

# CS 350: Assignment #5

Due on Thursday, March 17, 2016

*Bestavros 3:00 pm*

**Ben Pusey**

1.  $\lambda = 100 \text{ rps}$

$T_s = 0 - 19 \text{ ms}$  uniformly distributed

bound on F-Score unfairness

$$\frac{1}{2} \lambda T_g = \frac{1}{2} g \text{ by Little's Law}$$

To solve for  $g$  we must use M/G/1 formula

$$g = \frac{\rho^2 A}{1 - \rho} + \rho \quad A = \frac{1}{2} \left[ 1 + \left( \frac{\sigma T_s}{T_s} \right)^2 \right]$$

$$\rho = \frac{\lambda}{\mu} \quad \mu = \frac{1}{T_s}$$

$$T_s = \frac{0 + 19}{2} = 9.5 \text{ ms} = .0095 \text{ s}$$

$$\sigma T_s = \sqrt{\frac{1}{12} (b - a)^2} = \sqrt{\frac{1}{12} (19)^2} = 5.5 \text{ ms} = .0055 \text{ s}$$

$$A = \frac{1}{2} \left[ 1 + \left( \frac{.0055}{.0095} \right)^2 \right] = .67$$

$$\mu = \frac{1}{.0095} = 105$$

$$\rho = \frac{100}{105} = .95$$

$$g = \frac{.95^2 (.67)}{1 - .95} + .95 = 12.96$$

So the bound is  $(\frac{1}{2}) 12.96$  or  $6.5$

**Problem 2:****FCFS**

Low Lambda

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CLOCK: 100004.58204220934
Rho: 0.0098266704691827
IO System Time 0.012825090401281155
CPU System Time 0.3120715908599175
IO Tasks through system: 6029
CPU Tasks through system: 2995
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Normal Lambda

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CLOCK: 100000.00494249418
Rho: 0.9611812499997576
IO System Time 7.156759764158426
CPU System Time 7.436791653536298
IO Tasks through system: 600654
CPU Tasks through system: 300588
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$$\text{Slowdown I/O} = \frac{7.16}{.013} = 550$$

$$\text{Slowdown CPU} = \frac{7.44}{.31} = 24$$

The slowdown for I/O is much longer than the slowdown for cpu because I/O jobs have to wait for the much longer cpu jobs ahead of them, and their optimal system times are very low.

**SRTN**Low Lambdas

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CLOCK: 100022.04952596084  
Rho: 0.009745373282928895  
IO System Time 0.010058497904441464  
CPU System Time 0.30702980961107906  
IO Tasks through system: 6025  
CPU Tasks through system: 2993

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Normal Lambda

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CLOCK: 100000.23194325212  
Rho: 0.9589099005456027  
IO System Time 0.014010580748338622  
CPU System Time 1.44634390924814  
IO Tasks through system: 600748  
CPU Tasks through system: 298971

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$$\text{Slowdown I/O} = \frac{.014}{.01} = 1.4$$

$$\text{Slowdown CPU} = \frac{1.45}{.3} = 4.8$$

The slowdown for I/O is lower than the slowdown for CPU because I/O tasks have priority over CPU tasks since their service time is lower.

**RR**Low Lambdas

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CLOCK: 100020.82290120877  
Rho: 0.009837938942820677  
IO System Time 0.01049754071171709  
CPU System Time 0.30351451323802425  
IO Tasks through system: 5953  
CPU Tasks through system: 3070

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Normal Lambda

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CLOCK: 100000.02893491843  
Rho: 0.9615470237624546  
IO System Time 1.9791537542270279  
CPU System Time 7.383351984222936  
IO Tasks through system: 600051  
CPU Tasks through system: 300386

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$$\text{Slowdown I/O} = \frac{1.98}{.01} = 198$$

$$\text{Slowdown CPU} = \frac{7.4}{.3} = 24.66$$

The lower slowdown of I/O in comparison to the I/Os Fifo slowdown makes sense as the time quantum is .1 so a I/O task only has to wait .1 seconds for CPU tasks ahead of them as opposed to their full service time.

**HSN**Low Lambdas

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CLOCK: 100002.867397461  
Rho: 0.009363340133960237  
IO System Time 0.010516938701498349  
CPU System Time 0.2962037126131248  
IO Tasks through system: 6114  
CPU Tasks through system: 2971

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Normal Lambda

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CLOCK: 100000.07831014277  
Rho: 0.96003786142339  
IO System Time 0.1670700196483842  
CPU System Time 3.9945701548944794  
IO Tasks through system: 600314  
CPU Tasks through system: 300119

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Slowdown I/O =  $\frac{.17}{.01} = 17$

Slowdown CPU =  $\frac{4}{.3} = 13.33$

The slowdown for I/O and CPU are similar because HSN attempts to minimize the slowdown of the tasks in the system by servicing tasks with higher slowdowns.

**IO Slowdown Rankings**

1. SRTN slowdown = 1.4
2. HSN slowdown = 17
3. RR slowdown = 198
4. FCFS slowdown = 550



3.

a. The most obvious example would be if Q1 always has requests in it then requests in Q2 will never be serviced.

b. n - batch scan

c. A larger N makes the cache more efficient

d. A larger N is less fair though.

e.

□  $N=100$  The load on the system is low so we don't care about fairness. Setting N to 100 will also increase our systems performance.

□  $N=10$  The load is moderate so we should be more concerned about fairness.  $N=10$  provides a good balance on efficiency and fairness.

□  $N=1$  high load so fairness is a concern, caching provides little benefit here because there is little locality.

□  $N=10$  because the load is extremely high we should be concerned with fairness, however extremely high locality means caching will be effective, 10 provides a good balance between the two.



4.

a. yes starvation is possible. Class C jobs are the most susceptible because class B jobs are scheduled more frequently than class C jobs and have a higher priority.

b.

A event	$1/\text{minute}$	highest priority	5 seconds
B event	$1/\text{second}$	medium priority	.5 seconds
C event	$1/5 \text{ minutes}$	lowest priority	20 seconds

A worst case 5 seconds. it has the highest priority

B worst case 5.5 seconds. an A event arrives during the B events execution.

C worst case 50 seconds. I used geometric series to think about this

event C takes 20 seconds to do work which will generate 20 B events. Those 20 B events take 10s to service so they will generate an additional 10 B events those 10 B events take 5s to service so they will generate 5 more B events. This is a geometric series so the 20 seconds of work on A create 20 seconds of work on B events. If an A event arrives its 5 seconds of work will generate 5 seconds of work on B events. Total is  $20 + 20 + 5 + 5 = 50$ .



$$C \cdot 2X + \left( \left\lceil \frac{2X}{60} \right\rceil 5 \right) 2 = 300$$

$X = 125$  So 125 is the maximum amount of cpu time for C.

- d. if a Job C arrives just before a Job A the Job A will have to wait for Job C to finish using the resource R before it can begin to be serviced. Therefore the higher priority Job A must wait for the lower priority Job C.
- e. A Job C arrives before Job A. Then the C job is interrupted by Job B's causing Job C to take 40 seconds to complete Job A then runs for a total of 45 seconds.
- f. The above is an example of priority inversion even though Job A has higher priority than Job B. Job B's are serviced before Job A. A higher priority job is preempted by a lower priority one.
- g. Priority inheritance attempts to solve the problem of priority inversion. The priority of a low priority task is increased to the highest priority task waiting for the shared resource being consumed by the low priority task.
- h. Job B's can no longer interrupt Job C's using the resource so its worst case service time is 25 s