

Automated Unit Testing of Large Industrial Embedded Software using Concolic Testing

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Strong IT Industry in South Korea

Time-to-Market?

V.S Quality?





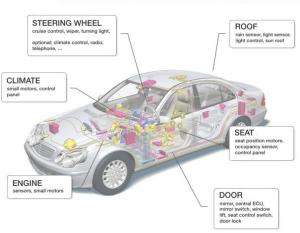


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Summary of the Talk

- Embedded SW is becoming larger and more complex
 - Ex. Android: 12 MLOC, Tizen > 6 MLOC
- Smartphone development period is very short
 - No time to manually test smartphones sufficiently
- Solution: Automated unit test generation for industrial embedded SW using CONBOL (CONcrete and symBOLic testing)
 - CONBOL automatically generates unit-test driver/stubs
 - CONBOL automatically generates test cases using concolic testing
 - CONBOL targets crash bugs (i.e. null pointer dereference, etc.)
- CONBOL detected 24 crash bugs in 4 MLOC Android SW in 16 hours











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- Motivation
- Background on concolic testing
- Overview of CONBOL
 - Unit test driver/stub generator
 - Pre-processor module
- Real-world application: Project S on Samsung smartphones
- Lessons learned and conclusion



Motivation

- Manual testing of SW is often ineffective and inefficient
 - Ineffectiveness: SW bugs usually exist in corner cases that are difficult to expect
 - Inefficiency: It is hard to generate a sufficient #
 of test cases in a given amount of project time
- For consumer electronics, these limitations become more threatening
 - Complex control logic
 - Large software size
 - Short development time
 - Testing platform limitation



Concolic Testing

- Combine concrete execution and symbolic execution
 - Concrete + Symbolic = Concolic
- In a nutshell, concrete execution over a concrete input guides symbolic execution
 - Symbolic execution is performed along with a concrete execution path
- Automated test case generation technique
 - Execute a target program on automatically generated test inputs
 - All possible execution paths are to be explored
 - Higher branch coverage than random testing



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Industrial Experience w/ Concolic Testing

Target platform: Samsung smartphone platforms

Testing Level	Target Programs	Results	Publication
Unit- testing	Busybox Is	Detected 4 bugs and covered 98% of branches	Kim et al. [ICST12]
	Samsung security library	Detected 1 memory bug and covered 73% of branches	Kim et al. [ICST12]
System- testing	Samsung Linux Platform (SLP) file manager	Detected 1 infinite loop bug and covered 20% of branches	Kim et al. [FSE11]
	10 Busybox utilities	Detected 1 bug in grep and covered 80% of branches	
	Libexif	Detected 6 bugs including 2 security bugs registered in Common Vulnerabilities and Exposures, and covered 43% of branches	Kim et al. [ICSE12]

Obstacles of Concolic Testing for Industrial Embedded SW

- 1. Each execution path can be very long, which causes a huge state space to analyze
 - Generating and running test cases on embedded platforms would take significant amount of time
- 2. Porting of a concolic testing tool to a target embedded OS can be difficult
 - Due to resource constraint of embedded platforms
- 3. Embedded SW often uses target-specific compiler extensions



Solutions of CONBOL

- 1. Automatically generate unit tests including test drivers/stubs
 - We can apply concolic testing on industrial embedded SW that has 4 MLOC
- 2. Test embedded SW on a host PC
 - Most unit functions can run on a host PC
 - Only a few unit functions are tightly coupled with target embedded platforms
- 3. Port target-specific compiler extensions to GCC compatible ones

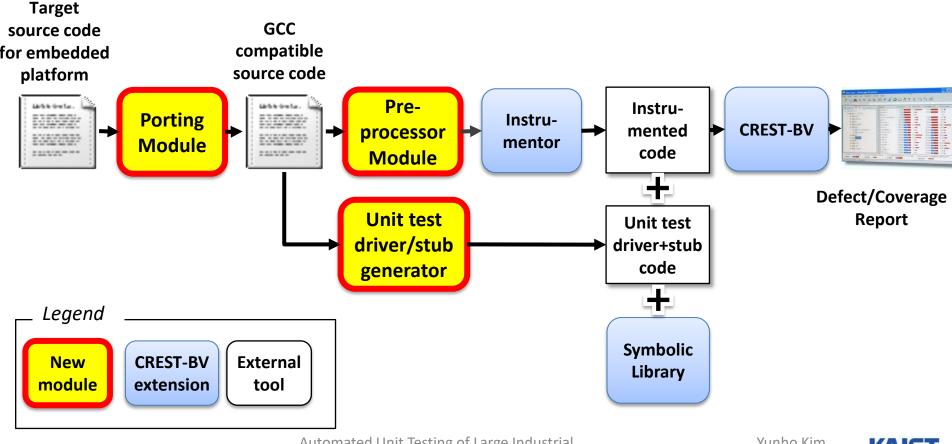


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Overview of CONBOL

 We have developed the CONcrete and symBOLic (CONBOL) framework: an automated concolic unit testing tool based-on CREST-BV for embedded SW



Porting Module

- The porting module automatically modifies the source code of unit functions so that the code can be compiled and executed at the host PC
 - 1. The porting module removes unportable functions
 - Inline ARM assembly code, hardware dependent code, unportable RVCT(RealView Compilation Tools) extensions
 - The porting module translates target code to be compatible with GCC and CIL(C Intermediate Language) which is an instrumentation tool

RVCT	Translation for GCC			
asm {}	Not Portable			
swi (0x01)	Not Portable			
align(8)	_attribute_((aligned(8)))			
packed	attribute((packed))			



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Unit Test Driver/Stub Generator(1/2)

- The unit test driver/stub generator automatically generates unit test driver/stub functions for unit testing of a target function
 - A unit test driver symbolically sets all visible global variables and parameters of the target function

Туре	Description	Code Example
Primitive	set a corresponding symbolic value	<pre>int a; SYM_int(a);</pre>
Array	set a fixed number of elements	<pre>int a[3]; SYM_int(a[0]); SYM_int(a[2]);</pre>
Structure	set NULL to all pointer fields and set symbolic value to all primitive fields	<pre>struct _st{int n,struct _st*p}a; SYM_int(a.n); a.p=NULL;</pre>
Pointer	allocate memory for a pointee and set a symbolic value of corresponding type of the pointee	<pre>int *a; a = malloc(sizeof(int)); SYM_int(*a);</pre>

 The test driver/stub generator replaces sub-functions invoked by the target function with symbolic stub functions



Unit Test Driver/Stub Generator(2/2)

Example of an automatically generated unit-test driver

```
01:typedef struct Node {
02: char c;
03: struct Node *next;
04: Node;
05: Node *head;
06:// Target unit-under-test
                                                           Unit Test Driver
07:void add last(char v) {
08: // add a new node containing v
   // to the end of the linked list
10:
                                                        Generate symbolic inputs
11:// Test driver for the target unit
                                                        for global variables and a
12:void test add last(){
                                                              parameter
     char v1;
13:
                                         Set global
   head = malloc(sizeof(Node));
14:
                                         variables
15:
   SYM char (head->c);
16:
   head->next = NULL;
                                                          Call target function
                      Set parameter
   SYM char(v1);
17:
     add last(v1); }
18:
```



Pre-processor Module

- The pre-processor module inserts probes for three heuristics to improve bug detection precision
 - 1. assert () insertion to detect more bugs
 - 2. Scoring of alarms to reduce false alarms
 - 3. Pre-conditions insertion to reduce false alarms



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Unit-testing Strategy to Reduce False Alarms

 CONBOL raises a false NPD alarm because ctx(line 6) is not correctly initialized by init ctx() (line 8)

```
    init ctx() is replaced with a symbolic stub function

 01:int init ctx(struct CONTEXT &ctx){
 02: ctx.f = malloc(...);
 03: ...
 04: return 0;}
 05:void f() {
 06: struct CONTEXT ctx;
                                init_ctx() is replaced with a
 07: int ret;
                                symbolic stub that does not initialize
 08: ret = init ctx(&ctx);
 09: if (ret == -1) {
                                ctx
 10:
         return; }
                                A false NPD alarm is raised at line 11
 11: if (ctx.f[1] > 0) {
 12:
         /* Some code */
                                 because ctx is not properly initialized
 13:
 14:}
```

 We are developing a technique to automatically identify subfunctions that should not be replaced with stub functions



Inserting assert () Statements

- The pre-processor module automatically inserts
 assert() to cause and detect the following three
 types of run-time failures
 - Out-of-bound memory access bugs(OOB)
 - Insert assert (0<=idx && idx<size) right before array access operations
 - Divide-by-zero bugs(DBZ)
 - Insert assert (denominator!=0) right before division operators whose denominator is not constant
 - Null-pointer-dereference bugs(NPD)
 - Insert assert (ptr!=NULL) right before pointer dereference operations



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Scoring of Alarms (1/2)

- CONBOL assigns a score to each alarm as follows:
 - 1. Every violated assertion(i.e., alarm) gets 5 as a default score.
 - 2. The score of the violated assertion increases by 1 if the assertions contains a variable x which is checked in the target function containing the assertion (e.g., if(x < y+1)...)
 - An explicit check of x indicates that the developer considers x important, and the assertion on x is important consequently.

```
01: void f(int x, int y) {
02:    int array[10];
03:    if (x < 15) {
04:        assert(x<10);
05:        array[x]++;
06:        assert(y<10);
07:        array[y]++;
08:}}</pre>
```

No	Туре	Location	Assert Expression	Score
1	ООВ	src.c:f():4	x<10	6(=5+1)
2	ООВ	src.c:f():6	y<10	5



Scoring of Alarms (2/2)

- 3. For each violated assertion assert (expr), the score of the assertion decreases by 1, if expr appears five or more times in other violated assertions in the entire target software.
 - Developers write code correctly most of the time: target code that is repeated frequently is not likely to be buggy

No	Type	Location	Assert Expression	Score
1	ООВ	src.c:f():1287	A.index - 1 >= 0	4(=5-1)
2	ООВ	src.c:g():1300	A.index - 1 >= 0	4(=5-1)
3	ООВ	src.c:h():1313	A.index - 1 >= 0	4(=5-1)
4	ООВ	src.c:x():1326	A.index - 1 >= 0	4(=5-1)
5	ООВ	src.c:y():1339	A.index - 1 >= 0	4(=5-1)

CONBOL reports alarms whose scores are 6 or above



Inserting Constraints to Satisfy Pre-conditions

- The pre-processor module automatically inserts
 assume() to avoid false alarms due to violation of
 implicit pre-conditions
 - Pre-conditions for array indexes
 - Insert array pre-conditions if the target function does not check an array index variable
 - Pre-conditions for constant parameters
 - Insert constant parameter pre-conditions if the parameter of the target function is one of some constant values for all invocations
 - Ex.) the third parameter of fseek() should be one of SEEK_SET, SEEK_CUR, Or SEEK_END
 - Pre-conditions for enum values
 - CONBOL considers an enum type as a special int type and generates concrete test cases defined in the corresponding the enum type



Inserting Constraints to Satisfy Pre-conditions(1/3)

- An automatically generated unit test driver can violate implicit pre-conditions of a target unit function
 - Violation of implicit pre-conditions raises false alarms

```
01:int array[10];
02:void get_ith_element(int i){
03: return array[i];
04:}
```

Line 3 can raise an OOB alarm because $\dot{\text{\footnote{1}}}$ can be greater than or equal to 10

```
05:// Test driver for get ith element()
```

```
06:void test_get_ith_element(){
07:    int i, idx;
08:    for(i=0; i<10; i++) {
09:        SYM_int(array[i]);
10:    }
11:    SYM_int(idx);
12:
13:    get_ith_element(idx);
14:}</pre>
```

However, developers often assume that get_ith_element() is always called under a pre-condition (0<=i && i<10)



Inserting Constraints to Satisfy Pre-conditions(3/3)

An example of pre-conditions for array index

```
Developers assume that callers of
01:int array[10];
                                       get ith element() performs sanity
02:void get ith element(int i) {
                                       checking of the parameter before they
     return array[i];
                                       invoke get ith element()
04:}
05:// Test driver for get ith element()
06:void test get ith element() {
07: int i, idx;
08: for (i=0; i<10; i++) {
09:
        SYM int(array[i]);
10:
11:
     SYM int(idx);
                                         assume (expr) enforces
12:
     assume(0<=idx && idx<10);
                                         symbolic values to satisfy expr
     get ith element(idx);
13:
14:}
```

Statistics of Project S

- Project S, our target program, is an industrial embedded software for smartphones developed by Samsung Electronics
 - Project S targets ARM platforms

ľ	Data		
Total lines of	About 4,000,000		
# of branches	S	397,854	
# of	Total	48,743	
functions	Having more than one branch	29,324	
# of files	Sources	7,243	
	Headers	10,401	



Test Experiment Setting

- CONBOL uses a DFS strategy used by CREST-BV in Kim et al. [ICSE12 SEIP]
- Termination criteria and timeout setting
 - Concolic unit testing of a target function terminates when
 - CONBOL detect a violation of an assertion, or
 - All possible execution paths are explored, or
 - Concolic unit testing spends 30 seconds (Timeout1)
 - In addition, a single test execution of a target unit should not spend more than 15 seconds (Timeout2)
- HW setting
 - Intel i5 3570K @ 3.4 GHz, 4GB RAM running Debian Linux 6.0.4 32bit



Results (1/2)

- Results of branch coverage and time cost
 - CONBOL tested 86.7%(=25,425) of target functions on a host PC
 - 13.3% of functions were not inherently portable to a host PC due to inline ARM assembly, direct memory access, etc
 - CONBOL covered 59.6% of branches in 15.8 hours

Statistics	Number
Total # of test cases generated	About 800,000
Branch coverage (%)	59.6
Execution time (hour)	15.8
# of functions reaching timtout1 (30s)	742
# of functions reaching timtout2 (15s)	134
Execution time w/o timeout (hour)	9.0



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Results (2/2)

- CONBOL raised 277 alarms
- 2 Samsung engineers (w/o prior knowledge on the target program) took 1
 week to remove 227 false alarms out of 277 alarms
 - We reported 50 alarms and 24 crash bugs were confirmed by the developers of Project S
- Pre-conditions and scoring rules filtered out 14.1% and 81.2% of likely false alarms, respectively
- Note that Coverity prevent could not detect any of these crash bugs

# of reported alarms	Out-of-bound		NULL-pointer- dereference		Divide-by-zero		Total	
	# of alarms	Ratio (%)	# of alarms	Ratio (%)	# of alarms	Ratio (%)	# of alarms	Ratio (%)
W/O any heuristics	3235	100.0	2588	100.0	61	100.0	5884	100.0
W/ inserted pre- conditions	2486	76.8	2511	97.0	58	95.1	5055	85.9
W/ inserted pre- conditions + scoring rules	220	6.8	42	1.6	15	24.6	277	4.7
Confirmed and fixed bugs	13	0.4	5	0.2	6	9.8	24	0.4

Recognition of Success of CONBOL at Samsung Electronics



- Bronze Award
 at Samsung
 Best Paper
 Award
- Oct's Best
 Practice Award



Lessons Learned

- Effective and efficient automated concolic unit testing approach for industrial embedded software
 - Detected 24 critical crash bugs in 4 MLOC embedded SW
- Samsung engineers were sensitive to false positives very much (>10 false/true alarms ratio)
 - False alarm reduction techniques are very important
- We have developed a new automated unit testing platform CONCERT which reduces false alarms by
 - Synthesizing realistic target unit contexts based on dynamic function correlation observed in system testing
 - Utilizing common dynamic invariants of various contexts



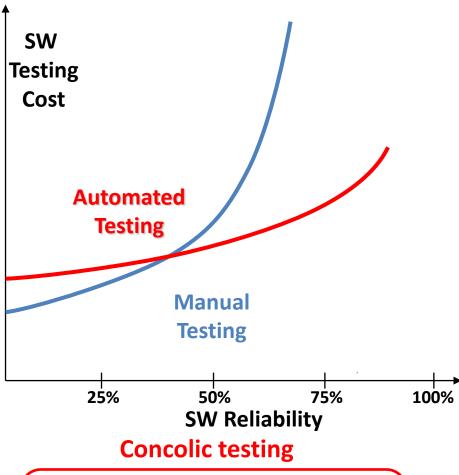
Conclusion

- Automated concolic testing is_effective and efficient for testing industrial embedded software including vehicle domain as well as consumer electronics domain
 - LG electronics introduced the technique from 2014 (c.f. ICSE SEIP 2015 paper)
 - Hyundai motors started to apply the technique from 2015
- Successful application of automated testing techniques requires <u>expertise</u> <u>of human engineers</u>

Traditional testing

- Manual TC gen
- Testing main scenarios
- System-level testing
- Small # of TCs





- Automated TC gen
- Testing exceptional scenarios
- Unit-level testing
- Large # of TCs