



Code as Art

Blog about system programming and not only

Say hello to x64 Assembly [part 2]

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Some days ago I wrote the first blog post - introduction to x64 assembly - [Say hello to x64 Assembly \[part 1\]](#) which to my surprise caused great interest:

Say hello to x64 Assembly, part 1 (0xax.blogspot.com)
158 points by nkurz 1 day ago | 40 comments

↑ 121 Say hello to Assembly [part 1] [Linux x84-64] (0xax.blogspot.com)
submitted 2 days ago by 0xAX

It motivates me even more to describe my way of learning. During this days I got many feedback from different people. There were many grateful words, but more important for me, there were many advices and adequate critics. Especially I want to say thank you words for great feedback to:

- [Fiennes](#)
- [Grienders](#)
- [nkurz](#)

And all who took a part in discussion at Reddit and Hacker News. There were many opinions, that first part was a not very clear for absolute beginner, that's why I will write more informative posts. So, let's start with second part of *Say hello to x64 assembly*.

Terminology and Concepts

As I wrote above, I got many feedback from different people that some parts of first [post](#) are not clear, that's why let's start from description of some terminology we will see in this and next parts.

Register - register is a small amount of storage inside processor. Main point of processor is data processing. Processor can get data from memory, but it is slow operation. That's why processor has own internal restricted set of data storage which name is - register.

Little-endian - we can imagine memory as one large array. It contains bytes. Each address stores one element of the memory "array". Each element is one byte. For example we have 4 bytes: *AA 56 AB FF*. In little-endian the least significant byte has the smallest address:

```
0 FF
1 AB
2 56
3 AA
```

where 0,1,2 and 3 are memory addresses.

Big-endian - big-endian stores bytes in opposite order than little-endian. So if we have *AA 56 AB FF* bytes sequence it will be:

```
0 AA
1 56
2 AB
3 FF
```

Syscall - is the way a user level program asks the operating system to do something for it. You can find syscall table - [here](#). **Stack** - processor has a very restricted set of registers. So stack is a continuous area of memory addressable special registers RSP,SS,RIP and etc... We will take a closer look on stack in next parts.

Section - every assembly program consists from sections. There are following sections:

- **data** - section is used for declaring initialized data or constants
- **bss** - section is used for declaring non initialized variables
- **text** - section is used for code

General-purpose registers - there are 16 general-purpose registers - rax, rbx, rcx, rdx, rbp, rsp, rsi, rdi, r8, r9, r10, r11, r12, r13, r14, r15. Of course, it is not only terms and concepts which related with assembly programming. If we will meet another strange and unfamiliar words in next blog posts, there will be explanations.

Data Types

The fundamental data types are bytes, words, doublewords, quadwords, and double quadwords. A byte is eight bits, a word is 2 bytes, a doubleword is 4 bytes, a quadword is 8 bytes and a double quadword is 16 bytes (128 bits).

Now we will work only with integer numbers, so let's see to it. There are two types of integer: unsigned and signed. Unsigned integers are unsigned binary numbers held in a byte, word, doubleword, and quadword. Their values range from 0 to 255 for an unsigned byte integer, from 0 to 65,535 for an unsigned word integer, from 0 to $2^{32} - 1$ for an unsigned doubleword integer, and from 0 to $2^{64} - 1$ for an unsigned quadword integer. Signed integers are signed binary numbers held in a byte, word and etc... The sign bit is set for negative integers and cleared for positive integers and zero. Integer values range from -128 to $+127$ for a byte integer, from $-32,768$ to $+32,767$ for a word integer, from -2^{31} to $+2^{31} - 1$ for a doubleword integer, and from -2^{63} to $+2^{63} - 1$ for a quadword integer.

Sections

As I wrote above, every assembly program consists of sections, it can be a data section, text section and bss section. Let's look at the *data* section. Its main purpose is to declare initialized constants. For example:

```
1 section .data
2     num1:    equ 100
3     num2:    equ 50
4     msg:     db "Sum is correct", 10
```

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Ok, it is almost all clear here. 3 constants with names num1, num2, msg and with values 100, 50 and "Sum is correct", 10. But what is it *db*, *equ*? Actual NASM number of pseudo-instructions:

- DB, DW, DD, DQ, DT, DO, DY and DZ - are used for declaring initialized data. For example:

```
1 ;; Initialize 4 bytes 1h, 2h, 3h, 4h
2 db 0x01,0x02,0x03,0x04
3
4 ;; Initialize word to 0x12 0x34
5 dw 0x1234
```

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- RESB, RESW, RESD, RESQ, REST, RESO, RESY and RESZ - are used for declaring non-initialized variables
- INCBIN - includes External Binary Files
- EQU - defines constant. For example:

```
1 ;; now one is 1
2 one equ 1
```

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- TIMES - Repeating Instructions or Data. (description will be in next posts)

Arithmetic operations

There is a short list of arithmetic instructions:

- ADD - integer add
- SUB - subtract
- MUL - unsigned multiply
- IMUL - signed multiply
- DIV - unsigned divide
- IDIV - signed divide
- INC - increment
- DEC - decrement
- NEG - negate

Some of them we will see at practice in this post. Others will be covered in next posts.

Control flow

Usually programming languages have ability to change order of evaluation (with if statement, case statement, goto and etc...) and assembly has it too. Here's some of it. There is *cmp* instruction for performing comparison between two values. It is used along with the conditional jump instruction for decision making example:

```
1 ;; compare rax with 50
2 cmp rax, 50
```

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cmp instruction just compares 2 values, but doesn't affect them and doesn't execute anything depend on result of comparison. For performing any actions after comparison there is *conditional jump* instructions. It can be one of it:

- JE - if equal
- JZ - if zero
- JNE - if not equal
- JNZ - if not zero
- JG - if first operand is greater than second
- JGE - if first operand is greater or equal to second
- JA - the same that JG, but performs unsigned comparison
- JAE - the same that JGE, but performs unsigned comparison

For example if we want to make something like if/else statement in C:

```
1 if (rax != 50) {
2     exit();
3 } else {
4     right();
5 }
```

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It will be in assembly:

```
1 ;; compare rax with 50
2 cmp rax, 50
3 ;; perform .exit if rax is not equal 50
4 jne .exit
5 jmp .right
```

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There is also unconditional jump with syntax:

JMP label

For example:

```
1 _start:
2     ;; ....
3     ;; do something and jump to .exit label
4     ;; ....
5     jmp .exit
6
7 .exit:
8     mov rax, 60
9     mov rdi, 0
10    syscall
```

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Here we have can have some code which will be after *_start* label, and all of this code will be executed, assembly transfer control to *.exit* label, and code after start to execute.

Often unconditional jump uses in loops. For example we have *label* and some code after it. This code executes anything, then we have condition and jump to this code if condition is not successfully. Loops will be covered in next parts.

Example

Let's see simple example. It will take two integer numbers, get sum of these numbers and compare it with predefined number. If predefined number is equal will print something on the screen, if not - just exit. Here is the source code of our example:

```
1  ;initialised data section
2  section .data
3      ; Define constants
4      num1:    equ 100
5      num2:    equ 50
6      ; initialize message
7      msg:     db "Sum is correct\n"
8
9  section .text
10
11      global _start
12
13      ;; entry point
14      _start:
15          ; set num1's value to rax
16          mov rax, num1
17          ; set num2's value to rbx
18          mov rbx, num2
19          ; get sum of rax and rbx, and store it's value in rax
20          add rax, rbx
21          ; compare rax and 150
22          cmp rax, 150
23          ; go to .exit label if rax and 150 are not equal
24          jne .exit
25          ; go to .rightSum label if rax and 150 are equal
26          jmp .rightSum
27
28      ; Print message that sum is correct
29      .rightSum:
30          ;; write syscall
31          mov     rax, 1
32          ;; file descriptor, standard output
33          mov     rdi, 1
34          ;; message address
35          mov     rsi, msg
36          ;; length of message
37          mov     rdx, 15
38          ;; call write syscall
39          syscall
40          ; exit from program
41          jmp .exit
42
43      ; exit procedure
44      .exit:
45          ; exit syscall
46          mov     rax, 60
47          ; exit code
48          mov     rdi, 0
49          ; call exit syscall
50          syscall
```

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Let's go through the source code. First of all there is *data* section with two constants *num1*, *num2* and variable *msg* with "Sum is correct\n" value. Now look There is begin of program's entry point. We transfer *num1* and *num2* values to general purpose registers *rax* and *rbx*. Sum it with *add* instruction. After execution instruction, it calculates sum of values from *rax* and *rbx* and store its value to *rax*. Now we have sum of *num1* and *num2* in the *rax* register.

Ok we have num1 which is 100 and num2 which is 50. Our sum must be 150. Let's check it with *cmp* instruction. After comparison *rax* and 150 we check result comparison, if *rax* and 150 are not equal (checking it with *jne*) we go to *.exit* label, if they are equal we go to *.rightSum* label.

Now we have two labels: *.exit* and *.rightSum*. First is just sets 60 to *rax*, it is exit system call number, and 0 to *rdi*, it is a exit code. Second is *.rightSum* is pre just prints *Sum is correct\n*. If you don't understand how it works, see first [post](#).

Conclusion

It was a second part of series 'say hello to x64 assembly', if you will have a questions/suggestions write me a comment.

All source code you can find as everytime - [here](#).