

Exam

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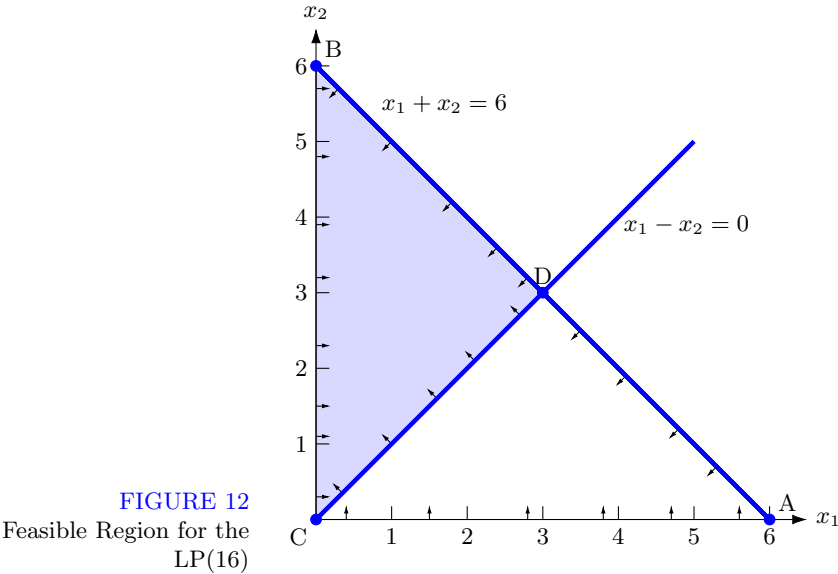


FIGURE 12  
Feasible Region for the  
LP(16)

TABLE 32  
Three Sets of Basic Variables Correspond to Corner point C

Basic Variables	Basic Feasible Solution	Corresponds to Extreme Point
$x_1, x_2$	$x_1 = x_2 = 3, s_1 = s_2 = 0$	D
$x_1, s_1$	$x_1 = 0, s_1 = 6, x_2 = s_2 = 0$	C
$x_1, s_2$	$x_1 = 6, s_2 = -6, x_2 = s_1 = 0$	Infeasible
$x_2, s_1$	$x_2 = 0, s_1 = 6, x_1 = s_2 = 0$	C
$x_2, s_2$	$x_2 = 6, s_2 = 6, s_1 = x_1 = 0$	B
$s_1, s_2$	$s_1 = 6, s_2 = 0, x_1 = x_2 = 0$	C

We can now discuss why the simplex algorithm often is an inefficient method for solving degenerate LPs. Suppose an LP is degenerate. Then there may be many sets(maybe hundreds) of basic variables that correspond to some nonoptimal extreme point.The simplex algorithm might encounter all these sets of basic variables before it finds that it was at a nonoptimal extreme point.This problem was illustrated(on a small scale) in solving(16): The simplex took two pivots before it found that point C was suboptimal. Fortunately,some degenerate LPs have a special structure that enables us to solve them by methods other than the simplex (see,foe example,the discussion of the assignment problem in Chapter 7).

P R O B L E M S  
Group A

- 1 Even if an LP's initial tableau is nondegenerate, later tableaus may exhibit degeneracy. Degenerate tableaus often occur in the tableau following a tie in the ratio test. To illustrate this,solve the following LP:
- Also graph the feasible region and show which extreme points correspond to more than one set of basic variables.
- 2 Find the optimal solution of the following LP:

$$\begin{aligned} \max \quad & z = 5x_1 + 3x_2 \\ \text{s.t.} \quad & 4x_1 + 2x_2 \leq 12 \\ & 4x_1 + x_2 \leq 10 \\ & x_1 + x_2 \leq 4 \\ & x_1, x_2 \leq 0 \end{aligned}$$

$$\begin{aligned} \min \quad & z = -x_1 - x_2 \\ \text{s.t.} \quad & x_1 + x_2 \leq 1 \\ & -x_1 + x_2 \leq 0 \\ & x_1, x_2 \geq 0 \end{aligned}$$