## [Solution] CS5222 Assignment 1 (due date: Sunday 23:59 pm, week 4)

## Chapter 1

- P6. This elementary problem begins to explore propagation delay and transmission delay, two central concepts in data networking. Consider two hosts, A and B, connected by a single link of rate *R* bps. Suppose that the two hosts are separated by *m* meters, and suppose the propagation speed along the link is *s* meters/sec. Host A is to send a packet of size *L* bits to Host B.
  - a. Express the propagation delay,  $d_{prop}$ , in terms of m and s.
  - b. Determine the transmission time of the packet, d<sub>trans</sub>, in terms of L and R.
  - Ignoring processing and queuing delays, obtain an expression for the endto-end delay.
  - d. Suppose Host A begins to transmit the packet at time t = 0. At time  $t = d_{trans}$ , where is the last bit of the packet?
  - e. Suppose  $d_{prop}$  is greater than  $d_{trans}$ . At time  $t = d_{trans}$ , where is the first bit of the packet?
  - f. Suppose  $d_{prop}$  is less than  $d_{trans}$ . At time  $t = d_{trans}$ , where is the first bit of the packet?
  - g. Suppose  $s = 2.5 \cdot 10^8$ , L = 120 bits, and R = 56 kbps. Find the distance m so that  $d_{\text{prop}}$  equals  $d_{\text{trans}}$ .

## **Problem 6**

- a)  $d_{prop} = m/s$  seconds.
- b)  $d_{trans} = L/R$  seconds.
- c)  $d_{end-to-end} = (m/s + L/R)$  seconds.
- d) The bit is just leaving Host A.
- e) The first bit is in the link and has not reached Host B.
- f) The first bit has reached Host B.
- g) Want

$$m = \frac{L}{R} s = \frac{120}{56 \times 10^3} (2.5 \times 10^8) = 536 \text{ km}.$$

- P31. In modern packet-switched networks, including the Internet, the source host segments long, application-layer messages (for example, an image or a music file) into smaller packets and sends the packets into the network. The receiver then reassembles the packets back into the original message. We refer to this process as *message segmentation*. Figure 1.27 illustrates the end-to-end transport of a message with and without message segmentation. Consider a message that is  $8 \cdot 10^6$  bits long that is to be sent from source to destination in Figure 1.27. Suppose each link in the figure is 2 Mbps. Ignore propagation, queuing, and processing delays.
  - a. Consider sending the message from source to destination without message segmentation. How long does it take to move the message from the source host to the first packet switch? Keeping in mind that each switch uses store-and-forward packet switching, what is the total time to move the message from source host to destination host?
  - b. Now suppose that the message is segmented into 800 packets, with each packet being 10,000 bits long. How long does it take to move the first packet from source host to the first switch? When the first packet is being sent from the first switch to the second switch, the second packet is being sent from the source host to the first switch. At what time will the second packet be fully received at the first switch?
  - c. How long does it take to move the file from source host to destination host when message segmentation is used? Compare this result with your answer in part (a) and comment.
  - d. In addition to reducing delay, what are reasons to use message segmentation?
  - e. Discuss the drawbacks of message segmentation.

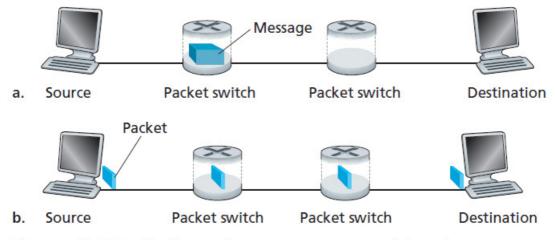


Figure 1.27 ◆ End-to-end message transport: (a) without message segmentation; (b) with message segmentation

## **Problem 31**

- a) Time to send message from source host to first packet switch =  $\frac{8\times10^6}{2\times10^6} \sec = 4\sec \text{ With store-and-forward switching, the total time to move}$   $\max \text{message from source host to destination host} = 4\sec\times3 \text{ hops} = 12\sec$
- b) Time to send  $1^{st}$  packet from source host to first packet switch = .  $\frac{1\times10^4}{2\times10^6} \sec = 5 \, m \sec$ . Time at which  $2^{nd}$  packet is received at the first switch = time at which  $1^{st}$  packet is received at the second switch =  $2\times5m \sec = 10 \, m \sec$
- c) Time at which  $1^{st}$  packet is received at the destination host =  $5 \, m \sec \times 3 \, hops = 15 \, m \sec$ . After this, every 5msec one packet will be received; thus time at which last  $(800^{th})$  packet is received =  $15 \, m \sec + 799 * 5m \sec = 4.01 \sec$ . It can be seen that delay in using message segmentation is significantly less (almost  $1/3^{rd}$ ).

d)

- i. Without message segmentation, if bit errors are not tolerated, if there is a single bit error, the whole message has to be retransmitted (rather than a single packet).
- ii. Without message segmentation, huge packets (containing HD videos, for example) are sent into the network. Routers have to accommodate these huge packets. Smaller packets have to queue behind enormous packets and suffer unfair delays.

e)

- i. Packets have to be put in sequence at the destination.
- ii. Message segmentation results in many smaller packets. Since header size is usually the same for all packets regardless of their size, with message segmentation the total amount of header bytes is more.

**Question**. WhatsApp and WeChat are smartphone real-time messaging systems. After doing some research on the Internet, for each of these systems, write one paragraph to briefly describe the protocols they use.

WhatsApp uses a customized version of the open standard Extensible Messaging and Presence Protocol (XMPP).

(https://en.wikipedia.org/wiki/WhatsApp)

It may be difficult to find the details of WeChat's protocol due to its proprietary nature and the lack of its protocol specifications. See for example "Fine-Grained Dissection of WeChat in Cellular Networks" by Q. Huang et al.

(www.cse.cuhk.edu.hk/~pclee/www/pubs/iwqos15chatdissect.pdf)