Running Time and Sorting

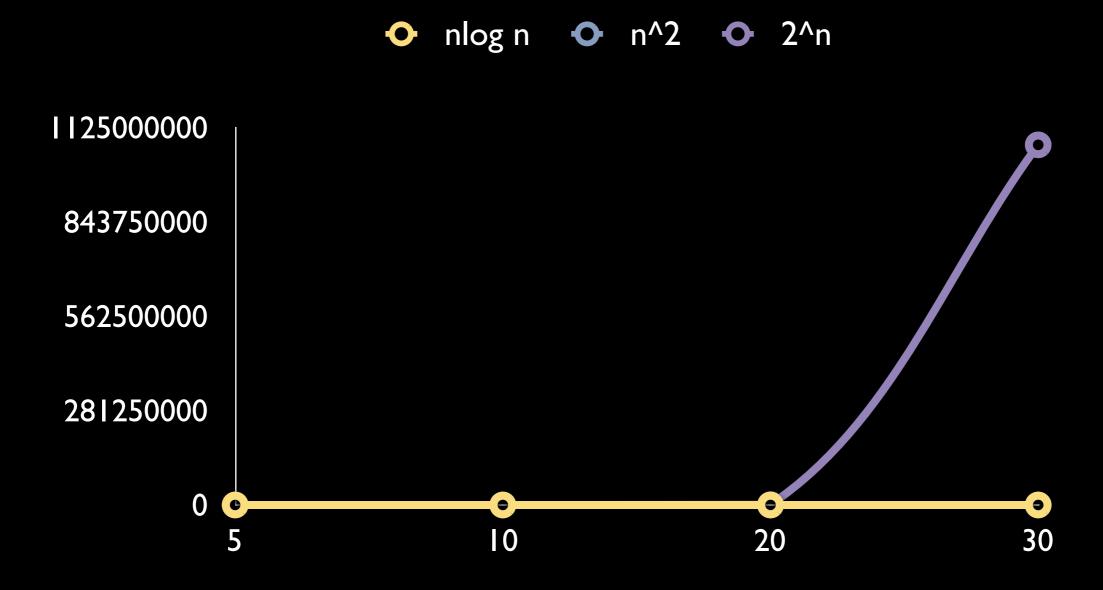
Running Time

- How fast does an algorithm run?
- Usually relative to the size of the input
- Is it reasonable? T(n), T(log n), T(n^2)...
- For even moderate input sizes is it unreasonable? T(2^n)

Examples of running times (in milli seconds)

input size	nlogn	n ²	2 ⁿ	
I	0		2	
10	10	100	1024	. an 19 years'
100	200	10000	1.26765E+30	x10^19 years
1000	3000	1000000	1.07151E+3017	1.4 x 10^290 years!

Function growth



Computing Running Time

- Which parts depend on input size?
- Usually ignore things that aren't input size dependent: constant amount of work on any input
- Places where time is often needed:
 - loops
 - recursive calls

Consider the following two loops:

```
for (int i = 0; i < n; i++)
{
    if list[i] > value
        remove list[i] from list
}
for (int j = 0; j < n; j + +)
{
    display(list[j])
}</pre>
```

```
for (int i = 0; i<n; i++)
{
    for(int j=0; j<n; j++)
    {
        display(i*j);
    }
}</pre>
```

Which one takes less time?

A

B

```
for (int i = 0; i < n; i++)
{
    if list[i] > value
        remove list[i] from list
}
for (int j = 0; j < n; j ++)
{
    display(list[j])
}</pre>
```

```
for (int i = 0; i<n; i++)
{
    for(int j=0; j<n; j++)
    {
        display(i*j);
    }
}</pre>
```

What are their rough running times?

A

B

```
for (int i = 0; i < n; i++)
{
    if list[i] > value
        remove list[i] from list
}
for (int j = 0; j < n; j ++)
{
    display(list[j])
}</pre>
```

```
for (int i = 0; i<n; i++)
{
    for(int j=0; j<n; j++)
    {
        display(i*j);
    }
}</pre>
```

```
for (int i = 0; i < n; i++)
{
    if list[i] > value
       remove list[i] from list
}
for (int j = 0; j < n; j ++)
{
    display(list[j])
}</pre>
```

```
for (int i = 0; i < n; i + +)
{
    for(int j = 0; j < n; j + +)
    {
        display(i*j);
    }
}</pre>
```

```
for (int i = 0; i < n; i++)
{
   if list[i] > value
     remove list[i] from list
}
for (int j = 0; j < n; j++)
{
   display(list[j])
}</pre>
```

```
for (int i = 0; i<n; i++)
{
    for(int j=0; j<n; j++)
    {
        display(i*j);
    }
}</pre>
```

```
for (int i = 0; i < n; i++)
{
    if list[i] > value
        remove list[i] from list
}
for (int j = 0; j < n; j++)
{
    display(list[j])
}
Total: n+n=2n</pre>
```

```
for (int i = 0; i<n; i++)
{
    for(int j=0; j<n; j++)
    {
        display(i*j);
    }
}</pre>
```

```
for (int i = 0; i < n; i++)
{
    if list[i] > value
        remove list[i] from list
}
for (int j = 0; j < n; j++)
{
    display(list[j])
}
Total: 2n</pre>
```

```
for (int i = 0; i < n; i++)
{
    if list[i] > value
        remove list[i] from list
}
for (int j = 0; j < n; j++)
        (display(list[j]))
        Total: 2n</pre>
```

Recursive Calls: Quicksort

```
Quicksort(list a, length x)
if(x<=1) return a;
int p = x-1;
list lower;
list higher;
higher.add(p)
for (int i = 0; i < x-1, i++)
      if(a[i] < p)
          lower.add(a[i]);
      else
          higher.add(a[i]);
Quicksort(lower, lower.length());
Quicksort(higher, higher.length());
return(lower+higher);
```

Recursive Calls: Quicksort

```
QUICKSORT(list a, length x)
             if(x<=1) return a;</pre>
             int p = x-1;
             list lower;
 what's
             list higher;
             higher.add(p)
going on?
             for (int i = 0; i < x-1, i++)
                  if(a[i] < p)
                      lower.add(a[i]);
                  else
                      higher.add(a[i]);
             QUICKSORT(lower, lower.length());
             QUICKSORT(higher, higher.length());
             return(lower+higher);
```

Running time of Quicksort

- Each time the pivot splits the list into two lists
- At each call there at most ~ n^2 comparisons
- consider how it works for: I 2 3 4 5
- how about: I 2 4 5 3

Example: quicksort

12345

Running time of Quicksort

- Each time the pivot splits the list into two lists
- -n^2 comparisons At each call there at most \sim n² comparisons
- consider how it works for: 1 2 3 4 5
- how about: 12453

Example: quicksort

12453

Running time of Quicksort

- Each time the pivot splits the list into two lists
- -n log n comparisons At each call there at most \sim n² comparisons
- consider how it works for: I 2 3 4 5
- how about: 12453

Which Sort is best?

- Many things to consider:
 - range of input
 - type of input
 - unique values or not?
 - integer, real numbers, alphabetical...
 - expected size of input

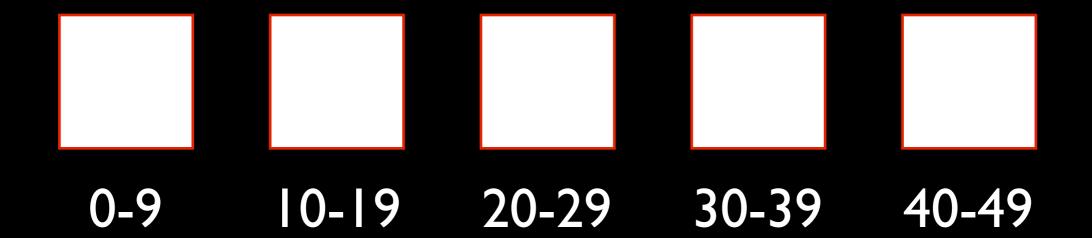
Can we sort better than T(n log n)?

- What if we know something about the data?
- Sorting integers: create a bucket for ranges of values
- if needed, use a comparison base sort in each bucket

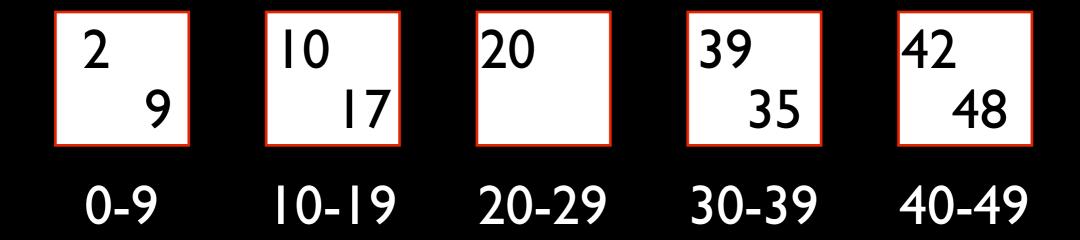
input: 20, 33, 2, 42, 48, 10, 17, 35, 9

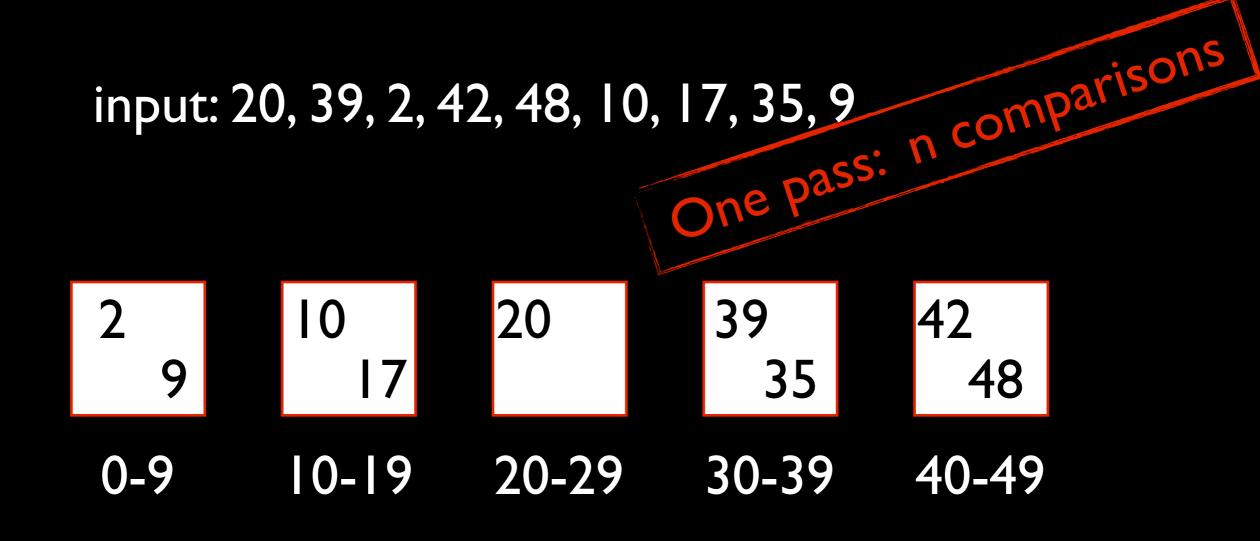
If we know the input is always integers from 0 to 49

input: 20, 33, 2, 42, 48, 10, 17, 35, 9



input: 20, 39, 2, 42, 48, 10, 17, 35, 9





Some Sorting Algorithms ...

Selection sort

Merge sort

Insertion sort

and many more...

- Quicksort
- Bubble sort
- Heap sort
- Bucket sort

Properties of sorting algorithms

- comparison based?
- speed (n logn vs. n²)
- memory usage (in place or not?)
- iterative or recursive?
- stability (do like-valued items stay in original order?)

Properties of some sorting algorithms

- Selection sort: in place, comparison, slow (~n²)
- Insertion sort: low memory needs, comparison, shifting needed (~n²)
- Quicksort: not stable, fast (~nlogn), low memory
- Bubble sort: simple, slow, low memory (~n²)
- Bucket sort: fast for data with little variation (~n)
- Merge sort: fast (~nlogn)