

Preliminary Quiz

STATE YOUR NAME AND UFVID and answer the following questions, please:

- 1. Define Algorithm.
- 2. Define Dijkstra's Control Structures.
- Define Von Neuman Architecture.
- 4. Describe the *Instruction Cycle*.
- Define the main characteristics of OOP.
- Define Recursion.
- 7. Write the <u>Java Code</u> for "Is there anybody out there?"
- 8. What does "a mod b" mean in words?

Goals of this Lecture

- Briefly summarize the introductory material (Section 1: Chapters 1.- 4 in several PPTs).
- Introduce the following fundamental notions:
 - The need for data structures within the context of computing,
 - The programming paradigms and to compare and contrast the various approaches to computing within these paradigms,
 - The core ideas behind data structures (which are nothing more than [presumably efficient] methods of information storage, manipulation and retrieval), and
 - The core ideas behind algorithms (which are nothing more than methods of addressing solvable problems where data structures are an integral part of algorithmic efficiency).

Goals of this Lecture

- **NOTE:** This is an **intermediate-level course** with emphasis on OO problem solving in **Java**.
- My approach is to briefly introduce data structures from the procedural viewpoint (i.e., a language like C), which is less abstract and, therefore, more amenable for developing the core concepts of data structures.
- Then, we increase the level of abstraction using OO techniques within the context of Java, which will form the main body of this course.

Preamble



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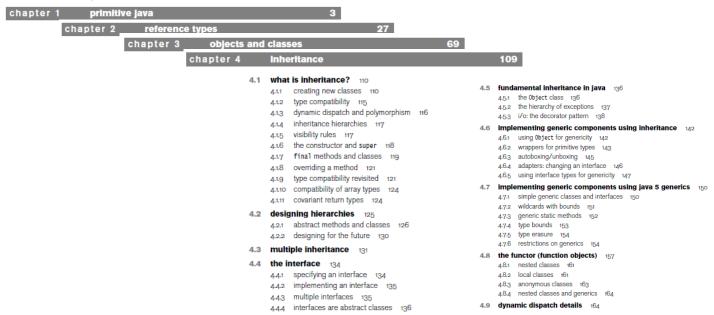


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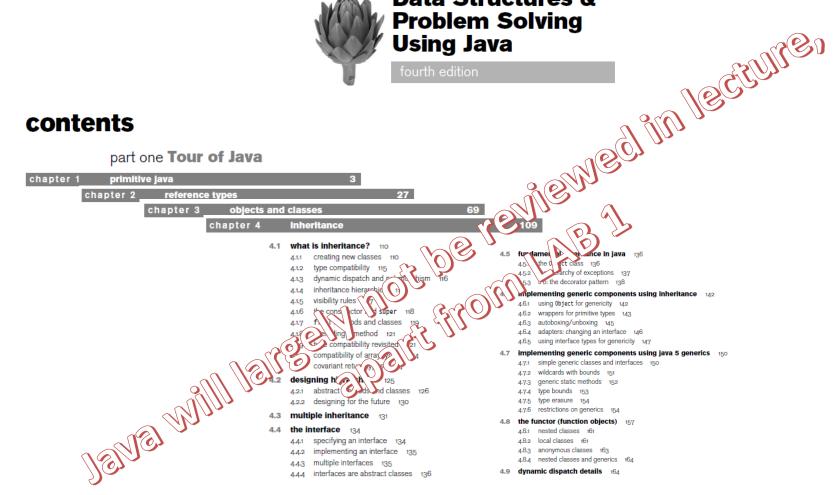


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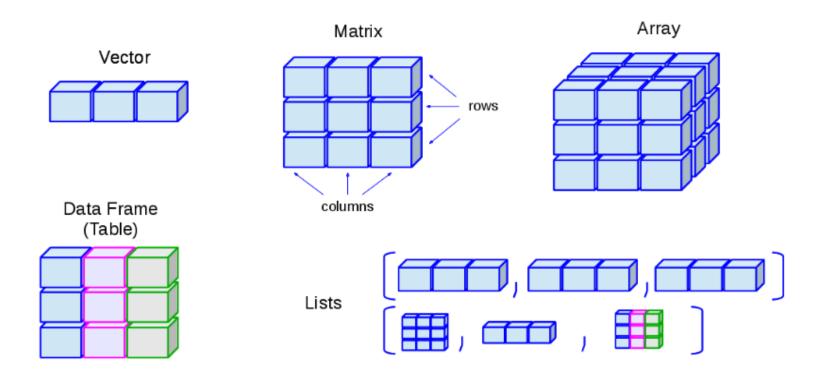




contents



Beyond Java: Introductory Notions of Data Structures and Algorithms



The study of *algorithms*, including

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The study of *algorithms*, including

- Their formal and mathematical properties,
- Their hardware realizations,
- Their linguistic realizations,
- Their applications.

Schneider and Gersting, 1999

Definition: Algorithm

Definition: Algorithm

A <u>well-ordered</u> collection of <u>unambiguous</u> and <u>effectively computable operations</u> that, when executed, <u>produces a result</u> and <u>halts</u> in a finite amount of time.

- Schneider and Gersting, 1999

Apply this definition to the quadratic equation an consider how to implement it with pseudocode.

Apply this definition to the quadratic equation an consider how to implement it with pseudocode.

 To be effectively computable, the discriminant must not be less than zero AND the leading coefficient must not be zero.

$$x=rac{-b\pm\sqrt{b^2-4ac}}{2a}.$$

Apply this definition to the quadratic equation an consider how to implement it with pseudocode.

- To be effectively computable, the discriminant must not be less than zero AND the leading coefficient must not be zero.
- Note that software development begins with defining the problem (requirements of SWE).

$$x=rac{-b\pm\sqrt{b^2-4ac}}{2a}$$
 .

So a generic BASIC-like pseudocode might be as below, where d is the discriminant.

```
1 let d = b * b - 4 * a * c
2 if d < 0 then
3 print "No real roots."
4 else if d = 0 then
5 print "Root is: "; -b / (2 * a); "."
6 else
7 print "Roots are: "; (-b + sqrt ( d )) / (2 * a); " and "; (-b - sqrt ( d )) / 8 end if</pre>
```



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```

Note that this is best done using Procedural Programming – OO is an unnecessary abstraction for most simple problems in mathematics.



Data Structures (Efficient Organization of Data)

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- Operating Systems (Managing the Computer)

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- Information Storage and Retrieval (DBMS)

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- Computational Complexity (Limitations of Computing)
- Architecture (Design of Computers)
- Intelligent Systems (AI)
- Automata Theory (Abstract Computers)
- Information Storage and Retrieval (DBMS)
- Software Engineering (generally large systems designed with professional engineering practices)

Post-Preamble Summary

- We have defined fundamental terminology:
 - Computing Science,
 - Algorithms,
- We have also provided an overview of the various areas of research in Computing Science and have seen where Data Structures fits into this taxonomy.

The Fundamental Motivation for Data Structures

The Philosophy of Data Structures

- The primary goal of this lecture is to motivate the <u>need</u> to study data structures.
- This involves understanding that unlimited computing power has logical barriers and that, hence, there is a fundamental requirement to understand the nature of how data may be manipulated by any computer and, at the very least, for the sake of efficient storage and retrieval of data.
- Computer programs are formal mathematical structures and, hence, these issues begin with the mathematical nature of algorithms.

Given a constant, n, and a variable, x, we can construct the following:

xⁿ

Monomial

Given a constant, n, and a variable, x, we can construct the following:

•
$$P(x) = a_n x^n + a_{n-1} x^{n-1} + ... + a_1 x + a_0$$

$$-P_2(x) =$$

$$-P_1(x)=$$

$$-P_0(x) =$$

$$a_1 x + ... + a_1 x + a_0$$

$$a_2 x^2 + a_1 x + a_0$$

$$a_1x + a_0$$

$$a_0$$

- Monomial
- Polynomial
- Quadratic
- Linear
- Constant

Given a constant, n, and a variable, x, we can construct the following:

$$P(x) = a_n x^n + A_n +$$

• $P(x) = a_n x^n + a_{n-1} x^{n-1} + ... + a_1 x + a_0$

$$a_2 x^2 + a_1 x + a_0$$

$$a_1x + a_0$$

$$a_0$$

- Monomial

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- Constant

n^x

Exponential

Given a constant, n, and a variable, x, we can construct the following:

 a_0

•
$$P(x) = a_n x^n + a_{n-1} x^{n-1} + ... + a_1 x + a_0$$

 $-P_2(x) = a_2 x^2 + a_1 x + a_0$
 $-P_1(x) = a_1 x + a_0$
 $-P_0(x) = a_0$

- *n*^x
- n!

- Monomial
- Polynomial
- Quadratic
- Linear
- Constant
- Exponential
- Factorial

Given a constant, n, and a variable, x, we can construct the following:

- xⁿ
- $P(x) = a_n x^n + a_{n-1} x^{n-1} + ... + a_1 x + a_0$

$$-P_2(x) =$$

$$-P_1(x)=$$

$$-P_{O}(x) =$$

$$a_2 x^2 + a_1 x + a_0$$

- $a_1 x + a_0$
 - a_0

- Monomial
- Polynomial
- Quadratic
- Linear
- Constant
- Exponential
- Factorial

- *n*^x
- n!
- And their inverse functions (such as $x^{1/n}$ or $\log x$)

Computational Time of Various Classes of Problems According to Input Size

TABLE 2 The Computer Time Used by Algorithms.							
Problem Size	Bit Operations Used						
п	$\log n$	n	$n \log n$	n^2			
10	$3 \times 10^{-11} \text{ s}$	10^{-10} s	$3 \times 10^{-10} \text{ s}$	10^{-9} s			
10^{2}	$7 \times 10^{-11} \text{ s}$	10^{-9} s	$7 \times 10^{-9} \text{ s}$	10^{-7} s			
10^{3}	$1.0 \times 10^{-10} \text{ s}$	10^{-8} s	$1 \times 10^{-7} \text{ s}$	10^{-5} s			
10^{4}	$1.3 \times 10^{-10} \text{ s}$	10^{-7} s	$1 \times 10^{-6} \text{ s}$	10^{-3} s			
10 ⁵	$1.7 \times 10^{-10} \text{ s}$	10^{-6} s	$2 \times 10^{-5} \text{ s}$	0.1 s			
10 ⁶	$2 \times 10^{-10} \text{ s}$	10^{-5} s	$2 \times 10^{-4} \text{ s}$	0.17 min			

Computational Time of Various Classes of Problems According to Input Size

Problem Size		Bit Operations Used				
n	log n	n	$n \log n$	n^2	2^n	n!
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10^{2}	$7 \times 10^{-11} \text{ s}$	10^{-9} s	$7 \times 10^{-9} \text{ s}$	10^{-7} s	$4 \times 10^{11} \text{ yr}$	*
10^{3}	$1.0 \times 10^{-10} \text{ s}$	10^{-8} s	$1 \times 10^{-7} \text{ s}$	10^{-5} s	*	*
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Informal Definition

Tractable Problem:

An algorithm that does not have a mathematical structure that is worse than a polynomial.

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An algorithm that does not have a mathematical structure that is worse than a polynomial.

Note: The idea here is that a tractable problem can be solved in *a reasonable time* with reasonable *storage* (*i.e.*, you can solve it in less than the age of the known universe).

Computational Complexity

- We will deal with this important subject in greater detail further on in this course.
- In particular, we will provide a formal definition of an intractable function which will require auxiliary definitions and an analysis of the growth of functions.
- Moreover, this will lead to a discussion of the logical limitations of computation and of science itself where, in the latter case, algorithms are non-existent for certain types of problems.

The Philosophy of Data Structures

PARKINSON'S LAW



The consequences of the forgoing subject-matter that are germane to this course.

PARKINSON'S LAW

"Work expands so as to fill the time available for its completion."

C. Northcote Parkinson

Naval historian

Cyril Northcote Parkinson was a British naval historian and author of some 60 books, the most famous of which was his best-seller Parkinson's Law, which led him to be also considered as an important scholar in public administration and management. Wikipedia



"Parkinson's Law of Algorithmic Complexity"

The complexity of algorithms expands at a rate that is roughly proportional to the increase in the available computational power.

C. Northcote Parkinson



"Parkinson's Law of Algorithmic Complexity"

We study data structures so that we can learn to write more efficient programs. But why must programs be efficient when new computers are faster every year? The reason is that our ambitions grow with our capabilities. Instead of rendering efficiency needs obsolete, the modern revolution in computing power and storage capability merely raises the efficiency stakes as we attempt more complex tasks.

Data Structures and Algorithm Analysis

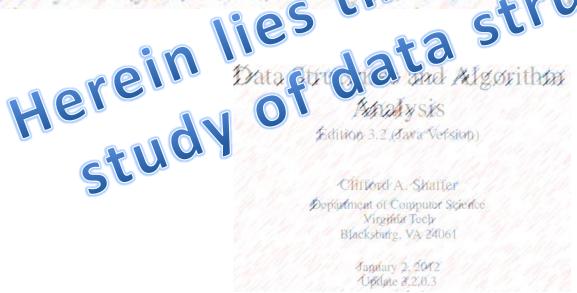
Edition 3.2 (Java Version)

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Definition: Data Structure

