

# COMP251: DATA STRUCTURES & ALGORITHMS

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\* Some slides from “Java Programming: Program Design Including Data Structures”  
by Chris Kiekintveld

# Recursion

# Recursion

- Sometimes, the best way to solve a problem is by solving a **smaller version** of the exact same problem first
- Recursion is a technique that solves a problem by solving a **smaller problem of the same type**
- A procedure that is defined in terms of itself

# Recursion

- When you turn that into a program in computer, you end up with functions that call themselves:

*Recursive Functions*

# Review: Calling Methods

```
int x(int n) {  
    int m = 0;  
    n = n + m + 1;  
    return n;  
}
```

What does y(3) return?

```
int y(int n) {  
    int m = 1;  
    n = x(n);  
    return m + n;  
}
```

# Calling Methods

- Methods can call other methods
- Can a method call itself?
- Yes! This is called a recursive method (function)

# Example

Trace the execution of test(4)

```
void test(int n) {  
    if (n > 0) {  
        System.out.println(n);  
        test(n-1);  
        System.out.println(n);  
    }  
}
```

# Example

What's behind this function ?

```
public int f(int a) {  
    if (a==0 || a==1)  
        return(1);  
    else  
        return(a * f(a-1));  
}
```



# Example

What's behind this function ?

```
public int f(int a) {  
    if (a==0 || a==1)  
        return(1);  
    else  
        return(a * f(a-1));  
}
```

It computes  $a!$  (factorial)

# Factorial

Factorial:

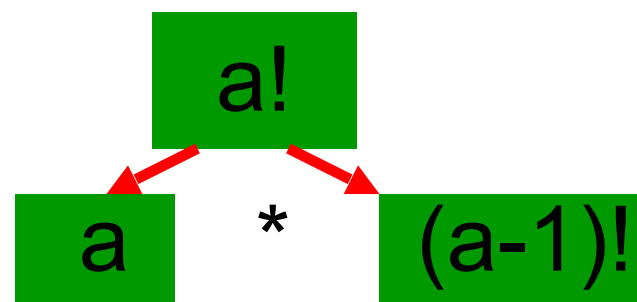
$$a! = 1 * 2 * 3 * \dots * (a-1) * a$$

Note:

$$a! = a * (a-1)!$$

remember:

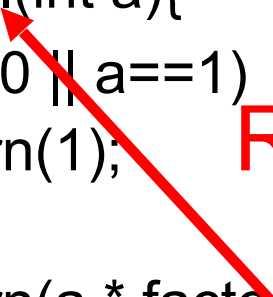
...splitting up the problem into a smaller problem of the same type...



# Tracing the example

```
public int factorial(int a){  
    if (a==0 || a==1)  
        return(1);  
    else  
        return(a * factorial(a-1));  
}
```

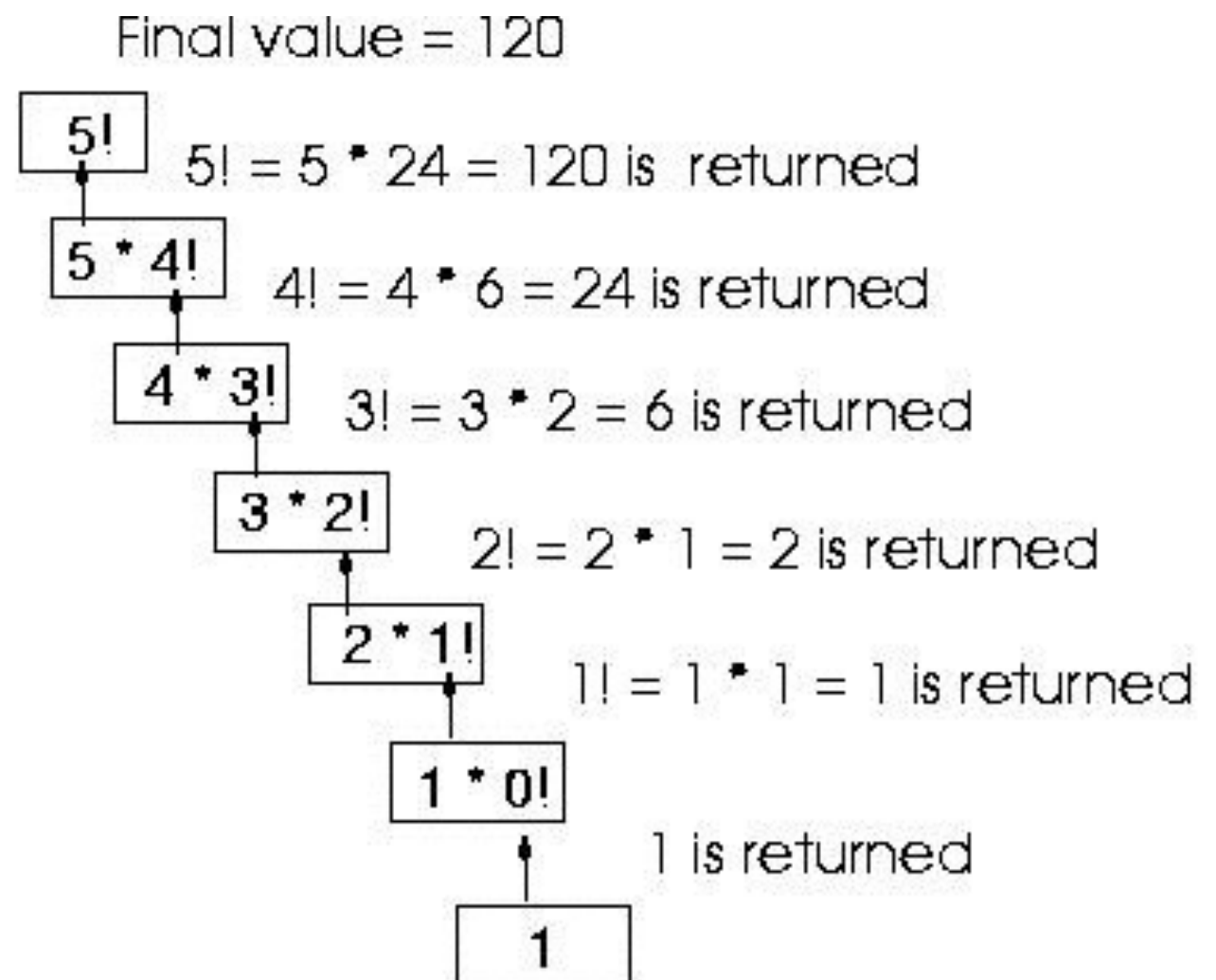
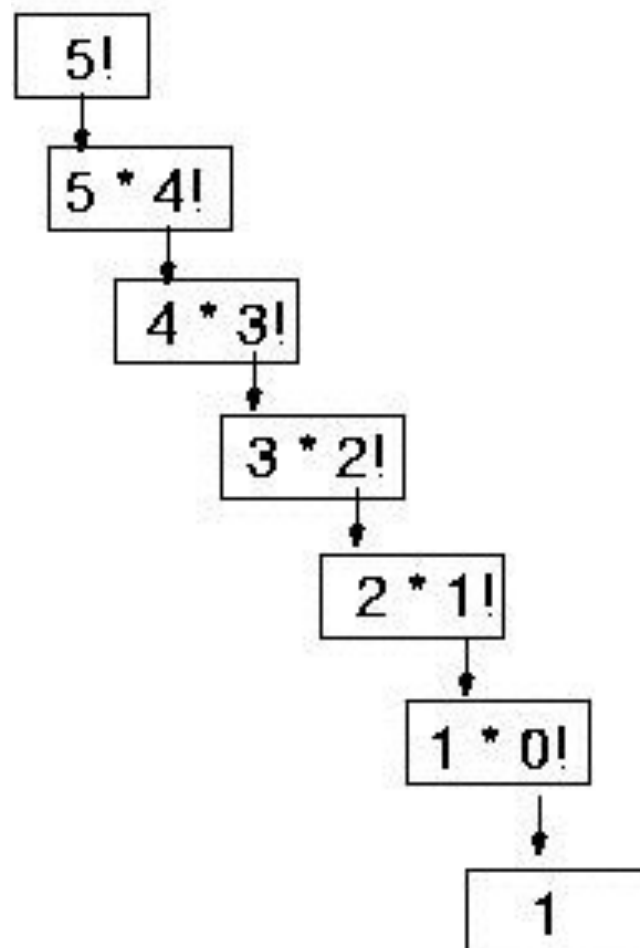
**RECURSION !**



# Tracing the example

```
public int factorial(int a){  
    if (a==0 || a==1)  
        return(1);  
    else  
        return(a * factorial(a-1));  
}
```

**RECURSION !**



# Tracing a Recursive Method

- As always, go line by line
- Recursive methods may have *many copies*
- Every method call creates a new copy and transfers flow of control to the new copy
- Each copy has its own:
  - code
  - parameters
  - local variables

# Tracing a Recursive Method

- After completing a recursive call:
  - Control goes back to the calling environment
  - Recursive call must execute completely before control goes back to the previous call
  - Execution in previous calls begins from point immediately following recursive call

# Watching the Stack

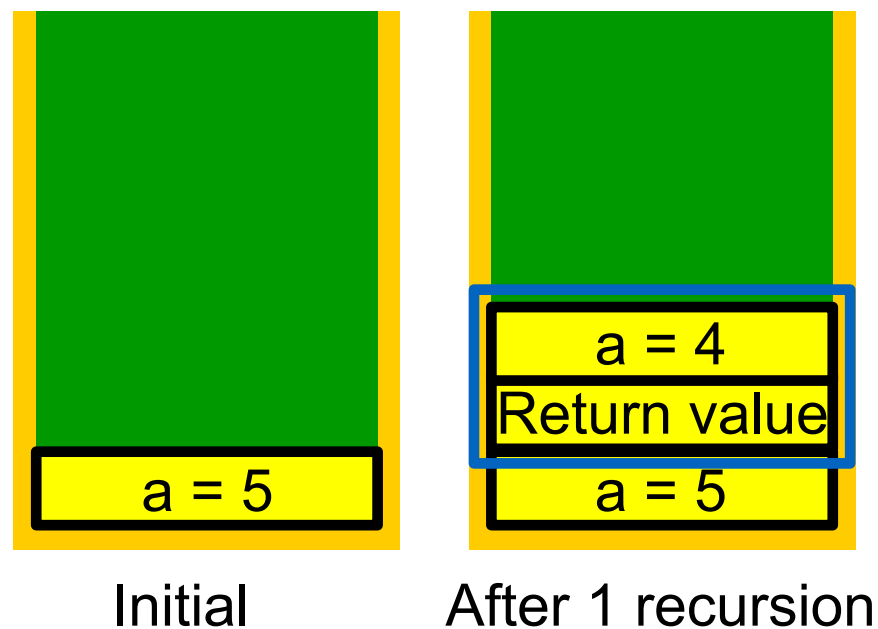
```
public int factorial(int a){  
    if (a==0 || a==1)  
        return(1);  
    else  
        return(a * factorial(a-1));  
}
```



Initial

# Watching the Stack

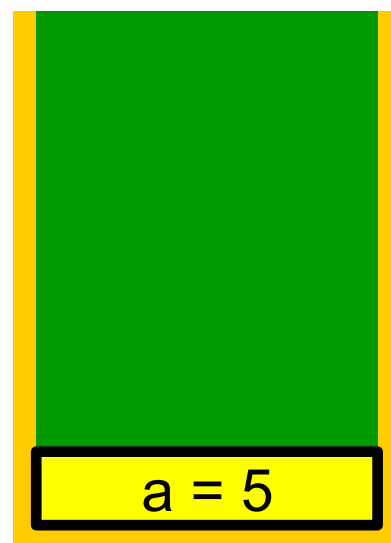
```
public int factorial(int a){  
    if (a==0 || a==1)  
        return(1);  
    else  
        return(a * factorial(a-1));  
}
```



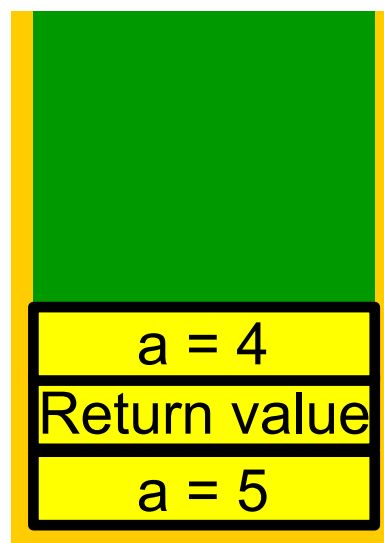


# Watching the Stack

```
public int factorial(int a){  
    if (a==0 || a==1)  
        return(1);  
    else  
        return(a * factorial(a-1));  
}
```

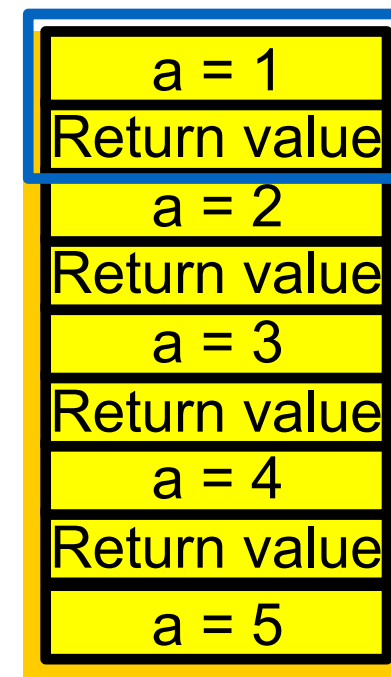


Initial



After 1 recursion

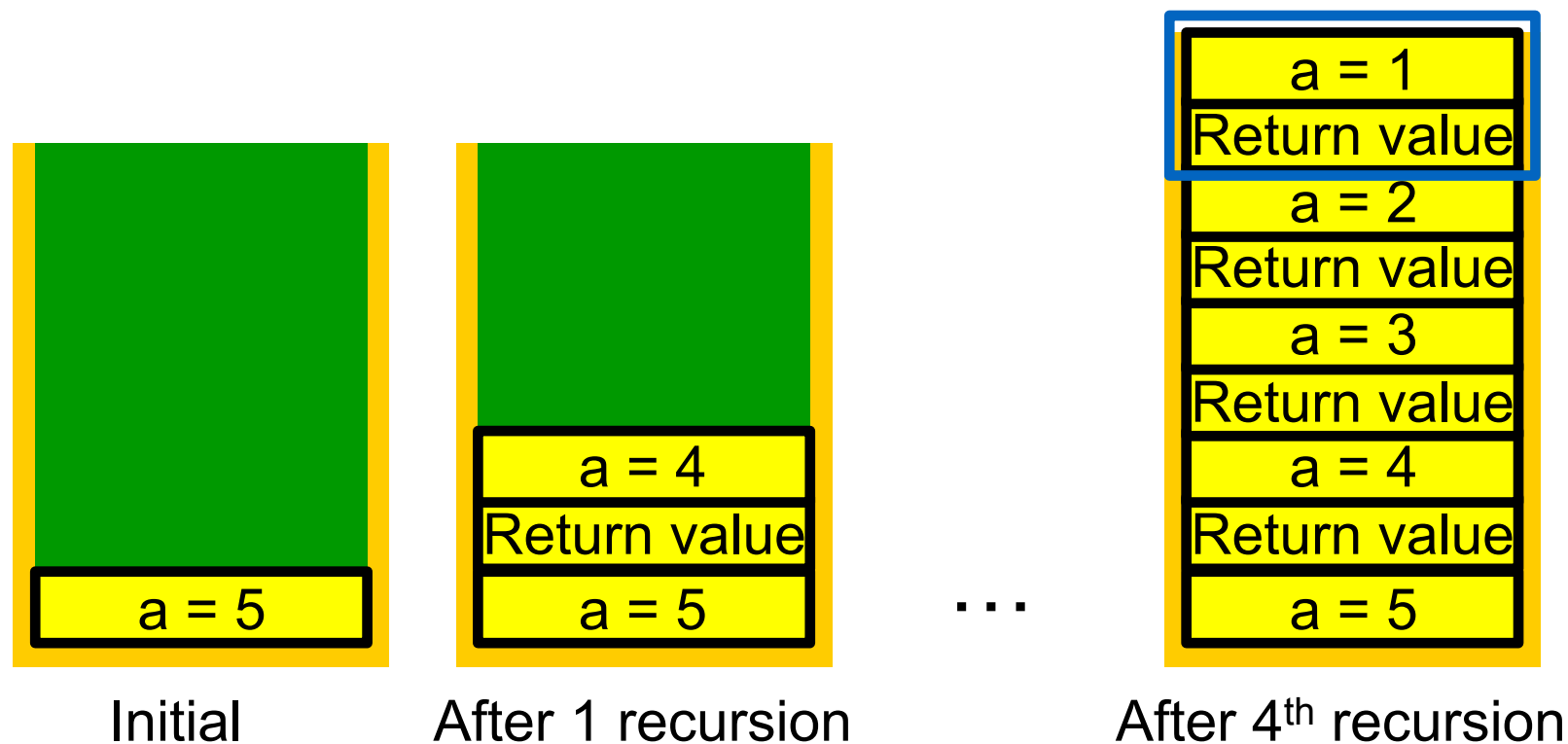
...



After 4<sup>th</sup> recursion

# Watching the Stack

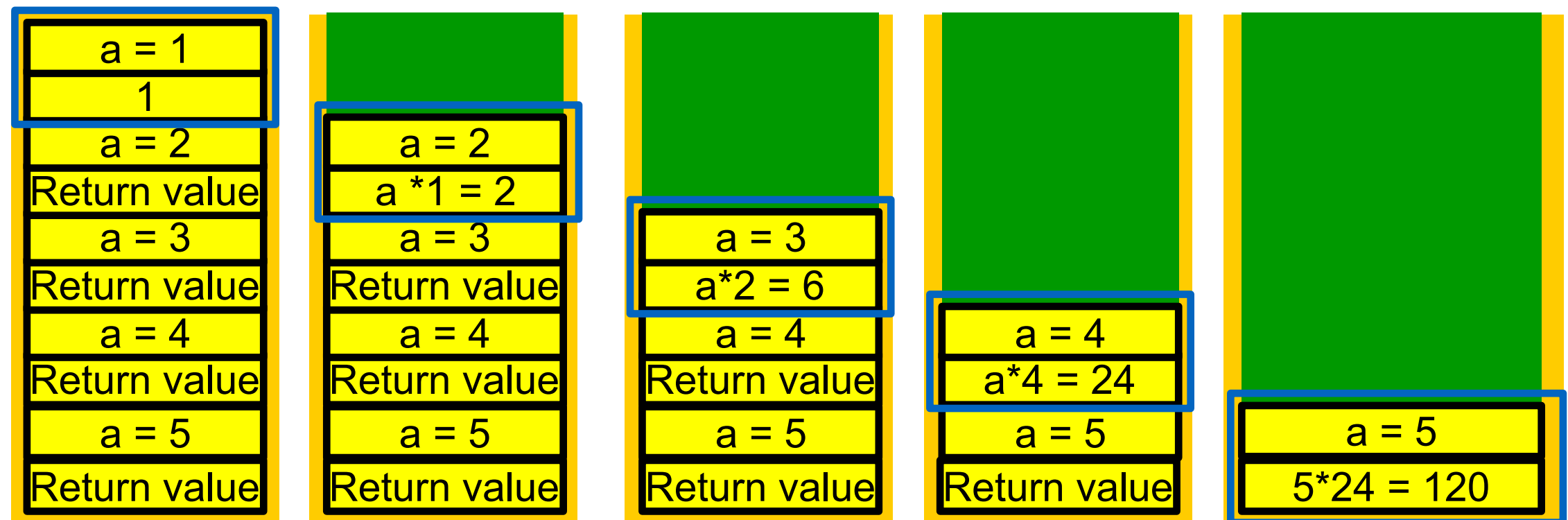
```
public int factorial(int a){  
    if (a==0 || a==1)  
        return(1);  
    else  
        return(a * factorial(a-1));  
}
```



Every call to the method creates a new set of local variables (Stack Frame)!

# Watching the Stack

```
public int factorial(int a){  
    if (a==0 || a==1)  
        return(1);  
    else  
        return(a * factorial(a-1));  
}
```



After 4<sup>th</sup> recursion

Result

# Iterative Factorial

Factorial can be write in iterative way

```
public int f(int a) {  
    int result = 1  
    for (int i = 1; i <= a; i++)  
        result *= i;  
    return result;  
}
```

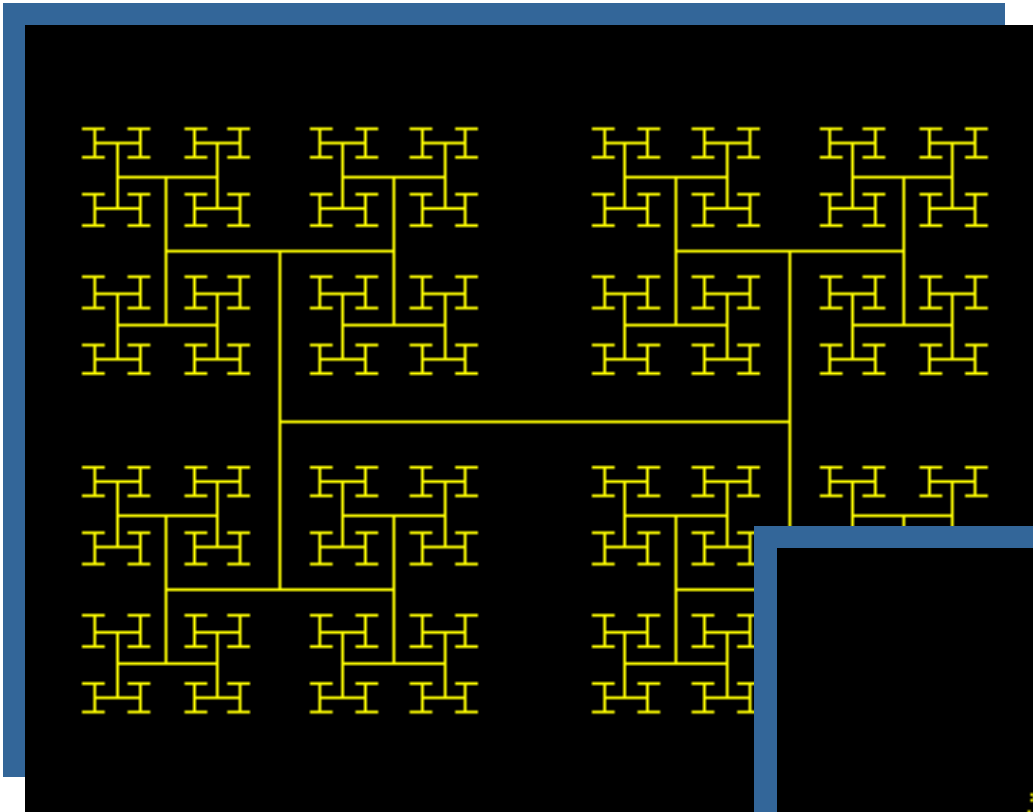
# Recursion vs. Iteration

- Iteration can be used in place of recursion
  - An iterative algorithm uses a *looping construct*
  - A recursive algorithm uses a *branching structure*
- Recursive solutions are often less efficient, in terms of both *time* and *space*, than iterative solutions
- Recursion can simplify the solution of a problem, often resulting in shorter, more easily understood source code
- (Nearly) every recursively defined problem can be solved iteratively
  - iterative optimization can be implemented after recursive design

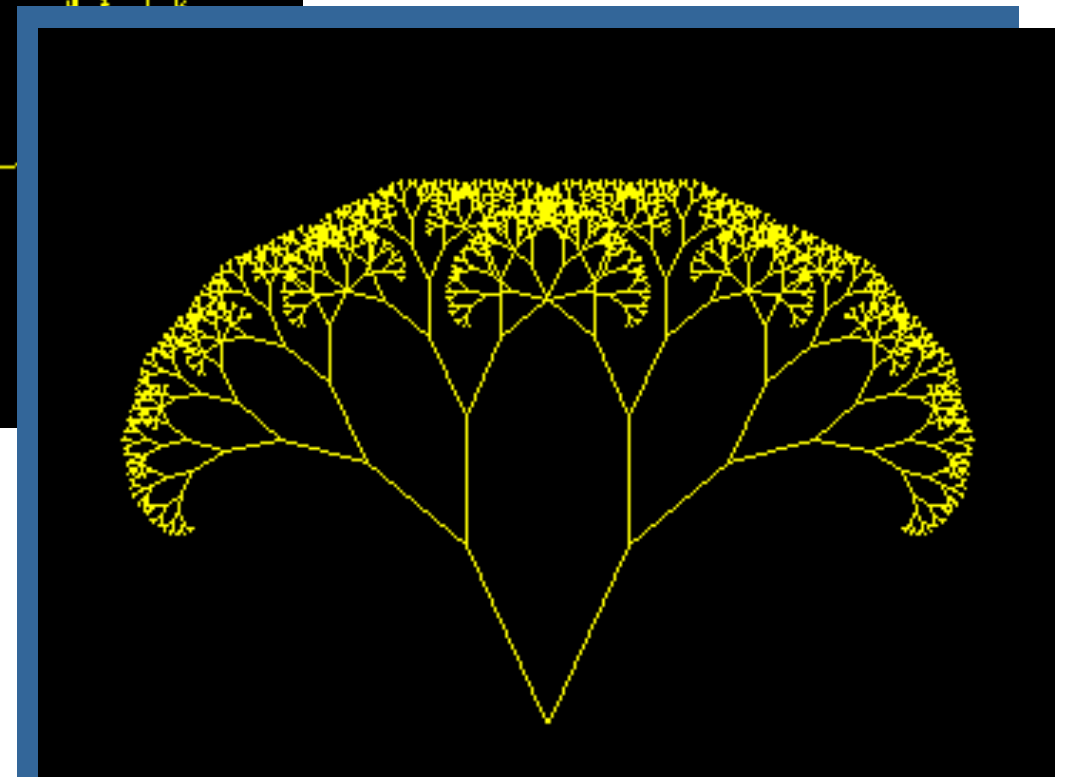
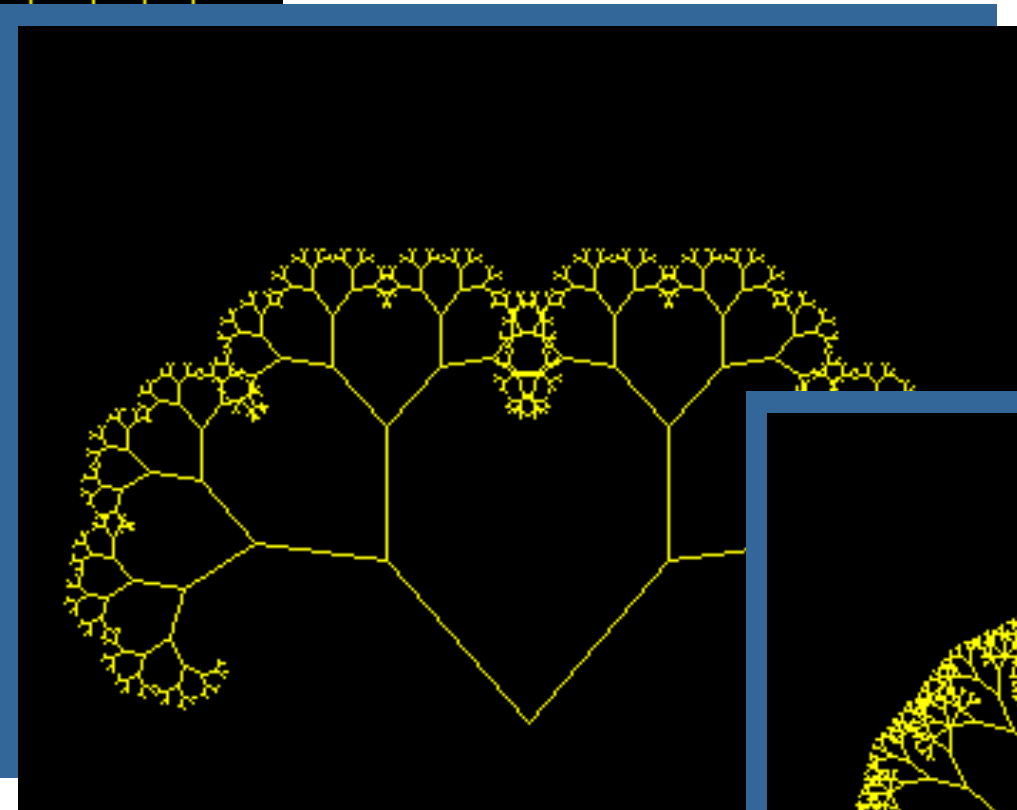
# Deciding whether to use a Recursive Function

- When the depth of recursive calls is relatively “shallow”
- The recursive version does about the same amount of work as the non-recursive version
- The recursive version is shorter and simpler than the non-recursive solution

# Examples: Fractal Tree



<http://id.mind.net/~zona/mmts/geometrySection/fractals/tree/treeFractal.html>



# Design Recursive Algorithms



# Designing Recursive Algorithms

- ♦ General strategy: “Divide and Conquer”
- ♦ Questions to ask yourself
  - ♦ How can we reduce the problem to smaller version of the same problem?
  - ♦ How does each call make the problem smaller?
  - ♦ What is the base case?
  - ♦ Will you *always* reach the base case?

# Properties of Recursive Functions

Problems that can be solved by recursion have these characteristics:

- One or more stopping cases have a simple, nonrecursive solution
- The other cases of the problem can be reduced (using recursion) to problems that are closer to stopping cases
- Eventually the problem can be reduced to only stopping cases, which are relatively easy to solve

Follow these steps to solve a recursive problem:

- Try to express the problem as a simpler version of itself
- Determine the stopping cases (**base case** of recursion)
- Determine the recursive steps (**recursive call**)

# Solution

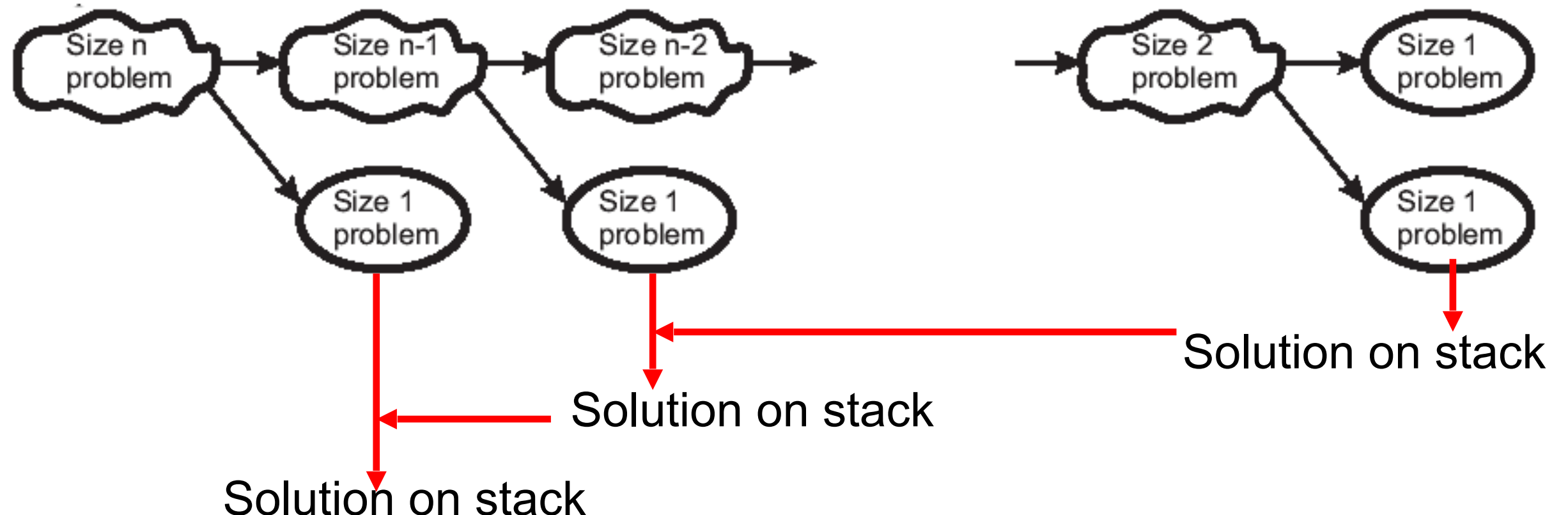
The recursive algorithms we write generally consist of an if statement:

IF

the stopping case is reached solve it

ELSE

split the problem into simpler cases using recursion



# Exercise

Write a recursive function that prints the numbers 1...n in descending order:

```
public void descending(int n) {  
  
  
  
  
  
  
  
  
  
}
```

# Exercise

Write a recursive function that prints the numbers 1...n in descending order:

```
public void descending(int n) {  
    if (n <= 0) return;  
    System.out.println(n);  
    descending(n-1);  
}
```

# Exercise

Write a recursive function to perform exponentiation  
return  $x^m$ , assuming  $m \geq 0$

```
public int exp(int x, int m) {  
  
  
  
  
  
  
  
  
}
```

# Exercise

Write a recursive function to perform exponentiation  
return  $x^m$ , assuming  $m \geq 0$

```
public int exp(int x, int m) {  
    if (m == 0) { return 1; }  
    if (m == 1) { return x; }  
    return x * exp(x, m-1);  
}
```

```
public static boolean p(string s, int i, int f)
    if (i < f) {
        if (s[i] == s[f]) {
            return p(s, i+1, f-1);
        } else {
            return false;
        }
    } else {
        return true;
    }
}
```

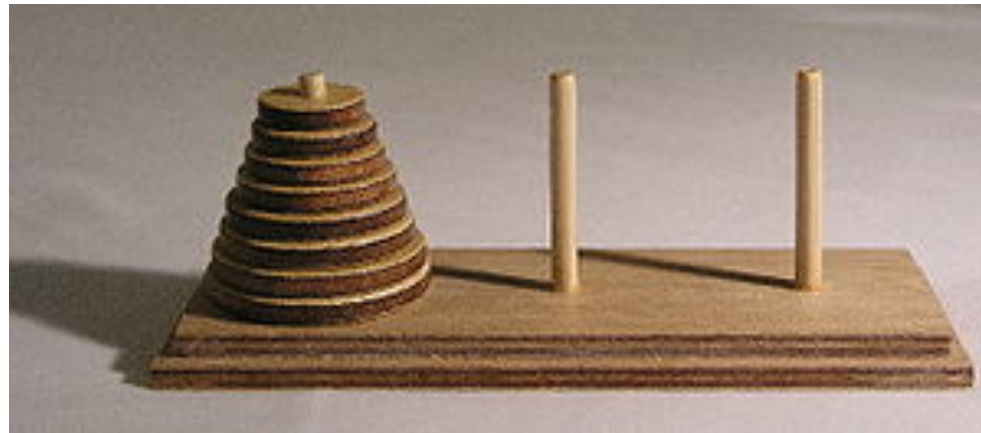
What does `p(s,0,s.length-1)` return

- a) if `s = "UTEP"`
- b) if `s = "SAMS"`
- c) if `s = "kayak"`
- d) if `s = "ABBA"`



# Towers of Hanoi

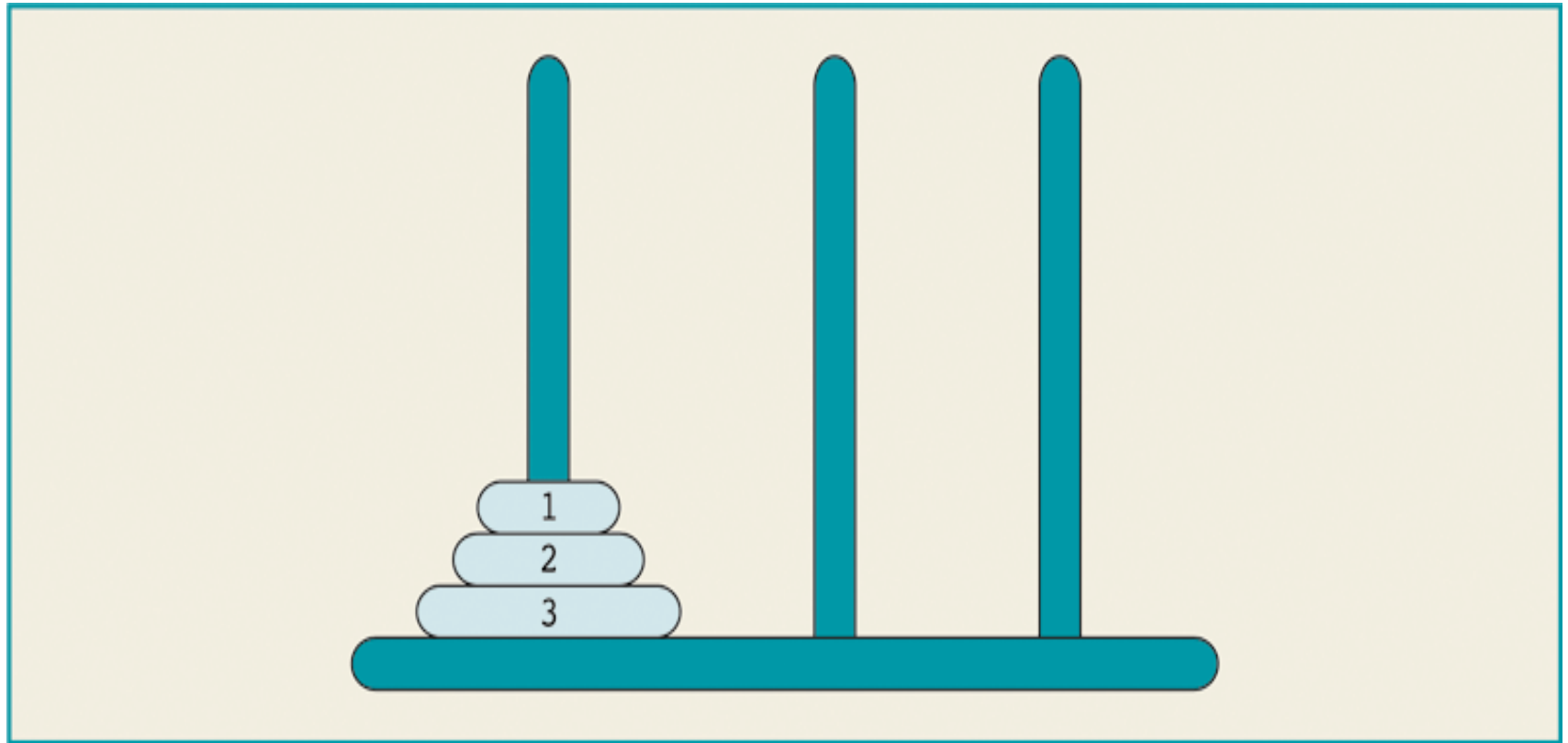
- ♦ The legend of the temple of Brahma
  - ♦ 64 golden disks on 3 pegs
  - ♦ The universe will end when the priest move all disks from the first peg to the last



# The Rules

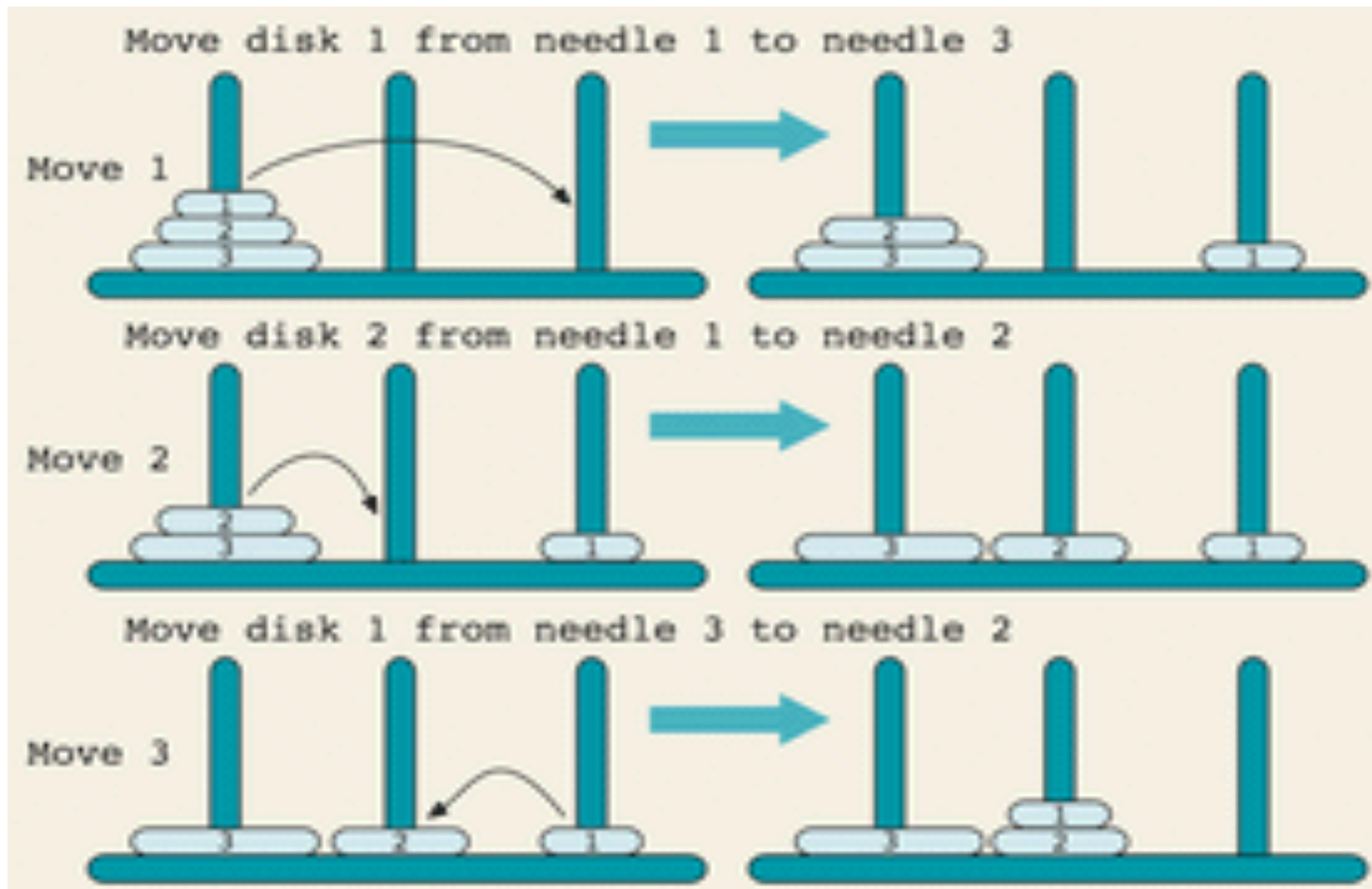
- ◆ Only move one disk at a time
- ◆ A move is taking one disk from a peg and putting it on another peg (on top of any other disks)
- ◆ Cannot put a larger disk on top of a smaller disk
- ◆ With 64 disks, at 1 second per disk, this would take roughly 585 billion years

# Towers of Hanoi: Three Disk Problem

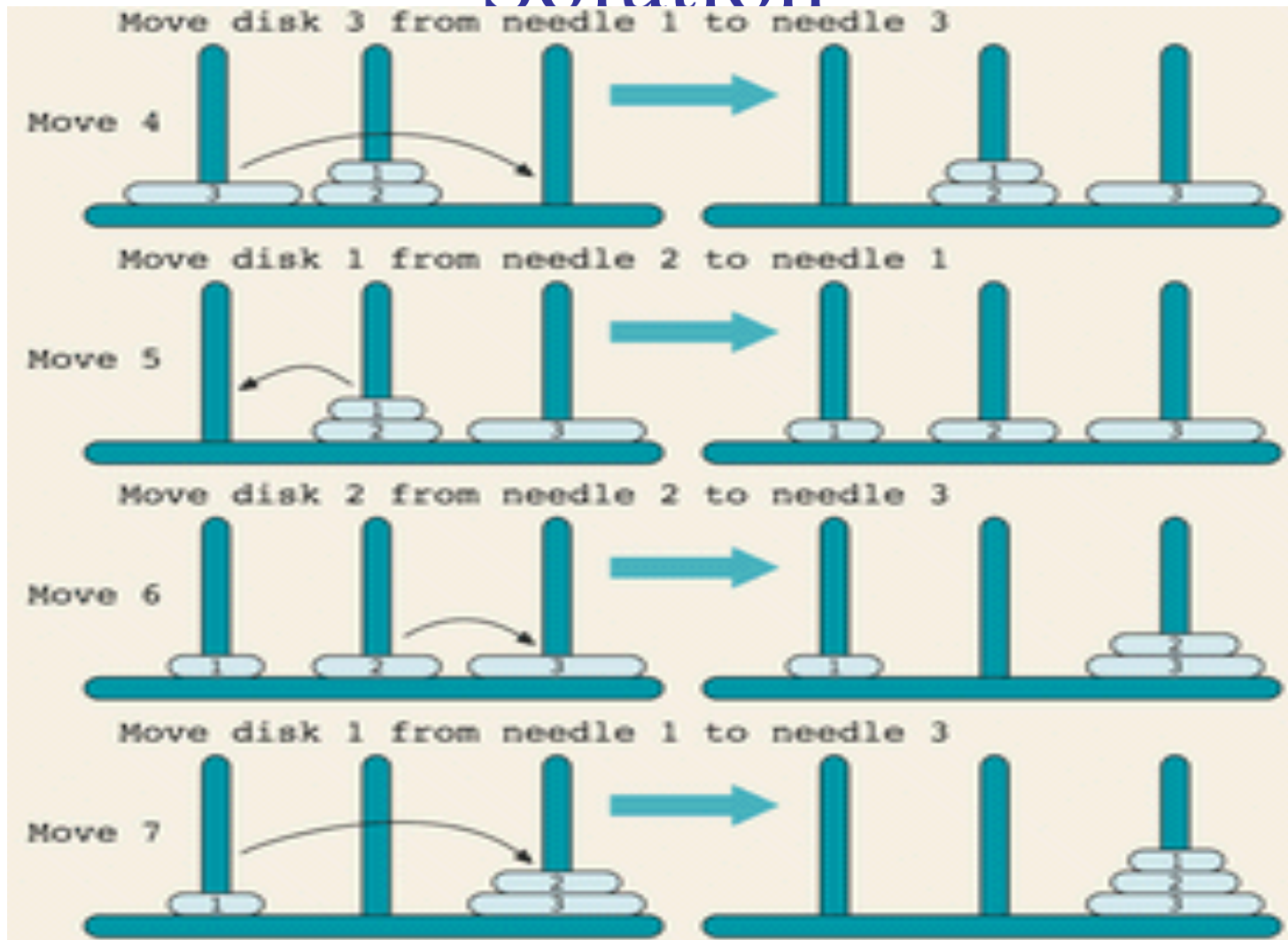


**Figure 14-6** Tower of Hanoi problem with three disks

# Towers of Hanoi: Three Disk Solution



# Towers of Hanoi: Three Disk Solution



# Four disk solution (courtesy wikipedia)



# Recursive algorithm idea

- ◆ Final step is to move the bottom disk from peg 1 to peg 3
- ◆ To do this, the other  $n-1$  disks must be on peg 2
- ◆ So, we need an way to move  $n-1$  disks from peg 1 to peg 2
- ◆ Base case: moving the smallest disk is easy (you can always move it to any peg in one step)

# Pseudocode

```
solveTowers (count, source, destination, spare) {  
  if (count == 1) {  
    move directly  
  }  
  else {  
    solveTowers(count-1, source, spare, destination)  
    solveTowers(1, source, destination, spare)  
    solveTowers(count-1, spare, destination, source)  
  }  
}
```