



CS 6375 (Extra Lecture)
Lagrange Multipliers,
Duality of SVMs, and Kernel SVMs

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The Strategy So Far...

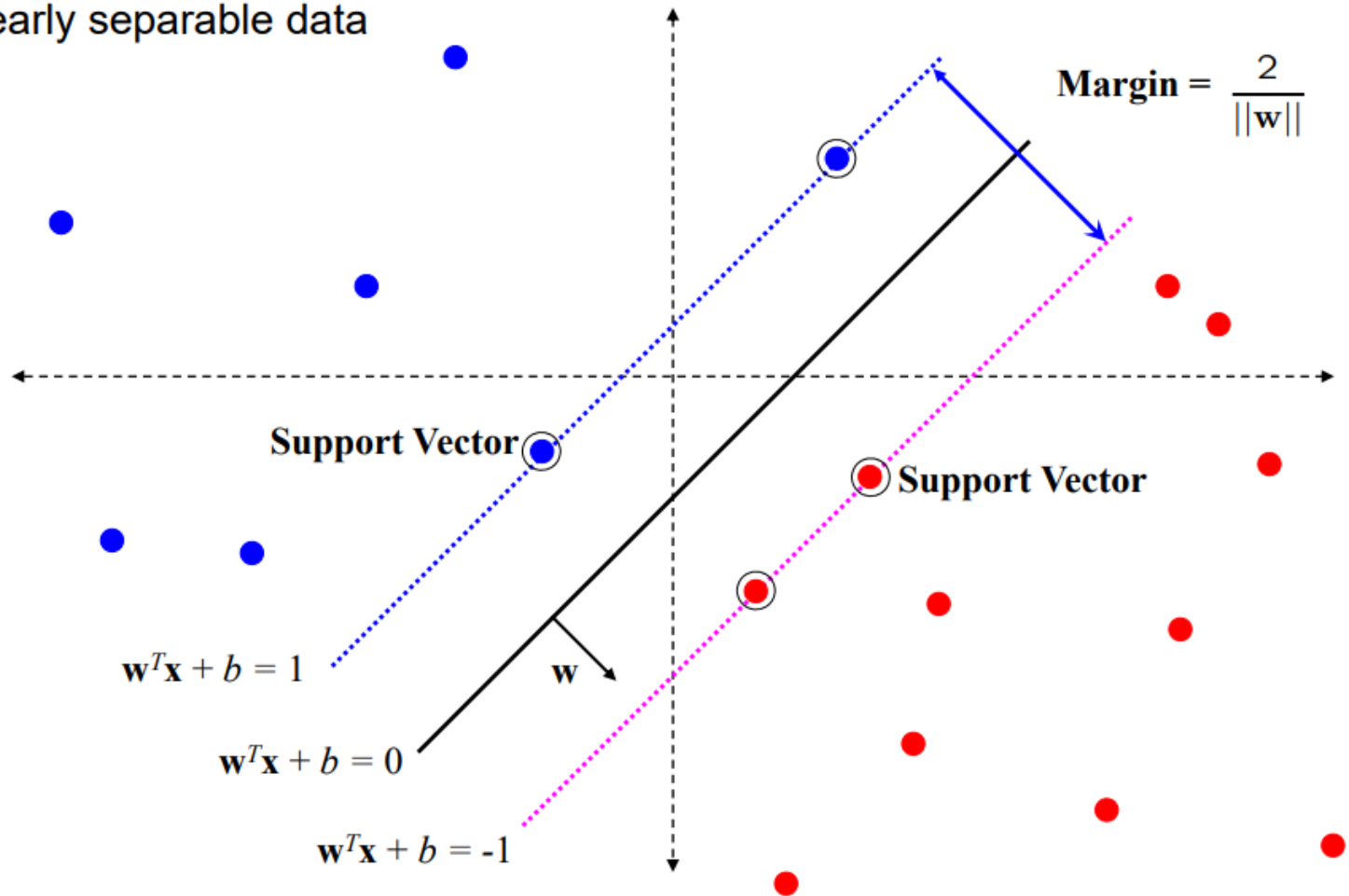


- Choose hypothesis space
- Construct loss function (ideally convex)
- Minimize loss to “learn” correct parameters

Recap: SVM



linearly separable data



- Recall: The SVM optimization problem:

$$\min_{w,b} \|w\|^2$$

such that

$$y^{(i)}(w^T x^{(i)} + b) \geq 1, \text{ for all } i$$

- This is a standard quadratic programming problem
 - Falls into the class of **convex optimization problems**
 - Can be solved with many specialized optimization tools (e.g., `quadprog()` in MATLAB)

Constrained Optimization



A mathematical detour, we'll come back to SVMs soon!

$$\min_{x \in \mathbb{R}^n} f_0(x)$$

subject to:

$$\begin{aligned} f_i(x) &\leq 0, & i &= 1, \dots, m \\ h_i(x) &= 0, & i &= 1, \dots, p \end{aligned}$$

Constrained Optimization



$$\min_{x \in \mathbb{R}^n} f_0(x)$$

f_0 is not necessarily convex

subject to:

$$\begin{aligned} f_i(x) &\leq 0, & i &= 1, \dots, m \\ h_i(x) &= 0, & i &= 1, \dots, p \end{aligned}$$

$$\min_{x \in \mathbb{R}^n} f_0(x)$$

subject to:

$$\begin{aligned} f_i(x) &\leq 0, & i &= 1, \dots, m \\ h_i(x) &= 0, & i &= 1, \dots, p \end{aligned}$$

Constraints do not need to
be linear

Example



$$\min_{x \in \mathbb{R}^3} x_1 \log x_1 + x_2 \log x_2$$

subject to:

$$x_1 + x_2 = 1$$

$$x_1 \geq 0$$

$$x_2 \geq 0$$

Example



$$\min_{x \in \mathbb{R}^3} x_1 \log x_1 + x_2 \log x_2$$

subject to:

$$1 - x_1 - x_2 = 0$$

$$-x_1 \leq 0$$

$$-x_2 \leq 0$$

$$L(x, \lambda, \nu) = f_0(x) + \sum_{i=1}^m \lambda_i f_i(x) + \sum_{i=1}^p \nu_i h_i(x)$$

- Incorporate constraints into a new objective function
- $\lambda \geq 0$ and ν are vectors of **Lagrange multipliers**
- The Lagrange multipliers can be thought of as enforcing soft constraints

Example



$$\min_{x \in \mathbb{R}^3} x_1 \log x_1 + x_2 \log x_2$$

subject to:

$$1 - x_1 - x_2 = 0$$

$$-x_1 \leq 0$$

$$-x_2 \leq 0$$

$$L(x_1, x_2, \nu_1, \lambda_1, \lambda_2)$$

$$= x_1 \log x_1 + x_2 \log x_2 + \nu_1 \cdot (1 - x_1 - x_2) - \lambda_1 x_1 - \lambda_2 x_2$$

- Construct a **dual function** by minimizing the Lagrangian over the primal variables

$$g(\lambda, \nu) = \inf_x L(x, \lambda, \nu)$$

- $g(\lambda, \nu) = -\infty$ whenever the Lagrangian is not bounded from below for a fixed λ and ν

Example



$$\min_{x \in \mathbb{R}^2} x_1 \log x_1 + x_2 \log x_2$$

subject to:

$$1 - x_1 - x_2 = 0$$

$$-x_1 \leq 0$$

$$-x_2 \leq 0$$

Example



$$\min_{x \in \mathbb{R}^3} x_1 \log x_1 + x_2 \log x_2$$

subject to:

$$1 - x_1 - x_2 = 0$$

$$-x_1 \leq 0$$

$$-x_2 \leq 0$$

$$\begin{aligned} L(x_1, x_2, \nu_1, \lambda_1, \lambda_2) \\ = x_1 \log x_1 + x_2 \log x_2 + \nu_1 \cdot (1 - x_1 - x_2) - \lambda_1 x_1 - \lambda_2 x_2 \end{aligned}$$

Example



$$\min_{x \in \mathbb{R}^3} x_1 \log x_1 + x_2 \log x_2$$

subject to:

$$1 - x_1 - x_2 = 0$$

$$-x_1 \leq 0$$

$$-x_2 \leq 0$$

The Primal Problem



$$\min_{x \in \mathbb{R}^n} f_0(x)$$

subject to:

$$\begin{aligned} f_i(x) &\leq 0, & i &= 1, \dots, m \\ h_i(x) &= 0, & i &= 1, \dots, p \end{aligned}$$

Equivalently,

$$\inf_x \sup_{\lambda \geq 0, \nu} L(x, \lambda, \nu)$$

Why are these equivalent?

$$\min_{x \in \mathbb{R}^n} f_0(x)$$

subject to:

$$\begin{aligned} f_i(x) &\leq 0, & i &= 1, \dots, m \\ h_i(x) &= 0, & i &= 1, \dots, p \end{aligned}$$

Equivalently,

$$\inf_x \sup_{\lambda \geq 0, \nu} L(x, \lambda, \nu)$$

$$\sup_{\lambda \geq 0, \nu} \left[f_0(x) + \sum_{i=1}^m \lambda_i f_i(x) + \sum_{i=1}^p \nu_i h_i(x) \right] = \infty$$

whenever x violates the constraints

$$\sup_{\lambda \geq 0, \nu} g(\lambda, \nu)$$

Equivalently,

$$\sup_{\lambda \geq 0, \nu} \inf_x L(x, \lambda, \nu)$$

- The dual problem is always concave, even if the primal problem is not convex
 - For each x , $L(x, \lambda, \nu)$ is a linear function in λ and ν
 - Minimum (or infimum) of linear functions is concave!

$$\sup_{\lambda \geq 0, \nu} \inf_x L(x, \lambda, \nu) \leq \inf_x \sup_{\lambda \geq 0, \nu} L(x, \lambda, \nu)$$

- Why?
 - $g(\lambda, \nu) \leq L(x, \lambda, \nu)$ for all x
 - $L(x', \lambda, \nu) \leq f_0(x')$ for any feasible x' , $\lambda \geq 0$
 - x is **feasible** if it satisfies all of the constraints
- Let x^* be the optimal solution to the primal problem and $\lambda \geq 0$

$$g(\lambda, \nu) \leq L(x^*, \lambda, \nu) \leq f_0(x^*)$$

Example: Solving the Dual Problem



$$\min_{x \in \mathbb{R}^3} x_1 \log x_1 + x_2 \log x_2$$

subject to:

$$1 - x_1 - x_2 = 0$$

$$-x_1 \leq 0$$

$$-x_2 \leq 0$$

More Examples



- Minimize $x^2 + y^2$ subject to $x + y \geq 1$
- Given a point $z \in \mathbb{R}^n$ and a hyperplane $w^T x + b = 0$, find the projection of the point z onto the hyperplane

- Under certain conditions, the two optimization problems are equivalent

$$\sup_{\lambda \geq 0, \nu} \inf_x L(x, \lambda, \nu) = \inf_x \sup_{\lambda \geq 0, \nu} L(x, \lambda, \nu)$$

- This is called **strong duality**
- If the inequality is strict, then we say that there is a **duality gap**
 - Size of gap measured by the difference between the two sides of the inequality

Slater's Condition



For any optimization problem of the form

$$\min_{x \in \mathbb{R}^n} f_0(x)$$

subject to:

$$\begin{aligned} f_i(x) &\leq 0, & i &= 1, \dots, m \\ Ax &= b \end{aligned}$$

where f_0, \dots, f_m are **convex functions**, strong duality holds if there exists an x such that

$$\begin{aligned} f_i(x) &< 0, & i &= 1, \dots, m \\ Ax &= b \end{aligned}$$

$$\min_w \frac{1}{2} \|w\|^2$$

such that

$$y_i(w^T x^{(i)} + b) \geq 1, \text{ for all } i$$

- Note that Slater's condition holds as long as the data is linearly separable

$$L(w, b, \lambda) = \frac{1}{2}w^T w + \sum_i \lambda_i (1 - y_i(w^T x^{(i)} + b))$$

Convex in w , so take derivatives to form the dual

$$\frac{\partial L}{\partial w_k} = w_k + \sum_i -\lambda_i y_i x_k^{(i)} = 0$$

$$\frac{\partial L}{\partial b} = \sum_i -\lambda_i y_i = 0$$

$$L(w, b, \lambda) = \frac{1}{2}w^T w + \sum_i \lambda_i (1 - y_i(w^T x^{(i)} + b))$$

Convex in w , so take derivatives to form the dual

$$w = \sum_i \lambda_i y_i x^{(i)}$$

$$\sum_i \lambda_i y_i = 0$$

$$\max_{\lambda \geq 0} -\frac{1}{2} \sum_i \sum_j \lambda_i \lambda_j y_i y_j x^{(i)T} x^{(j)} + \sum_i \lambda_i$$

such that

$$\sum_i \lambda_i y_i = 0$$

- By strong duality, solving this problem is equivalent to solving the primal problem
 - Given the optimal λ , we can easily construct w (b can be found by **complementary slackness...**)

Complementary Slackness



- Suppose that there is zero duality gap
- Let x^* be an optimum of the primal and (λ^*, ν^*) be an optimum of the dual

$$\begin{aligned} f_0(x^*) &= g(\lambda^*, \nu^*) \\ &= \inf_x \left[f_0(x) + \sum_{i=1}^m \lambda_i^* f_i(x) + \sum_{i=1}^p \nu_i^* h_i(x) \right] \\ &\leq f_0(x^*) + \sum_{i=1}^m \lambda_i^* f_i(x^*) + \sum_{i=1}^p \nu_i^* h_i(x^*) \\ &= f_0(x^*) + \sum_{i=1}^m \lambda_i^* f_i(x^*) \\ &\leq f_0(x^*) \end{aligned}$$

Complementary Slackness



- This means that

$$\sum_{i=1}^m \lambda_i^* f_i(x^*) = 0$$

- As $\lambda \geq 0$ and $f_i(x_i^*) \leq 0$, this can only happen if $\lambda_i^* f_i(x^*) = 0$ for all i
- Put another way,
 - If $f_i(x^*) < 0$ (i.e., the constraint is not tight), then $\lambda_i^* = 0$
 - If $\lambda_i^* > 0$, then $f_i(x^*) = 0$
 - ONLY applies when there is no duality gap

Dual SVM (Obtaining b)



$$\max_{\lambda \geq 0} -\frac{1}{2} \sum_i \sum_j \lambda_i \lambda_j y_i y_j x^{(i)T} x^{(j)} + \sum_i \lambda_i$$

such that

$$\sum_i \lambda_i y_i = 0$$

- By complementary slackness, $\lambda_i^* > 0$ means that $x^{(i)}$ is a support vector (can then solve for b using w)
- In particular,

$$b = y_i - w \cdot x_i$$

for any i where $\lambda_i > 0$

$$\max_{\lambda \geq 0} -\frac{1}{2} \sum_i \sum_j \lambda_i \lambda_j y_i y_j x^{(i)T} x^{(j)} + \sum_i \lambda_i$$

such that

$$\sum_i \lambda_i y_i = 0$$

- Takes $O(n^2)$ time just to evaluate the objective function
 - Active area of research to try to speed this up

- For some feature vectors, we can compute the inner products quickly, even if the feature vectors are very large
- This is best illustrated by example

- Let $\phi(x_1, x_2) = \begin{bmatrix} x_1 x_2 \\ x_2 x_1 \\ x_1^2 \\ x_2^2 \end{bmatrix}$

- $$\begin{aligned} \phi(x_1, x_2)^T \phi(z_1, z_2) &= x_1^2 z_1^2 + 2x_1 x_2 z_1 z_2 + x_2^2 z_2^2 \\ &= (x_1 z_1 + x_2 z_2)^2 \\ &= (x^T z)^2 \end{aligned}$$

The Kernel Trick



- For some feature vectors, we can compute the inner products quickly, even if the feature vectors are very large
- This is best illustrated by example

- Let $\phi(x_1, x_2) = \begin{bmatrix} x_1 x_2 \\ x_2 x_1 \\ x_1^2 \\ x_2^2 \end{bmatrix}$

- $$\begin{aligned} \phi(x_1, x_2)^T \phi(z_1, z_2) &= x_1^2 z_1^2 + 2x_1 x_2 z_1 z_2 + x_2^2 z_2^2 \\ &= (x_1 z_1 + x_2 z_2)^2 \\ &= (x^T z)^2 \end{aligned}$$

Reduces to a dot product in the original space

- The same idea can be applied for the feature vector ϕ of all polynomials of degree (exactly) d
 - $\phi(x)^T \phi(z) = (x^T z)^d$
- More generally, a **kernel** is a function $k(x, z) = \phi(x)^T \phi(z)$ for some feature map ϕ
- Rewrite the dual objective

$$\max_{\lambda \geq 0, \sum_i \lambda_i y_i = 0} -\frac{1}{2} \sum_i \sum_j \lambda_i \lambda_j y_i y_j k(x^{(i)}, x^{(j)}) + \sum_i \lambda_i$$

Examples of Kernels



- Polynomial kernel of degree exactly d
 - $k(x, z) = (x^T z)^d$
- General polynomial kernel of degree d for some c
 - $k(x, z) = (x^T z + c)^d$
- Gaussian kernel for some σ (RBF Kernel)
 - $k(x, z) = \exp\left(\frac{-\|x-z\|^2}{2\sigma^2}\right)$
 - The corresponding ϕ is infinite dimensional!
- Sigmoid Kernel
 - $k(x, z) = \tanh(\gamma \cdot x^T z + r)$

- Consider the Gaussian kernel

$$\begin{aligned}\exp\left(\frac{-\|x - z\|^2}{2\sigma^2}\right) &= \exp\left(\frac{-(x - z)^T(x - z)}{2\sigma^2}\right) \\ &= \exp\left(\frac{-\|x\|^2 + 2x^T z - \|z\|^2}{2\sigma^2}\right) \\ &= \exp\left(-\frac{\|x\|^2}{2\sigma^2}\right) \exp\left(-\frac{\|z\|^2}{2\sigma^2}\right) \exp\left(\frac{x^T z}{\sigma^2}\right)\end{aligned}$$

- Use the Taylor expansion for $\exp()$

$$\exp\left(\frac{x^T z}{\sigma^2}\right) = \sum_{n=0}^{\infty} \frac{(x^T z)^n}{\sigma^{2n} n!}$$

- Consider the Gaussian kernel

$$\begin{aligned}\exp\left(\frac{-\|x - z\|^2}{2\sigma^2}\right) &= \exp\left(\frac{-(x - z)^T(x - z)}{2\sigma^2}\right) \\ &= \exp\left(\frac{-\|x\|^2 + 2x^T z - \|z\|^2}{2\sigma^2}\right) \\ &= \exp\left(-\frac{\|x\|^2}{2\sigma^2}\right) \exp\left(-\frac{\|z\|^2}{2\sigma^2}\right) \exp\left(\frac{x^T z}{\sigma^2}\right)\end{aligned}$$

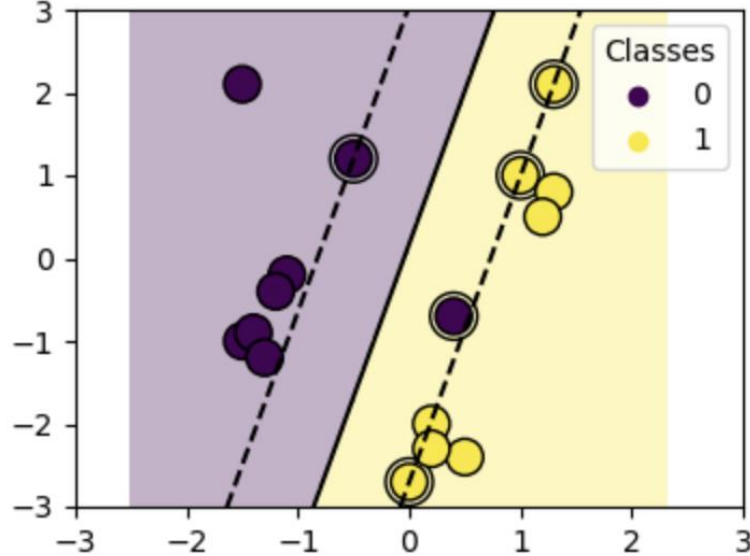
- Use the Taylor expansion for $\exp()$

$$\exp\left(\frac{x^T z}{\sigma^2}\right) = \sum_{n=0}^{\infty} \frac{(x^T z)^n}{\sigma^{2n} n!}$$

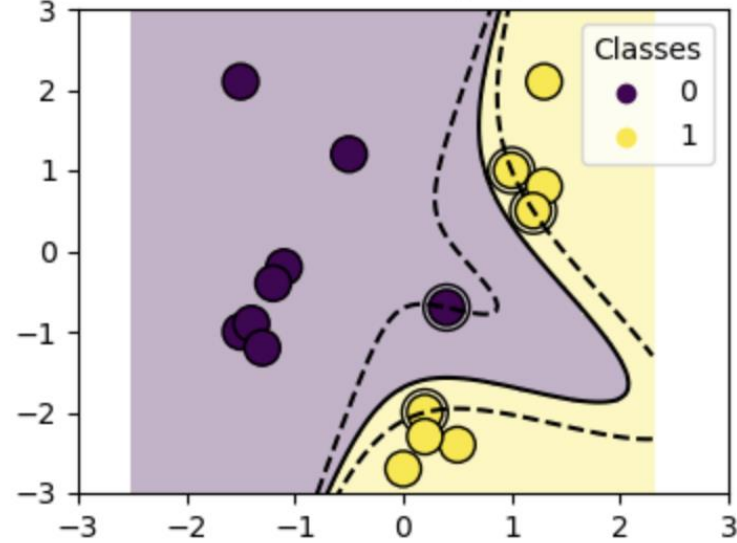
Polynomial kernels of every degree!

Illustrating Different Kernels

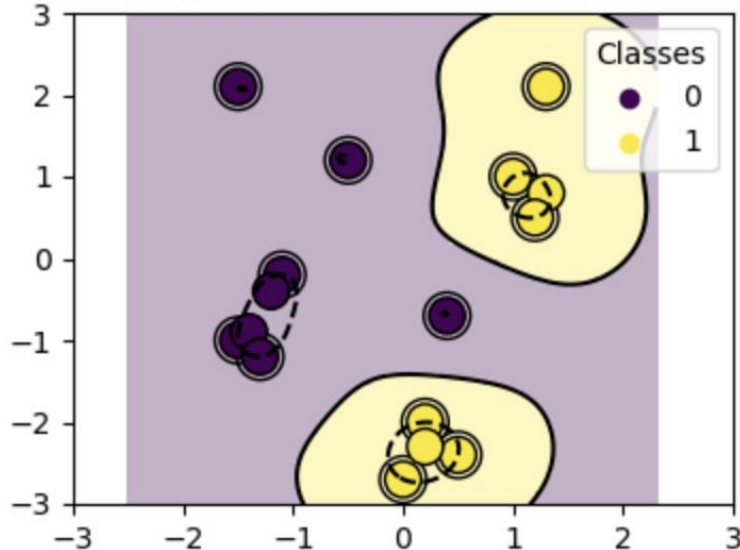
Decision boundaries of linear kernel in SVC



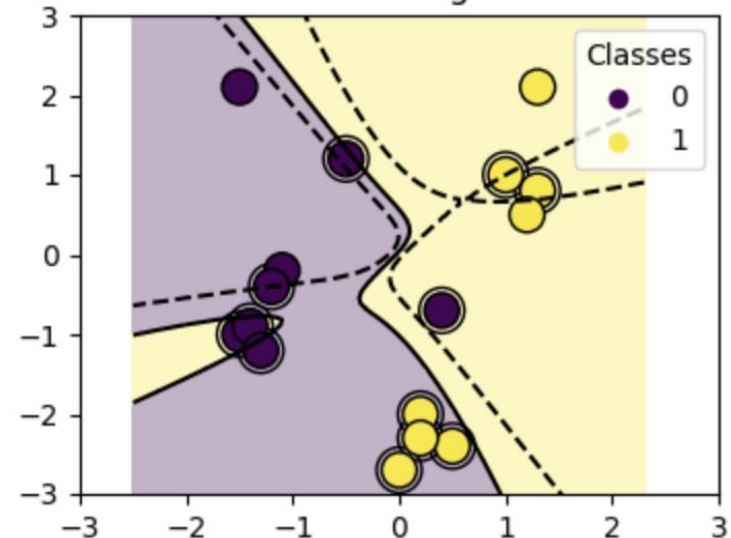
Decision boundaries of poly kernel in SVC



Decision boundaries of rbf kernel in SVC



Decision boundaries of sigmoid kernel in SVC



- Bigger feature space increases the possibility of overfitting
 - Large margin solutions may still generalize reasonably well
- Alternative: add “penalties” to the objective to disincentivize complicated solutions