

<https://www.topcoder.com/community/data-science/data-science-tutorials/an-introduction-to-recursion-part-1/>

<https://www.topcoder.com/community/data-science/data-science-tutorials/an-introduction-to-recursion-part-2/>

Go through above links to get a very fine idea of recursion .

Let us discuss a nice example . There is a 2D array filled with some values and we are at (1,1) cell ( assume 1-based indexing ) . Let us define some directions , w.r.t (i,j) cell

(i+1,j) be direction 1

(i+1,j+1) be direction 2

Now from (i,j) cell we can go in either of these two directions and as we go down we take the sum of values we visit . Now we want to maximize this sum . Suppose the array is :

	1	2	3	4
1	5			
2	10	15		
3	30	20	15	
4	100	200	300	50

Maximum sum which we can create is  $5 + 15 + 20 + 300 = 340$

Now a possible solution is to create all possible sequences and output the maximum one . Now how do we get an idea that this problem can be solved with recursion ?? Now suppose the max. sum which we can create by starting from (2,1) be x and (2,2) be y . Then the our answer is  $\max(5+x, 5+y)$  . Now our problem is to know the values of x and y . And similarly to get value x we need max. sum sequence starting from (3,1) and (3,2) and for y we need starting ones from (3,2) and (3,3) . So this problem can be broken down into similar subproblems with similar property . So this can be solved with recursion

Here is the pseudocode for that:

i,j represent the current cell and n,m represent the size of matrix.

```
Int max_sum( int i , int j ){  
    If(i>n || j>m) return 0;  
    Return max( a[i][j] + max_sum(i+1,j) , a[i][j] + max_sum(i+1,j+1) );  
}
```

This solution has exponential complexity but this can be solved in a much better way in  $O(N^2)$  with dynamic programming approach but let's leave that for some other tutorial .