CSC236 Fall 2016
Assignment #1: induction
due October 7th, 10 p.m.

The aim of this assignment is to give you some practice with various forms of induction. For each question below you will present a proof by induction. For full marks you will need to make it clear to the reader that the base case(s) is/are verified, that the inductive step follows for each element of the domain (typically the natural numbers), where the inductive hypothesis is used and that it is used in a valid case.

Your assignment must be typed to produce a PDF document (hand-written submissions are not acceptable). You may work on the assignment in groups of 1, 2, or 3, and submit a single assignment for the entire group on MarkUs

1. Consider the Fibonacci-esque function g:

$$g(n) = \begin{cases} 1, & \text{if } n = 0 \\ 3, & \text{if } n = 1 \\ g(n-2) + g(n-1) & \text{if } n > 1 \end{cases}$$

Use complete induction to prove that if n is a natural number greater than 1, then $2^{n/2} \leq g(n) \leq 2^n$. You may **not** derive or use a closed-form for $g(n)$ in your proof.

Sample solution: Proof, using complete induction.

inductive step: Let n be a typical natural number greater than 1 and assume $H(n)$: Every natural number $i \in \{2, \ldots, n-1\}$ satisfies $2^{i/2} \leq g(i) \leq 2^i$.

show that inductive conclusion follows: We'll derive $C(n)$: $2^{n/2} \leq g(n) \leq 2^n$.

Base cases: $1 < n < 4$: $g(2) = 4$ and $g(3) = 7$ # by the definition of $g(2)$ and $g(3)$.

$$2^{2/2} = 2 \leq 4 = g(2) \leq 4 = 2^2 \qquad \text{and} \qquad 2^{3/2} = 2\sqrt{2} \leq 7 = g(3) \leq 8 = 2^3$$

$C(2)$ and $C(3)$ follow from our assumptions in this case.

Case $n \geq 4$: By assumptions $H(n-2)$ and $H(n-1)$ # $n \geq 4$ implies $2 \leq n-2, n-1 < n$:

$$2^{(n-2)/2} \leq g(n-2) \leq 2^{n-2} \qquad \text{and} \qquad 2^{(n-1)/2} \leq g(n-1) \leq 2^{n-1}.$$

Substituting these inequalities into the definition of $g(n)$ # by definition of $g(n), n \geq 4 > 0$:

$$g(n) = g(n-2) + g(n-1) \geq 2^{(n-2)/2} + 2^{(n-1)/2} = (1 + \sqrt{2})2^{(n-2)/2} \geq 2 \times 2^{(n-2)/2} = 2^{n/2}$$
$$g(n) = g(n-2) + g(n-1) \leq 2^{n-2} + 2^{n-1} = (1+2)2^{n-2} \leq 2^2 \times 2^{n-2} = 2^n$$

$C(n)$ follows from our assumptions in this case.

In all cases $H(n)$ implies $C(n)$.

2. Suppose B is a set of binary strings where each binary string is of length n. n is positive (greater than 0), and no two strings in B differ in fewer than 2 positions. Use simple induction to prove that B has no more than 2^{n-1} elements.

Sample solution: Proof, using simple induction.

verify base: There are two binary strings of length 1: "0" and "1", and they differ from each other in exactly one position (i.e. fewer than 2). That means that the only sets of binary strings of length 1 that contain no pairs that differ in fewer than 2 positions are {"1"} and {"0"}, which each have no more than $1 = 2^{1-1}$ elements, verifying the claim in this case.

inductive step: Let n be a typical natural number greater than 0. Assume $H(n)$: any set of binary strings of length n containing no pairs that differ in fewer than 2 positions must have no more than 2^{n-1} elements.

derive conclusion $C(n)$: We must show that from $H(n)$ follows $C(n)$: any set of binary strings of length $n+1$ containing no pairs that differ in fewer than 2 positions must have no more than 2^n strings.

Let B be an arbitrary set of binary strings of length $n+1$ that contains no pairs that differ in fewer than 2 positions.

Let B_1 be the subset of B consisting of those elements with 1 in the first position, and B_2 be the subset of B consisting of those elements of B with 0 in the first position. B_1 and B_2 partition B, since every element of B has **either** a 1 **or** a 0 in the first position, and no element of B has **both** a 1 **and** a 0 in the first position.

From B_1 construct B_1', consisting of the strings of B_1 with the leading 1 removed. Similarly, from B_2 construct B_2', consisting of the strings of B_2 with the leading 0 removed.

Sets B_1' and B_2' contain strings of length n, and contain no pairs that differ in fewer than 2 positions, since removing the leading 1s or 0s cannot change the number of positions in which elements differ. By assumption $H(n)$, both B_1' and B_2' have no more than 2^{n-1} elements each.

Elements of B_1 are in 1–1 correspondence with those of B_1', since you can transform one into the other by prepending a leading 1, or removing a leading 1. Similarly elements of B_2 are in 1–1 correspondence with those of B_2'. $|B_1| = |B_1'|$ and $|B_2| = |B_2'|$, since they are in 1–1 correspondence.

$|B| = |B_1| + |B_2| \leq 2^{n-1} + 2^{n-1} = 2^n$, since B_1 and B_2 partition B. This is what $C(n)$ claims.

3. Define T as the smallest set of strings such that:

   (a) "b" $\in T$

   (b) If $t_1, t_2 \in T$, then $t_1 +$ "ene" $+ t_2 \in T$, where the $+$ operator is string concatenation.

   Use structural induction to prove that if $t \in T$ has n "b" characters, then t has $2n-2$ "e" characters.

Sample solution: Proof, using structural induction.

verify basis: "b" $\in T$ # from definition. "b" has 1 character "b" and $2(1) - 2 = 0$ "e" characters. This verifies the claim for the basis.

inductive step: Let $t_1, t_2 \in T$ and assume $H(\{t_1, t_2)$: If t_1 has n_1 "b" characters and t_2 has n_2 "b" characters, then t_1 has $2n_1 - 2$ "e" characters and t_2 has $2n_2 - 2$ "e" characters.

show that inductive conclusion follows from assumptions: We'll derive $C(t_1 +$ "ene" $+ t_2)$: If $t_1 +$ "ene" $+ t_2$ has $n_{1,2}$ "b" characters, then it has $2n_{1,2} - 2$ "e" characters.

$t_1 +$ "ene" $+ t_2 \in T$ # by definition of T, where $+$ is string concatenation.

Let n_1, m_1 be the number of "b" (respectively "e") characters in t_1, and n_2, m_2 be the number of "b" (respectively "e") characters in t_2. $t_1 + "ene" + t2$ has $n_1 + n_2$ "b" characters # Concatenating "ene" doesn't increase the number of "b" characters.
Let $n_{1,2}$ be the number of "b" characters in $t_1 + "ene" + t_2$. Then $n_{1,2} = n_1 + n_2$ # since no "b" characters are added by concatenating "ene".
$t_1 + "ene" + t_2$ has $2n_1 - 2$ $+2n_2 - 2$ $+2$ "e" characters # by assumptions $H(t_1), H(t_2)$, and two "e" characters in "ene".
Summing up $t_1 + "ene" + t_2$ has

$$2n_1 - 2 + 2n_2 - 2 + 2 = 2(n_1 + n_2) - 2 - 2 + 2 = 2(n_1 + n_2) - 2 = 2n_{1,2} - 2$$

... "e" characters. Conclusion $C(t_1 + "ene" + t_2)$ follows in this case.

4. On page 79 of the Course Notes the quantity $\phi = (1 + \sqrt{5})/2$ is shown to be closely related to the Fibonacci function. You may assume that $1.61803 < \phi < 1.61804$. Complete the steps below to show that ϕ is irrational.

   (a) Show that $\phi(\phi - 1) = 1$.

   Sample solution: Substituting the expression for ϕ:

   $$\phi(\phi - 1) = \left(\frac{1 + \sqrt{5}}{2}\right)\left(\frac{\sqrt{5} - 1}{2}\right) = \frac{4}{4} = 1$$

   (b) Rewrite the equation in the previous step so that you have ϕ on the left-hand side, and on the right-hand side a fraction whose numerator and denominator are expressions that may only have integers, $+$ or $-$, and ϕ. There are two different fractions, corresponding to the two different factors in the original equation's left-hand side. Keep both fractions around for future consideration.

   Sample solution: I can choose to divide 1 by either ϕ or $\phi - 1$, yielding:

   $$\phi = \frac{1}{\phi - 1} \qquad \phi = \frac{1 + \phi}{\phi}$$

   (c) Assume, for a moment, that ϕ is the ratio of two natural numbers. Let $m, n \in \mathbb{N}$ such that $\phi = n/m$. Re-write the right-hand side of both equations in the previous step so that you end up with fractions whose numerators and denominators are expressions that may only have integers, $+$ or $-$, m and n.

   Sample solution: Substitute n/m for ϕ on the right-hand side, and then simplify:

   $$\phi = \frac{1}{\phi - 1} \longrightarrow \phi = \frac{m}{n - m} \qquad \phi = \frac{1 + \phi}{\phi} \longrightarrow \frac{m + n}{n}$$

   (d) Use your assumption from the previous part to construct a non-empty subset of the natural numbers that contains m. Use the Principle of Well-Ordering, plus one of the two expressions for ϕ from the previous step to derive a contradiction.

   Sample solution: Let $F \subseteq \mathbb{N}$ be defined by:

   $$F = \{m' \in \mathbb{N} \mid \exists n' \in \mathbb{N}, \phi = n'/m'\}.$$

   By assumption in (c), F is non-empty, since it has at least one member, m. By PWO F has a smallest element, let it be m_0, with its corresponding n_0 so that $m_0, n_0 \in \mathbb{N}$ and $\phi = n_0/m_0$..

Rewriting the equation for ϕ and using the assumption $1.61803 < \phi < 1.61804$ yields:

$$
\begin{array}{rcll}
\phi & = & \dfrac{n_0}{m_0} & \\
\phi m_0 & = & n_0 & \text{\#multiply both sides by } m_0 \\
m_0 & < & n_0 < 2m_0 & \text{\# multiply } 1.61803 < \phi < 1.61804 \text{ by } m \\
0 & < & n_0 - m_0 < m_0 & \text{\#subtract } m_0 \text{ from both inequalities.}
\end{array}
$$

$n_0 - m_0 \in \mathbb{N}$ #integers closed under $-$ and difference is non-negative.
$n_0 - m_0 \in F$, since $\phi = n_0/m_0 = m_0/(n_0 - m_0)$ and $n_0 - m_0 < m_0$.
Contradiction $\rightarrow\leftarrow$. m_0 is the smallest element of F, by construction.

(e) Combine your assumption and contradiction from the previous step into a proof that ϕ cannot be the ratio of two natural numbers. Extend this to a proof that ϕ is irrational.

Sample solution: Proof (by contradiction) that ϕ is irrational.
Assume, for the sake of contradiction, that ϕ is rational.
Let $z_1, z_2 \in \mathbb{Z}, \phi = z_1/z_2$ # by definition of ϕ is rational
Let $n, m \in \mathbb{N}, m/n = \phi$. # Since $\phi = z_1/z_2 > 0$ the numerator and denominator have the same sign. If $z_1, z_2 > 0$, let $n = z_1, m = z_2$. Otherwise, if $z_1, z_2 < 0$, let $n = -z_1, m = -z_2$.
Contradiction $\longrightarrow\longleftarrow$. From the previous part, there are no natural numbers m, n such that $\phi = m/n$.
ϕ is irrational, since assuming otherwise leads to a contradiction.

5. Consider the function f, where $3 \div 2 = 1$ (integer division, like 3//2 in Python):

$$f(n) = \begin{cases} 1 & \text{if } n = 0 \\ f^2(n \div 3) + 3f(n \div 3) & \text{if } n > 0 \end{cases}$$

Use complete induction to prove that for every natural number n greater than 2, $f(n)$ is a multiple of 7. **NB**: Think carefully about which natural numbers you are justified in using the inductive hypothesis for.

Sample solution: Proof by complete induction.

inductive step: Let n be a typical natural number greater than 2, and assume $H(n)$: that $f(i)$ is a multiple of 7 for natural numbers $2 < i < n$.

show that inductive conclusion follows: We'll derive $C(n) : f(n)$ is a multiple of 7.

Base case $2 < n < 6$: $n > 0$, so by the definition of $f(n)$:

$$f(n) = f^2(n \div 3) + 3f(n \div 3) = f^2(1) + 3f(1) = 28 \qquad \# f(1) = 4 \text{ from definition}$$

28 is a multiple of 7, so $C(n)$ follows in this case.

Base case $6 \leq n < 9$: $n > 0$, so by the definition of $f(n)$:

$$f(n) = f^2(n \div 3) + 3f(n \div 3) = f^2(2) + 3f(2) = 28 \qquad \# f(2) = 4 \text{ from definition}$$

28 is a multiple of 7, so $C(n)$ follows in this case.

Case $n \geq 9$: $n > n \div 3 > 2$, so by assumption $H(n \div 3)$ we know that $f(n \div 3)$ is a multiple of 7. Let $k \in \mathbb{N}$ be a natural number such that $f(n \div 3) = 7k$.

$$f(n) = f^2(n \div 3) + 3f(n \div 3) = 49k^2 + 21k = 7(7k^2 + 3k)$$

$7(7k^2 + 3k)$ is a multiple of 7, so the conclusion $C(n)$ is verified in this case.