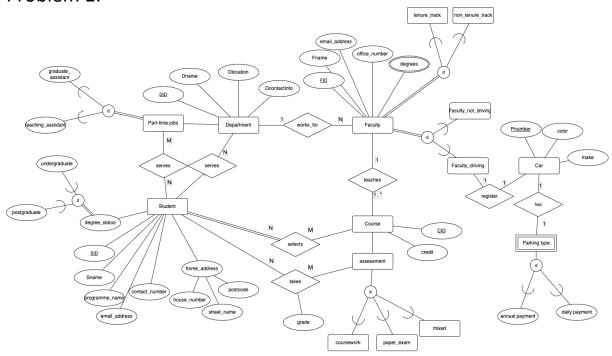
Problem 1:



Problem 2:

1. the strong entities:

Publisher (URL) Publisher (<u>URL</u>, Pname, Paddress)

Library_Item (IID) Library_Item (<u>IID</u>, Title)
Borrower (BID) Borrower (<u>BID</u>, Bname)

2. the weak entity:

without any weak entity in this case.

3. one-to-one (1:1) relationship:

without any one-to-one relationship types in this case.

4. one-to-many (1:N) relationship:

Since Publishes and Borrow are both one-many relationships, it does not require a table. We just need to put the primary key from many-side into the 1-side, as foreign keys.

Publishes Publisher (<u>URL</u>, Pname, Paddress, IID*)
Borrow Borrower (BID, Bname, IID*, dates)

5. many-to-many (N:M) relationship:

Since Reserved and Write are both many-to-many relationship, creating a new table and pulling primary keys from others as composite primary key.

Write Write (Bnumber*, AID*, URL*)
Reserved (BID*, IID*, Order ID*)

6. Mapping of multivalued attributes

Phone numbers Publisher phone numbers (PNID*, Phone number)

7. Mapping of N-ary relationship types

without N-ary relationship types

8. Options for mapping Specialization or Generalization: Borrower (BID, Bname, Staff, Student) Staff (BID*) Student (BID*)

allows to store more information in the future about public sector Publisher (URL, Pname, Paddress, Publisher type) Public_sector (Publisher_type*) Private sector (Publisher type*)

Library Item (IID, Title) Book (IID*) Journal (<u>IID</u>*, date, Issue_no, Volume) Magazine (IID*, date, Issue no) Conference_proceedings (IID*, date, Location)

Collecting all the tables together:

Borrower (BID, Bname, IID*, dates) Write (Bnumber*, AID*, URL*) Reserved (BID*, IID*, Order ID*) Publisher_phone_numbers (PNID*, Phone_number) Authors (AID) Borrower (BID, Bname, Staff, Student) Staff (BID*) Student (BID*) Publisher (URL, Pname, Paddress, Publisher_type, IID*) Public sector (Publisher type*) Private sector (Publisher type*) Library Item (IID, Title) Book (IID*) Journal (IID*, date, Issue no, Volume)

Problem 3:

Magazine (IID*, date, Issue no)

Conference_proceedings (IID*, date, Location)

3.1

table Library Item, it allows the library to store more type of items such as CDs, Newspaper and so on. However, it causes too many tables.

table Borrower, it allows the library to save more borrower, even though there are normally 2 type of people, neither staff or student. it might cause too many tables as well. table Publisher, we can see that we have 2 higher level entities which are public sector and private sector, each one may have more data itself.

The table Book is a good candidate for horizontal partitioning.

IID	Bname	Desc	Horizontal Partition(Hash based)								
1	а		IID	Bname	Desc		Ð	Bname	Desc		
2	b		1	а			2	b			
3	С		3	С			4	d			
4	d		User table-Shard1				User table-Shard2				

There are normally many books. This partition reduces the speed of reading operations.

select * from Book where IID = 1 and IID = 3;

select * from Book where IID = 2 and IID = 4;

3.3

Table Publisher is a good candidate for vertical partitioning

URL	Pname	Paddress	Publisher_type	URL	Pname	Paddress		URL	Publisher_type	
url1	PN1	addr1	public	url1	PN1	addr1		url1	public	
url2	PN2	addr2	private	url2	PN2	addr2		url2	private	
url3	PN3	addr3	public	url3	PN3	addr3		url3	public	
url4	PN4	addr4	private	url4	PN4	addr4		url4	private	
Vertical Partition(Range based)				Uer table				Extension table		

Sometimes, a few basic information of publisher might be checked frequently, or we don't have to know the certain part out of user information.

select URL, Pname, Paddress from Publisher;

Problem 4:

4.1

 $U_R = \{ F \}$ $N = U - U_R = \{ A, B \}$

 $AB^{+} = \{ A, B, F \} = R$

Thus, N is a minimal key. The candidate key is { A, B }.

4.2

relation R is not in BCNF. R is atomic so it is in the 1NF. B is partial of AB so there is a partial dependency cause that it is not in the 2NF.

To be BCNF, it must be a 3NF and for any dependency A--->B, A should be a super key.

let D be the set of all attributes of R

 $D = \{ A, B, F \}$

B --> F

 $D = \{ A, B \} => Rnew = \{ F, B \}$

It is now in a collection of BCNF relations.