

# LECTURE NOTES ON ALGEBRAIC LOGIC

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## 1. ALGEBRAS AND EQUATIONS

We begin by reviewing some fundamentals of general algebraic systems.

### Definition 1.1.

- (i) A *type* is a map  $\rho: \mathcal{F} \rightarrow \mathbb{N}$ , where  $\mathcal{F}$  is a set of function symbols. In this case,  $\rho(f)$  is said to be the *arity* of the function symbol  $f$ , for every  $f \in \mathcal{F}$ . Function symbols of arity zero are called *constants*.
- (ii) An *algebra* of type  $\rho$  is a pair  $A = \langle A; F \rangle$  where  $A$  is a nonempty set and  $F = \{f^A : f \in \mathcal{F}\}$  is a set of operations on  $A$  whose arity is determined by  $\rho$ , in the sense that each  $f^A$  has arity  $\rho(f)$ . The set  $A$  is called the *universe* of  $A$ .

When  $\mathcal{F} = \{f_1, \dots, f_n\}$ , we shall write  $\langle A; f_1^A, \dots, f_n^A \rangle$  instead of  $\langle A; F \rangle$ . In this case, we often drop the superscripts, and write simply  $\langle A; f_1, \dots, f_n \rangle$ .

Classical examples of algebras are groups and rings. For instance, the type of groups  $\rho_G$  consists of a binary symbol  $+$ , a unary symbol  $-$ , and a constant symbol  $0$ . Then a group is an algebra  $\langle G; +, -, 0 \rangle$  of type  $\rho_G$  in which  $+$  is associative,  $0$  is a neutral element for  $+$ , and  $-$  produces inverses.

Lattices, Heyting algebras, and modal algebras are also algebras in the above sense. For instance, the type of lattices  $\rho_L$  consists of two binary symbols  $\wedge$  and  $\vee$  and a lattice is an algebra  $\langle A; \wedge, \vee \rangle$  of type  $\rho_L$  that satisfies the idempotent, commutative, associative, and absorption laws. Similarly, the type of Heyting algebras  $\rho_H$  consists of three binary operations symbols  $\wedge, \vee$ , and  $\rightarrow$  and of two constant symbols  $0$  and  $1$ . Then a Heyting algebra is an algebra  $\langle A; \wedge, \vee, \rightarrow, 0, 1 \rangle$  such that  $\langle A; \wedge, \vee, 0, 1 \rangle$  is a bounded lattice and, for every  $a, b, c \in A$ ,

$$a \wedge b \leq c \iff a \leq b \rightarrow c. \quad (\text{residuation law})$$

Boolean algebras can be viewed as the Heyting algebras that satisfy the following equational version of the *excluded middle law*:

$$x \vee (x \rightarrow 0) \approx 1.$$

In this case, the complement operation  $\neg x$  can be defined as  $x \rightarrow 0$ .

Perhaps less obviously, even algebraic structures whose operations are apparently *external* can be viewed as algebras in the sense of the above definition. For instance, modules over a ring  $R$  can be viewed as algebras whose type  $\rho_R$  extends that of groups with the unary symbols  $\{\lambda_r : r \in R\}$ . From this point of view, a module over  $R$  is an

algebra  $\langle G; +, -, 0, \{\lambda_r : r \in R\} \rangle$  of type  $\rho_R$  such that  $\langle G; +, -, 0 \rangle$  is an abelian group and, for every  $r, s \in R$  and  $a, c \in G$ ,

$$\begin{aligned}\lambda_r(a + c) &= \lambda_r(a) + \lambda_r(c) \\ \lambda_{r+s}(a) &= \lambda_r(a) + \lambda_s(a) \\ \lambda_r(\lambda_s(a)) &= \lambda_{r \cdot s}(a) \\ \lambda_1(a) &= a.\end{aligned}$$

Given a type  $\rho: \mathcal{F} \rightarrow \mathbb{N}$  and a set of variables  $X$  disjoint from  $\mathcal{F}$ , the set of *terms of type  $\rho$  over  $X$*  is the least set  $T_\rho(X)$  such that

- (i)  $X \subseteq T_\rho(X)$ ;
- (ii) if  $c \in \mathcal{F}$  is a constant, then  $c \in T_\rho(X)$ ; and
- (iii) if  $\varphi_1, \dots, \varphi_{\rho(f)} \in T_\rho(X)$  and  $f \in \mathcal{F}$ , then  $f\varphi_1 \dots \varphi_{\rho(f)} \in T_\rho(X)$ .

For the sake of readability, we shall often write  $f(\varphi_1, \dots, \varphi_{\rho(f)})$  instead of  $f\varphi_1 \dots \varphi_{\rho(f)}$ . Similarly, if  $f$  is a binary operation  $+$ , we often write  $\varphi_1 + \varphi_2$  instead of  $f(\varphi_1, \varphi_2)$ .

Given a term  $\varphi \in T_\rho(X)$ , we write  $\varphi(x_1, \dots, x_n)$  to indicate that the variables occurring in  $\varphi$  are among  $x_1, \dots, x_n$ . Furthermore, given an algebra  $A$  of type  $\rho$  and elements  $a_1, \dots, a_n \in A$ , we define an element

$$\varphi^A(a_1, \dots, a_n)$$

of  $A$ , by recursion on the construction of  $\varphi$ , as follows:

- (i) if  $\varphi$  is a variable  $x_i$ , then  $\varphi^A(a_1, \dots, a_n) := a_i$ ;
- (ii) if  $\varphi$  is a constant  $c$ , then  $c^A$  is the interpretation of  $c$  in  $A$ ;
- (iii) if  $\varphi = f(\psi_1, \dots, \psi_m)$ , then

$$\varphi^A(a_1, \dots, a_n) := f^A(\psi_1^A(a_1, \dots, a_n), \dots, \psi_m^A(a_1, \dots, a_n)).$$

An *equation of type  $\rho$  over  $X$*  is an expression of the form  $\varphi \approx \psi$ , where  $\varphi, \psi \in T_\rho(X)$ . Such an equation  $\varphi \approx \psi$  is *valid* in an algebra  $A$  of type  $\rho$ , if

$$\varphi^A(a_1, \dots, a_n) = \psi^A(a_1, \dots, a_n), \text{ for every } a_1, \dots, a_n \in A,$$

in which case we say that  $A$  *satisfies*  $\varphi \approx \psi$ .

For instance, groups are precisely the algebras of type  $\rho_G$  that satisfy the equations

$$x + (y + z) \approx (x + y) + z \quad x + 0 \approx x \quad 0 + x \approx x \quad x + -x \approx 0 \quad -x + x \approx 0.$$

Similarly, lattices are the algebras of type  $\rho_L$  that satisfy the equations

$$\begin{array}{lll} x \wedge x \approx x & x \vee x \approx x & (\text{idempotent laws}) \\ x \wedge y \approx y \wedge x & x \vee y \approx y \vee x & (\text{commutative laws}) \\ x \wedge (y \wedge z) \approx (x \wedge y) \wedge z & x \vee (y \vee z) \approx (x \vee y) \vee z & (\text{associative laws}) \\ x \wedge (y \vee x) \approx x & x \vee (y \wedge x) \approx x. & (\text{absorption laws}) \end{array}$$

## 2. BASIC CONSTRUCTIONS

Algebras of the same type are called *similar* and can be compared by means of maps that preserve their structure.

**Definition 2.1.** Given two similar algebras  $A$  and  $B$ , a *homomorphism* from  $A$  to  $B$  is a map  $f: A \rightarrow B$  such that, for every  $n$ -ary operation  $g$  of the common type and  $a_1, \dots, a_n \in A$ ,

$$f(g^A(a_1, \dots, a_n)) = g^B(f(a_1), \dots, f(a_n)).$$

An injective homomorphism is called an *embedding* and, if there exists an embedding from  $A$  to  $B$ , we say that  $A$  *embeds* into  $B$ . Lastly, a surjective embedding is called an *isomorphism*. Accordingly,  $A$  and  $B$  are said to be *isomorphic* if there exists an isomorphism between them, in which case we write  $A \cong B$ .

A simple induction on the construction of terms shows that, for every pair of algebras  $A$  and  $B$  of type  $\rho$  and every term  $\varphi(x_1, \dots, x_n)$  of  $\rho$ , if  $f$  is a homomorphism from  $A$  to  $B$ , then

$$f(\varphi^A(a_1, \dots, a_n)) = \varphi^B(f(a_1), \dots, f(a_n)),$$

for every  $a_1, \dots, a_n \in A$ . Therefore homomorphisms preserve not only basic operations, but also arbitrary terms.

In the particular case where  $A$  and  $B$  are lattices, a homomorphism from  $A$  to  $B$  is a map  $f: A \rightarrow B$  such that, for every  $a, c \in A$ ,

$$f(a \wedge^A c) = f(a) \wedge^B f(c) \quad \text{and} \quad f(a \vee^A c) = f(a) \vee^B f(c).$$

For instance, the inclusion map from the lattice  $\langle \mathbb{N}; \leq \rangle$  into the lattice  $\langle \mathbb{Z}; \leq \rangle$  is an injective homomorphism, that is, an embedding. Similarly, given two sets  $Y \subseteq X$ , the inclusion map from the powerset lattice  $\langle \mathcal{P}(Y); \subseteq \rangle$  to the powerset lattice  $\langle \mathcal{P}(X); \subseteq \rangle$  is also an embedding. On the other hand, if  $Y \subsetneq X$ , the map

$$(-) \cap Y: \mathcal{P}(X) \rightarrow \mathcal{P}(Y)$$

that sends every  $Z \subseteq X$  to  $Z \cap Y$  is a noninjective homomorphism from  $\langle \mathcal{P}(X); \subseteq \rangle$  to  $\langle \mathcal{P}(Y); \subseteq \rangle$ .

**Definition 2.2.** Let  $A$  and  $B$  be algebras of the same type  $\rho: \mathcal{F} \rightarrow \mathbb{N}$ . Then  $A$  is said to be a *subalgebra* of  $B$  if  $A \subseteq B$  and  $f^A$  is the restriction of  $f^B$  to  $A$ , for every  $f \in \mathcal{F}$ . In this case, we write  $A \leq B$ .

Given a class of algebras  $K$ , let

$$\mathbb{I}(K) := \{A : A \cong B \text{ for some } B \in K\}$$

$$\mathbb{S}(K) := \{A : A \leq B \text{ for some } B \in K\}.$$

When  $K = \{A\}$ , we write  $\mathbb{I}(A)$  and  $\mathbb{S}(A)$  as a shorthand for  $\mathbb{I}(\{A\})$  and  $\mathbb{S}(\{A\})$ , respectively. The following observation is an immediate consequence of the definitions.

**Proposition 2.3.** Let  $A$  and  $B$  be algebras of the same type. Then  $A \in \mathbb{IS}(B)$  if and only if there exists an embedding  $f: A \rightarrow B$ . In this case,  $A$  is isomorphic to the unique subalgebra of  $B$  with universe  $f[A]$ .

As we mentioned, homomorphisms can be used to compare similar algebras.

**Definition 2.4.** Given two similar algebras  $A$  and  $B$ , we say that  $A$  is a *homomorphic image* of  $B$  if there exists a surjective homomorphism  $f: B \rightarrow A$ .

Accordingly, given a class of algebras  $K$ , we set

$$\mathbb{H}(K) := \{A : A \text{ is a homomorphic image of some } B \in K\}.$$

As usual, when  $K = \{A\}$ , we write  $\mathbb{H}(A)$  as a shorthand for  $\mathbb{H}(\{A\})$ .

Observe that every (not necessarily surjective) homomorphism  $f: A \rightarrow B$  induces a homomorphic image of  $A$ .

**Proposition 2.5.** *If  $f: A \rightarrow B$  is a homomorphism, then  $f[A]$  is the universe of a subalgebra of  $B$  that, moreover, is a homomorphic image of  $A$ .*

*Proof.* Observe that  $f[A]$  is nonempty, because  $A$  is. Then consider an  $n$ -ary function symbol  $g$  of the common type of  $A$  and  $B$  and  $b_1, \dots, b_n \in f[A]$ . Clearly, there are  $a_1, \dots, a_n \in A$  such that  $f(a_i) = b_i$ , for every  $i \leq n$ . Since  $f$  is a homomorphism from  $A$  to  $B$ , we obtain

$$g^B(b_1, \dots, b_n) = g^B(f(a_1), \dots, f(a_n)) = f(g^A(a_1, \dots, a_n)) \in f[A].$$

Hence, we conclude that  $f[A]$  is the universe of a subalgebra  $f[A]$  of  $B$ .

Furthermore,  $f: A \rightarrow f[A]$  is a homomorphism, because for every basic  $n$ -ary function symbol  $g$  of the common type and  $a_1, \dots, a_n \in A$ ,

$$f(g^A(a_1, \dots, a_n)) = g^B(f(a_1), \dots, f(a_n)) = g^{f[A]}(f(a_1), \dots, f(a_n)),$$

where the first equality follows from the assumption that  $f: A \rightarrow B$  is a homomorphism. Since the map  $f: A \rightarrow f[A]$  is surjective, we conclude that  $f[A] \in \mathbb{H}(A)$ .  $\square$

In view of the above result, when  $f: A \rightarrow B$  is a homomorphism, we denote by  $f[A]$  the unique subalgebra of  $B$  with universe  $f[A]$ .

For instance, let  $f: \mathbb{Z} \rightarrow \mathbb{R}$  be the absolute value map, that is, the function defined by the rule

$$f(n) := \text{the absolute value of } n.$$

Observe that  $f$  is a nonsurjective homomorphism from the lattice of integers to that of reals. Furthermore, the homomorphic image  $f[\langle \mathbb{Z}; \leq \rangle]$  of  $\langle \mathbb{Z}; \leq \rangle$  is the lattice of natural numbers  $\langle \mathbb{N}; \leq \rangle$ , which, in turn, is a subalgebra of lattice of reals.

Notably, the homomorphic images of an algebra  $A$  can be “internalized” as special equivalence relations on  $A$  as follows.

**Definition 2.6.** A *congruence* of an algebra  $A$  is an equivalence relation  $\theta$  on  $A$  such that, for every basic  $n$ -ary operation  $f$  of  $A$  and  $a_1, \dots, a_n, c_1, \dots, c_n \in A$ ,

$$\text{if } \langle a_1, c_1 \rangle, \dots, \langle a_n, c_n \rangle \in \theta, \text{ then } \langle f^A(a_1, \dots, a_n), f^A(c_1, \dots, c_n) \rangle \in \theta. \quad (1)$$

In this case, we often write  $a \equiv_\theta c$  as a shorthand for  $\langle a, c \rangle \in \theta$ . The poset of congruences of  $A$  ordered under the inclusion relation will be denoted by  $\text{Con}(A)$ .

A simple induction on the construction of terms shows that, for every congruence  $\theta$  of  $A$  and every term  $\varphi(x_1, \dots, x_n)$ ,

$$\text{if } \langle a_1, c_1 \rangle, \dots, \langle a_n, c_n \rangle \in \theta, \text{ then } \langle \varphi^A(a_1, \dots, a_n), \varphi^A(c_1, \dots, c_n) \rangle \in \theta,$$

for every  $a_1, \dots, a_n \in A$ . Therefore, congruences preserve not only basic operations, but also arbitrary terms. Furthermore, a simple argument shows that  $\text{Con}(A)$  is an inductive closure system and, therefore, an algebraic lattice whose maximum is the total relation  $A \times A$  and whose minimum is the identity relation  $\text{id}_A := \{ \langle a, a \rangle : a \in A \}$ .

**Example 2.7** (Boolean algebras). Recall that a *filter* of a Boolean algebra  $A$  is a nonempty upset  $F \subseteq A$  closed under binary meets. We denote by  $\text{Fi}(A)$  the poset of filters of  $A$  ordered under the inclusion relation. It is easy to see  $\text{Fi}(A)$  is an inductive closure system and, therefore, an algebraic lattice. Furthermore, the lattices  $\text{Fi}(A)$  and  $\text{Con}(A)$  are isomorphic via the inverse isomorphisms

$$\Omega^A(-): \text{Fi}(A) \rightarrow \text{Con}(A) \quad \text{and} \quad \tau(-): \text{Con}(A) \rightarrow \text{Fi}(A)$$

defined by the rules

$$\begin{aligned} \Omega^A(F) &:= \{ \langle a, c \rangle \in A \times A : a \rightarrow c, c \rightarrow a \in F \} \\ \tau(\theta) &:= \{ a \in A : \langle a, 1 \rangle \in \theta \}. \end{aligned}$$

Because of this, every congruence  $\theta$  of a Boolean algebra  $A$  is induced by some filter  $F$ , in the sense that  $\theta = \Omega^A F$ . This correspondence between filters and congruences generalizes straightforwardly to all Heyting algebras.  $\square$

**Example 2.8** (Modal algebras). A *modal algebra* is an algebra  $A = \langle A; \wedge, \vee, \neg, \Box, 0, 1 \rangle$  such that  $\langle A; \wedge, \vee, \neg, 0, 1 \rangle$  is a Boolean algebra and  $\Box$  is a unary operation such that

$$\Box(a \wedge c) = \Box a \wedge \Box c \quad \text{and} \quad \Box 1 = 1,$$

for every  $a, c \in A$ . An *open filter* of a modal algebra  $A$  is a filter of the Boolean reduct of  $A$  that, moreover, is closed under the operation  $\Box$ . The poset of open filters of  $A$  ordered under the inclusion relation will be denoted by  $\text{Op}(A)$ . It forms an inductive closure system and, therefore, an algebraic lattice. Furthermore, the lattices  $\text{Op}(A)$  and  $\text{Con}(A)$  are isomorphic via the inverse isomorphisms described in Example 2.7. Because of this, every congruence of a modal algebra  $A$  has the form  $\theta = \Omega^A F$ , for some open filter  $F$ .  $\square$

**Example 2.9** (Groups). Similarly, it is well known that the lattice of congruences of a group is isomorphic to that of its normal subgroups. Because of this, every congruence of a group is induced by some normal subgroup.  $\square$

As we mentioned, there is a tight correspondence between the homomorphic images and the congruences of an algebra  $A$ . On the one hand, every congruence  $\theta$  of  $A$  gives rise to a homomorphic image  $A/\theta$  of  $A$ . Let  $\mathcal{F}$  be the set of function symbols of  $A$ . Given  $\theta \in \text{Con}(A)$  and a basic  $n$ -ary function symbol  $f \in \mathcal{F}$ , let  $f^{A/\theta}$  be the  $n$ -ary operation on  $A/\theta$  defined by the rule

$$f^{A/\theta}(a_1/\theta, \dots, a_n/\theta) := f^A(a_1, \dots, a_n)/\theta.$$

Notice that  $f^{A/\theta}$  is well-defined, by condition (1). As a consequence, the structure

$$A/\theta := \langle A/\theta; \{f^{A/\theta} : f \in \mathcal{F}\} \rangle$$

is a well-defined algebra of the type as  $A$ . Furthermore,  $A/\theta \in \mathbb{H}(A)$ , because the map  $\pi_\theta: A \rightarrow A/\theta$ , defined, for every  $a \in A$ , as  $\pi_\theta(a) := a/\theta$ , is a surjective homomorphism from  $A$  to  $A/\theta$ . To prove this, consider  $a_1, \dots, a_n \in A$ . We have

$$\begin{aligned} \pi_\theta(f^A(a_1, \dots, a_n)) &= f^A(a_1, \dots, a_n)/\theta \\ &= f^{A/\theta}(a_1/\theta, \dots, a_n/\theta) \\ &= f^{A/\theta}(\pi_\theta(a_1), \dots, \pi_\theta(a_n)), \end{aligned}$$

where the second equality follows from the definition of the operation  $f^{A/\theta}$ .

**Corollary 2.10.** *If  $\theta$  is a congruence of an algebra  $A$ , then  $A/\theta$  is a well-defined homomorphic image of  $A$ .*

In view of the above result, every congruence  $\theta$  of an algebra  $A$  induces a homomorphic image of  $A$ , namely  $A/\theta$ . The converse is also true, as we proceed to explain.

**Definition 2.11.** The *kernel* of a homomorphism  $f: A \rightarrow B$  is the binary relation

$$\text{Ker}(f) := \{\langle a, c \rangle \in A \times A : f(a) = f(c)\}.$$

**Proposition 2.12.** *The kernel of a homomorphism  $f: A \rightarrow B$  is a congruence of  $A$ .*

*Proof.* It is obvious that  $\text{Ker}(f)$  is an equivalence relation on  $A$ . Therefore, to prove that  $\text{Ker}(f)$  is a congruence of  $A$ , it suffices to show that it preserves the basic operations of  $A$ . Consider a basic  $n$ -ary operation  $g$  of  $A$  and  $a_1, \dots, a_n, c_1, \dots, c_n \in A$  such that  $\langle a_1, c_1 \rangle, \dots, \langle a_n, c_n \rangle \in \text{Ker}(f)$ . By the definition of  $\text{Ker}(f)$ ,

$$f(a_i) = f(c_i), \text{ for every } i \leq n.$$

It follows that  $g^B(f(a_1), \dots, f(a_n)) = g^B(f(c_1), \dots, f(c_n))$ . Since  $f: A \rightarrow B$  is a homomorphism, this yields

$$f(g^A(a_1, \dots, a_n)) = g^B(f(a_1), \dots, f(a_n)) = g^B(f(c_1), \dots, f(c_n)) = f(g^A(c_1, \dots, c_n)).$$

Hence, we conclude that  $\langle g^A(a_1, \dots, a_n), g^A(c_1, \dots, c_n) \rangle \in \text{Ker}(f)$ , as desired.  $\square$

The behaviour of kernels is governed by the next principle.

**Fundamental Homomorphism Theorem 2.13.** *If  $f: A \rightarrow B$  is a homomorphism with kernel  $\theta$ , then there exists a unique embedding  $g: A/\theta \rightarrow B$  such that  $f = g \circ \pi_\theta$ .*

*Proof.* We begin by proving the existence of  $g$ . Let  $g: A/\theta \rightarrow B$  be the map defined as  $g(a/\theta) := f(a)$ , for every  $a \in A$ . To show that  $g$  is well-defined, consider  $a, c \in A$  such that  $a/\theta = c/\theta$ . Since  $\theta = \text{Ker}(f)$ , this means that  $f(a) = f(c)$ , as desired. Furthermore, the definition of  $g$  guarantees that  $f = g \circ \pi_\theta$ .

Now, observe  $g$  is injective, because, for every  $a, c \in A$  such that  $g(a/\theta) = g(c/\theta)$ , we have  $f(a) = f(c)$ , that is,  $\langle a, c \rangle \in \text{Ker}(f) = \theta$  and, therefore,  $a/\theta = c/\theta$ . Moreover, for every basic  $n$ -ary operation  $p$  of  $A$  and  $a_1, \dots, a_n \in A$ , we have

$$\begin{aligned} g(p^{A/\theta}(a_1/\theta, \dots, a_n/\theta)) &= g(p^A(a_1, \dots, a_n)/\theta) \\ &= f(p^A(a_1, \dots, a_n)) \\ &= p^B(f(a_1), \dots, f(a_n)) \\ &= p^B(g(a_1/\theta), \dots, g(a_n/\theta)). \end{aligned}$$

The first equality above follows from the definition of  $A/\theta$ , the second and the last from the definition of  $g$ , and the third from the assumption that  $f: A \rightarrow B$  is a homomorphism. Hence, we conclude that  $g: A/\theta \rightarrow B$  is a homomorphism and, therefore, an embedding, as desired.

The uniqueness of  $g$  follows from the fact that, if a map  $g^*$  satisfies the condition in the statement of the theorem, then, for every  $a \in A$ ,

$$f(a) = g^* \circ \pi_\theta(a) = g^*(a/\theta),$$

that is,  $g^*$  coincides with  $g$ .  $\square$

**Corollary 2.14.** *If  $f: A \rightarrow B$  is a homomorphism, then  $f[A] \cong A/\text{Ker}(f)$ . In particular, if  $f$  is surjective,  $B \cong A/\text{Ker}(f)$ .*

*Proof.* In the proof of the Fundamental Homomorphism Theorem we showed that the map  $g: A/\text{Ker}(f) \rightarrow B$ , defined by the rule  $g(a/\text{Ker}(f)) := f(a)$ , is an embedding of  $A/\text{Ker}(f)$  into  $B$ . As  $g$  can be viewed as a surjective embedding of  $A/\text{Ker}(f)$  into  $f[A]$ , we conclude that  $f[A] \cong A/\text{Ker}(f)$ .  $\square$

At this stage, it should be clear that if  $\theta$  is a congruence on an algebra  $A$ , then  $\pi_\theta: A \rightarrow A/\theta$  is a surjective homomorphism whose kernel is  $\theta$ . Similarly, if  $f: A \rightarrow B$  is a surjective homomorphism, then  $A/\text{Ker}(f) \cong B$ , by Corollary 2.14. As a consequence, for every class of algebras  $K$ ,

$$\mathbb{H}(K) = \mathbb{I}\{A/\theta : A \in K \text{ and } \theta \in \text{Con}(A)\}. \quad (2)$$

Now, recall that the Cartesian product of a family of sets  $\{A_i : i \in I\}$  is the set

$$\prod_{i \in I} A_i := \{f: I \rightarrow \bigcup_{i \in I} A_i : f(i) \in A_i, \text{ for all } i \in I\}.$$

In particular, if  $I$  is empty, then  $\prod_{i \in I} A_i$  is the singleton containing only the empty map.

**Definition 2.15.** The *direct product* of a family of similar algebras  $\{A_i : i \in I\}$  is the unique algebra of the common type whose universe is the Cartesian product  $\prod_{i \in I} A_i$  and such that, for every basic  $n$ -ary operation symbol  $f$  and every  $\vec{a}_1, \dots, \vec{a}_n \in \prod_{i \in I} A_i$ ,

$$f^{\prod_{i \in I} A_i}(\vec{a}_1, \dots, \vec{a}_n)(i) = f^{A_i}(\vec{a}_1(i), \dots, \vec{a}_n(i)), \text{ for every } i \in I.$$

We denote this algebra by  $\prod_{i \in I} A_i$ .

In this case, for every  $j \in I$ , the projection map on the  $j$ -th component  $p_j: \prod_{i \in I} A_i \rightarrow A_j$ , defined by the rule  $p_j(\vec{a}) := \vec{a}(j)$ , is a surjective homomorphism from  $\prod_{i \in I} A_i$  to  $A_j$ .

Given a class of similar algebras  $K$ , we set

$$\mathbb{P}(K) := \{A : A \text{ is a direct product of a family } \{B_i : i \in I\} \subseteq K\}.$$

As usual, when  $K = \{A\}$ , we write  $\mathbb{P}(A)$  as a shorthand for  $\mathbb{P}(\{A\})$ .

Notice that up to isomorphism, there exists a unique one-element algebra of a given type. Because of this, one-element algebras are called *trivial*. Accordingly, when the set of indexes  $I$  is empty, the direct product  $\prod_{i \in I} A_i$  is the trivial algebra of the given type. It follows that  $\mathbb{P}(K)$  contains always a trivial algebra, for every class of similar algebras  $K$ .

**Example 2.16** (Powerset algebras). Boolean algebras of the form  $\langle \mathcal{P}(X); \cap, \cup, -, \emptyset, X \rangle$  are called *powerset Boolean algebras*. Let  $B$  be the two-element Boolean algebra and observe that  $\mathbb{I}\mathbb{P}(B)$  is the class of algebras isomorphic to some powerset Boolean algebra. To prove this, observe that every powerset Boolean algebra  $\mathcal{P}(X)$  is isomorphic to a direct product of  $B$  via the *characteristic function*  $f_X: \mathcal{P}(X) \rightarrow \prod_{x \in X} B_x$ , defined by the rule

$$f(Y)(x) := \begin{cases} 1 & \text{if } x \in Y \\ 0 & \text{if } x \notin Y, \end{cases}$$

where  $Y \in \mathcal{P}(X)$  and  $x \in X$ . By the same token, every direct product  $\prod_{i \in I} B_i$  of  $B$  is isomorphic to the powerset Boolean algebra  $\mathcal{P}(I)$  via the isomorphism  $f_I$ .  $\square$

We close this section by reviewing the subdirect product construction.

**Definition 2.17.** A subalgebra  $B$  of a direct product  $\prod_{i \in I} A_i$  is said to be a *subdirect product* of  $\{A_i : i \in I\}$  if the projection map  $p_i$  is surjective, for every  $i \in I$ . Similarly, an embedding  $f: B \rightarrow \prod_{i \in I} A_i$  is said to be *subdirect* when  $f[B]$  is a subdirect product of the family  $\{A_i : i \in I\}$ .

Given a class of similar algebras  $K$ , we set

$$\mathbb{P}_{\text{SD}}(K) := \{A : A \text{ is a subdirect direct product of a family } \{B_i : i \in I\} \subseteq K\}.$$

As usual, when  $K = \{A\}$ , we write  $\mathbb{P}_{\text{SD}}(A)$  as a shorthand for  $\mathbb{P}_{\text{SD}}(\{A\})$ . Clearly,  $\mathbb{P}_{\text{SD}}(K) \subseteq \mathbb{SP}(K)$ . Furthermore,  $\mathbb{P}_{\text{SD}}(K)$  contains always a trivial algebra.

**Example 2.18** (Distributive lattices). Let  $\text{DL}$  be the class of distributive lattices and  $B$  be the two-element distributive lattice. Birkhoff's Representation Theorem states that  $\text{DL} = \mathbb{IP}_{\text{SD}}(B)$ . The inclusion  $\mathbb{IP}_{\text{SD}}(B) \subseteq \text{DL}$  follows from the fact that  $\text{DL}$  is closed under  $\mathbb{I}, \mathbb{S}$ , and  $\mathbb{P}$ . For the other inclusion, consider a distributive lattice  $A$  and let  $I$  be the set of its prime filters. By Birkhoff's Representation Theorem, the map

$$\gamma: A \rightarrow \prod_{F \in I} B_F,$$

defined, for every  $a \in A$  and  $F \in I$ , by the rule

$$\gamma(a)(F) := \begin{cases} 1 & \text{if } a \in F \\ 0 & \text{if } a \notin F, \end{cases}$$

is a well-defined subdirect embedding. \(\square\)

**Example 2.19** (Boolean algebras). Similarly, Stone's Representation Theorem states that the class of Boolean algebras coincides with  $\mathbb{IP}_{\text{SD}}(B)$ , where  $B$  the two-element Boolean algebra. \(\square\)

The next result provides a general recipe to construct subdirect products.

**Proposition 2.20.** Let  $A$  be an algebra and  $\{\theta_i : i \in I\} \subseteq \text{Con}(A)$ . Then the map

$$f: A / \bigcap_{i \in I} \theta_i \rightarrow \prod_{i \in I} A / \theta_i,$$

defined, for every  $a \in A$  and  $j \in I$ , as

$$f(a / \bigcap_{i \in I} \theta_i)(j) := a / \theta_j,$$

is a subdirect embedding.

*Proof.* For the sake of readability, set  $B := A / \bigcap_{i \in I} \theta_i$ . To prove that  $f$  is injective, consider  $a, c \in A$  such that  $\langle a, c \rangle \notin \bigcap_{i \in I} \theta_i$ . Then there exists  $j \in I$  such that  $\langle a, c \rangle \notin \theta_j$  and, therefore,

$$f(a / \bigcap_{i \in I} \theta_i)(j) := a / \theta_j \neq c / \theta_j = f(c / \bigcap_{i \in I} \theta_i)(j).$$

It follows that  $f(a / \bigcap_{i \in I} \theta_i) \neq f(c / \bigcap_{i \in I} \theta_i)$ . Thus,  $f$  is injective. Moreover, by the definition of  $f$ , the composition  $p_i \circ f: B \rightarrow A / \theta_i$  is surjective, for every  $i \in I$ .



It only remains to prove that  $f$  is a homomorphism. Consider an  $n$ -ary basic operation  $g$  and  $a_1, \dots, a_n \in A$ . For every  $j \in I$ , we have

$$\begin{aligned}
 f(g^B(a_1 / \bigcap_{i \in I} \theta_i, \dots, a_n / \bigcap_{i \in I} \theta_i))(j) &= f(g^A(a_1, \dots, a_n) / \bigcap_{i \in I} \theta_i)(j) \\
 &= g^A(a_1, \dots, a_n) / \theta_j \\
 &= g^{A/\theta_j}(a_1 / \theta_j, \dots, a_n / \theta_j) \\
 &= g^{A/\theta_j}(f(a_1 / \bigcap_{i \in I} \theta_i)(j), \dots, f(a_n / \bigcap_{i \in I} \theta_i)(j)) \\
 &= g^{\prod_{i \in I} A/\theta_i}(f(a_1 / \bigcap_{i \in I} \theta_i), \dots, f(a_n / \bigcap_{i \in I} \theta_i))(j).
 \end{aligned}$$

It follows that

$$f(g^B(a_1 / \bigcap_{i \in I} \theta_i, \dots, a_n / \bigcap_{i \in I} \theta_i)) = g^{\prod_{i \in I} A/\theta_i}(f(a_1 / \bigcap_{i \in I} \theta_i), \dots, f(a_n / \bigcap_{i \in I} \theta_i)). \quad \square$$

### 3. PREVARIETIES

**Definition 3.1.** A class of similar algebras closed  $\mathbb{I}$ ,  $\mathbb{S}$ , and  $\mathbb{P}$  is said to be a *prevariety*.

Given a class of similar algebras  $K$ , the least prevariety extending  $K$  is  $\mathbb{ISP}(K)$  and is called the prevariety *generated* by  $K$ . Our aim will be to prove that prevarieties are precisely the classes of algebras axiomatized by a certain kind of infinitary formulas.

When no confusion shall arise, given a sequence  $\vec{a}$  and a set  $A$ , we write  $\vec{a} \in A$  to indicate that the elements of the sequence  $\vec{a}$  belong to  $A$ .

**Definition 3.2.** A *generalized quasi-equation* of type  $\rho$  is an expression  $\Phi$  of the form

$$\left( \bigwedge_{i \in I} \varphi_i(\vec{x}) \approx \psi_i(\vec{x}) \right) \implies \varepsilon(\vec{x}) \approx \delta(\vec{x}),$$

where  $\{\varphi_i \approx \psi_i : i \in I\} \cup \{\varepsilon \approx \delta\}$  is a set of equations of type  $\rho$ . Then  $\Phi$  is *valid* in an algebra  $A$  of type  $\rho$  when so is its universal closure, that is, for every  $\vec{a} \in A$ ,

$$\text{if } (\varphi_i^A(\vec{a}) = \psi_i^A(\vec{a}), \text{ for all } i \in I), \text{ then } \varepsilon^A(\vec{a}) = \delta^A(\vec{a}).$$

In this case, we often say that  $A$  *satisfies*  $\Phi$ .

Notice that, in the above definition, the set of indexes  $I$  can be arbitrarily large and that the same applies to the sequence of variables  $\vec{x}$  that appear in the equations of  $\Phi$ . This motivates the following.

**Definition 3.3.** A generalized quasi-equation is said to be

- (i) a *quasi-equation* when the index set  $I$  is finite; and
- (ii) an *equation* when the index set  $I$  is empty.

*Remark 3.4.* It might seem that we are using the term *equations* to refer to two distinct kinds of expressions, namely those of the form  $\varepsilon \approx \delta$  and  $\emptyset \implies \varepsilon \approx \delta$ . This is not a problem, however, because these expressions are synonyms, in the sense that an algebra satisfies  $\varepsilon \approx \delta$  if and only if it satisfies  $\emptyset \implies \varepsilon \approx \delta$ . Because of this, we will continue to denote equations by  $\varepsilon \approx \delta$ , while keeping in mind that they are special instances of generalized quasi-equations.  $\square$

**Definition 3.5.** Let  $\rho: \mathcal{F} \rightarrow \mathbb{N}$  be a type and  $X$  a set of variables disjoint from  $\mathcal{F}$ . The *term algebra*  $T_\rho(X)$  of type  $\rho$  over  $X$  is the unique algebra of type  $\rho$  whose universe is  $T_\rho(X)$  and with basic  $n$ -ary operations  $f$  defined, for every  $\varphi_1, \dots, \varphi_n \in T_\rho(X)$ , as

$$f^{T_\rho(X)}(\varphi_1, \dots, \varphi_n) := f(\varphi_1, \dots, \varphi_n).$$

Term algebras have the following fundamental property.

**Proposition 3.6.** Let  $A$  be an algebra of type  $\rho$  and  $X$  a set of variables. Every function  $f: X \rightarrow A$  extends uniquely to a homomorphism  $f^*: T_\rho(X) \rightarrow A$ .

*Proof.* The unique extension  $f^*$  is defined, for every  $\varphi(x_{\alpha_1}, \dots, x_{\alpha_n}) \in T_\rho(X)$ , as

$$f^*(\varphi) = \varphi^A(f(x_{\alpha_1}), \dots, f(x_{\alpha_n})). \quad \square$$

*Exercise 3.7.* Prove the above proposition.  $\square$

**Theorem 3.8.** A class of similar algebras is a prevariety if and only if it can be axiomatized by a class of generalized quasi-equations.

*Proof.* The “if” part follows from the fact that the validity of generalized quasi-equations persists under the formation of isomorphic copies, subalgebras, and direct product. To prove the converse, consider a prevariety  $K$  and let  $\Sigma$  be the proper class of generalized quasi-equations valid in it. Let  $K^+$  be the class of algebras in which the generalized quasi-equations in  $\Sigma$  are valid. Clearly,  $K \subseteq K^+$ . To prove the other inclusion, consider an algebra  $A \in K^+$ . Let also  $X$  be a set of variables for which there exists a surjective map  $f: X \rightarrow A$ . By Proposition 3.6,  $f$  extends to a surjective homomorphism  $f^*: T_\rho(X) \rightarrow A$ . Together with Corollary 2.14, this yields

$$A \cong T_\rho(X) / \text{Ker}(f^*). \quad (3)$$

Now, consider an arbitrary pair  $\langle \varphi, \psi \rangle \in (T_\rho(X) \times T_\rho(X)) \setminus \text{Ker}(f^*)$ . Notice that the elements of  $T_\rho(X) \times T_\rho(X)$  are ordered pairs of terms and, therefore, can be viewed as equations under the identification of  $\langle \varepsilon, \delta \rangle$  with  $\varepsilon \approx \delta$ . In this way,  $\text{Ker}(f^*)$  becomes a set of equations in variables  $X$ . Bearing this in mind, consider the generalized quasi-equation

$$\Phi := \left( \& \text{Ker}(f^*) \right) \implies \varphi \approx \psi.$$

We will prove that  $\Phi$  fails in  $A$ . For the sake of readability we will denote by  $\vec{x}$  the sequence of all variables in  $X$ . Observe that every element  $\varepsilon \in T_\rho(X)$  is of the form  $\varepsilon(\vec{x})$ . Then consider the assignment  $f: X \rightarrow A$ . We will denote by  $f(\vec{x})$  the sequence obtained by applying  $f$  component-wise to  $\vec{x}$ . For every pair  $\langle \varepsilon, \delta \rangle \in \text{Ker}(f^*)$ , we have

$$\varepsilon^A(f(\vec{x})) = \varepsilon^A(f^*(\vec{x})) = f^*(\varepsilon(\vec{x})) = f^*(\delta(\vec{x})) = \delta^A(f^*(\vec{x})) = \delta^A(f(\vec{x})).$$

The equalities above can be justified as follows. The first and the last holds because  $f^*$  extends  $f$ , the second and the fourth because  $f^*: T_\rho(X) \rightarrow A$  is a homomorphism, and the third because  $\langle \varepsilon, \delta \rangle \in \text{Ker}(f^*)$ . On the other hand, since  $\langle \varphi, \psi \rangle \notin \text{Ker}(f^*)$ , a similar argument shows

$$\varphi^A(f(\vec{x})) \neq \psi^A(f(\vec{x})).$$

Thus,  $A$  refutes  $\Phi$ , as desired.

Since  $A \in K^+$ , this implies that there exists some algebra  $C_{\varphi,\psi} \in K$  and an assignment  $g_{\varphi,\psi}: X \rightarrow C_{\varphi,\psi}$  such that

$$\varepsilon^{C_{\varphi,\psi}}(g_{\varphi,\psi}(\vec{x})) = \delta^{C_{\varphi,\psi}}(g_{\varphi,\psi}(\vec{x})), \text{ for all } \langle \varepsilon, \delta \rangle \in \text{Ker}(f^*), \text{ and } \varphi^{C_{\varphi,\psi}}(g_{\varphi,\psi}(\vec{x})) \neq \psi^{C_{\varphi,\psi}}(g_{\varphi,\psi}(\vec{x})).$$

Recall that  $g_{\varphi,\psi}$  extends uniquely to a homomorphism  $g_{\varphi,\psi}^*: T_\rho(X) \rightarrow C_{\varphi,\psi}$ . Moreover, from the above display it follows

$$g_{\varphi,\psi}^*(\varepsilon) = g_{\varphi,\psi}^*(\delta), \text{ for all } \langle \varepsilon, \delta \rangle \in \text{Ker}(f^*), \text{ and } g_{\varphi,\psi}^*(\varphi) \neq g_{\varphi,\psi}^*(\psi).$$

Consequently,

$$\text{Ker}(f^*) \subseteq \text{Ker}(g_{\varphi,\psi}^*) \text{ and } \langle \varphi, \psi \rangle \notin \text{Ker}(g_{\varphi,\psi}^*).$$

It follows that

$$\text{Ker}(f^*) = \bigcap \{ \text{Ker}(g_{\varphi,\psi}^*) : \langle \varphi, \psi \rangle \in (T_\rho(X) \times T_\rho(X)) \setminus \text{Ker}(f^*) \}.$$

By Proposition 2.20, this yields

$$T_\rho(X)/\text{Ker}(f^*) \in \mathbb{IP}_{\text{SD}}(\{T_\rho(X)/\text{Ker}(g_{\varphi,\psi}^*) : \langle \varphi, \psi \rangle \in (T_\rho(X) \times T_\rho(X)) \setminus \text{Ker}(f^*)\}). \quad (4)$$

Moreover, from Corollary 2.14 and the fact that  $K$  is closed under  $\mathbb{I}$  and  $\mathbb{S}$  it follows that

$$T_\rho(X)/\text{Ker}(g_{\varphi,\psi}^*) \in \mathbb{IS}(C_{\varphi,\psi}) \subseteq K,$$

for every  $\langle \varphi, \psi \rangle \in (T_\rho(X) \times T_\rho(X)) \setminus \text{Ker}(f^*)$ . Consequently, (4) simplifies to

$$T_\rho(X)/\text{Ker}(f^*) \in \mathbb{IP}_{\text{SD}}(K) \subseteq K,$$

where the last inclusion follows from the fact that  $K$  is a prevariety. Together with (3), this yields  $A \in \mathbb{I}(K) \subseteq K$ .  $\square$

*Remark 3.9.* In view of Theorem 3.8, prevarieties are classes of algebras axiomatized by classes of generalized quasi-equations. It is therefore natural to wonder whether there exists a prevariety that cannot be axiomatized by a set (as opposed to proper class) of generalized quasi-equations. It turns out that the answer to this question depends on the set theory we live in, as the nonexistence of such a prevariety is equivalent to *Vopěnka's Principle*.  $\square$

Nonetheless, prevarieties axiomatizable by a set of generalized quasi-equations admit a relatively transparent description, as we proceed to explain. Given an infinite cardinal  $\kappa$  and a class of algebras  $K$ , let

$$\mathbb{U}_\kappa(K) := \{A : B \in K, \text{ for all } \kappa\text{-generated } B \leq A\}.$$

**Definition 3.10.** Let  $\kappa$  be an infinite cardinal. A  $\kappa$ -generalized quasi-variety is a prevariety closed under  $\mathbb{U}_\kappa$ .

When  $\kappa = \aleph_0$ , we often say that  $K$  is simply a *generalized quasi-variety*. Given a class of similar algebras  $K$ , the least  $\kappa$ -generalized quasi-variety extending  $K$  is  $\mathbb{U}_\kappa \mathbb{ISP}(K)$  and is called the  $\kappa$ -generalized quasi-variety *generated* by  $K$ .

**Theorem 3.11.** Let  $\kappa$  be an infinite cardinal. A class of similar algebras is a  $\kappa$ -generalized quasi-variety if and only if it can be axiomatized by a set of generalized quasi-equations in which at most  $\kappa$  variables occur.

*Proof.* The “if” part follows from the fact that the validity of generalized quasi-equations in  $\leq \kappa$  variables persist under the  $\mathbb{I}, \mathbb{S}, \mathbb{P}, \mathbb{U}_\kappa$ . To prove the converse, consider a  $\kappa$ -generalized quasi-variety  $K$ . Then let  $X$  be a set of variables of cardinality  $\kappa$  and  $\Sigma$  the class of generalized quasi-equations written with variables in  $X$ . Since  $X$  is a set, so is  $\Sigma$ . It only remains to prove that  $K$  coincides with the class  $K^+$  of algebras satisfying the generalized quasi-equations in  $\Sigma$ . Clearly,  $K \subseteq K^+$ . To prove the other inclusion, consider an algebra  $A \in K^+$ . We need to prove that  $A \in K$ . Since  $K$  is closed under  $\mathbb{U}_\kappa$ , it suffices to show that all the  $\kappa$ -generated subalgebras of  $A$  belong to  $K$ .

Accordingly, let  $B$  be a  $\kappa$ -generated subalgebra of  $A$  and  $Y \subseteq B$  a set of generators for  $B$  of size  $\leq \kappa$ . There exists a surjective map  $f: X \rightarrow Y$ . By Proposition 3.6,  $f$  extends to a surjective homomorphism  $f^*: T_\rho(X) \rightarrow B$ . Now, we repeat the argument in the proof of Theorem 3.8, obtaining  $B \in K$ , as desired.  $\square$

**Corollary 3.12.** *A prevariety can be axiomatized by a set of generalized quasi-equations if and only if it is a  $\kappa$ -generalized quasi-variety, for some infinite cardinal  $\kappa$ .*

*Exercise\* 3.13.* Let  $K$  be a class of similar algebras and  $\kappa$  an infinite cardinal. Prove that the prevariety and the  $\kappa$ -generalized quasi-variety generated by  $K$  are, respectively,  $\mathbb{ISP}(K)$  and  $\mathbb{U}_\kappa \mathbb{ISP}(K)$ .  $\square$

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