.wad cot.sledge.wad fa.rue.plot.hated.wad gent.ti.ty.friend.noe did.toothy.co.rd shadowiest.tad.wad

# An example printout of the edge data encoded between a host and domain in this list is:

For nodes 74.92.32.18, braved.racier.wad: date: 2013-03-17 00:05:38.806021

```
data:
? braved.racier.wad A
! braved.racier.wad CNAME collective.racier.wad
! collective.racier.wad CNAME collective.racier.wad.eco.racier.wad.friend.noe
! collective.racier.wad.eco.racier.wad.friend.noe CNAME marathon.racier.wad.wryness.noe
! marathon.racier.wad.wryness.noe CNAME moody.g.eyeballing.noe
! moody.g.eyeballing.noe A 59.195.249.218
date: 2013-03-17 00:10:38.720933
data:
? braved.racier.wad A
! braved.racier.wad CNAME collective.racier.wad
! collective.racier.wad CNAME collective.racier.wad.eco.racier.wad.friend.noe
! collective.racier.wad.eco.racier.wad.friend.noe CNAME marathon.racier.wad.wryness.noe
! marathon.racier.wad.wryness.noe CNAME moody.g.eyeballing.noe
! moody.g.eyeballing.noe A 59.195.249.218
date: 2013-03-17 00:15:38.834693
data:
? braved.racier.wad A
! braved.racier.wad CNAME collective.racier.wad
! collective.racier.wad CNAME collective.racier.wad.eco.racier.wad.friend.noe
! collective.racier.wad.eco.racier.wad.friend.noe CNAME marathon.racier.wad.wryness.noe
! marathon.racier.wad.wryness.noe CNAME moody.g.eyeballing.noe
! moody.g.eyeballing.noe A 59.195.249.218
date: 2013-03-17 00:20:38.755377
data:
? braved.racier.wad A
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! collective.racier.wad CNAME collective.racier.wad.eco.racier.wad.friend.noe
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date: 2013-03-17 00:25:38.757728
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! marathon.racier.wad.wryness.noe CNAME moody.g.eyeballing.noe
! moody.g.eyeballing.noe A 59.195.249.218
date: 2013-03-17 00:30:39.008724
data:
? braved.racier.wad A
! braved.racier.wad CNAME collective.racier.wad
! collective.racier.wad.CNAME collective.racier.wad.eco.racier.wad.friend.noe
! collective.racier.wad.eco.racier.wad.friend.noe CNAME marathon.racier.wad.wryness.noe
! marathon.racier.wad.wryness.noe CNAME moody.g.eyeballing.noe
! moody.g.eyeballing.noe A 59.195.249.218
```

As you can see, the interactions between this host and domain occurred extremely regularly (At 5-minute intervals) and always involved multiple CNAME responses. Our Belief Propagation algorithm was also able to find a sizable list of domains and hosts that we can further inspect, and observe over a longer time period.

# 5 Conclusion

In conclusion our methodology seems to be correct as a proof of concept. However for real world applications our system must be made more robust by being able to handle months of data instead of just a day. To do this we would have to build a software stack involving databases to filter out the normal DNS traffic of a company and find rarely contacted domains. Our methods for finding C&C domains could also be improved by expanding the parameters we are using to define suspicious activity and using more sophisticated metrics and methods than, for example, just the standard deviation of communication intervals between hosts and domains to identify automated domains.

### References

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- [3] Alina Oprea, Zhou Li, Ting-Fang Yen, Sang H. Chin and Sumayah Alrwais. Detection of Early-Stage Enterprise Infection by Mining Large-Scale Log Data