



Extra Topic (Reading Material)

Lagrange Multipliers & Duality of SVMs

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The Strategy So Far...

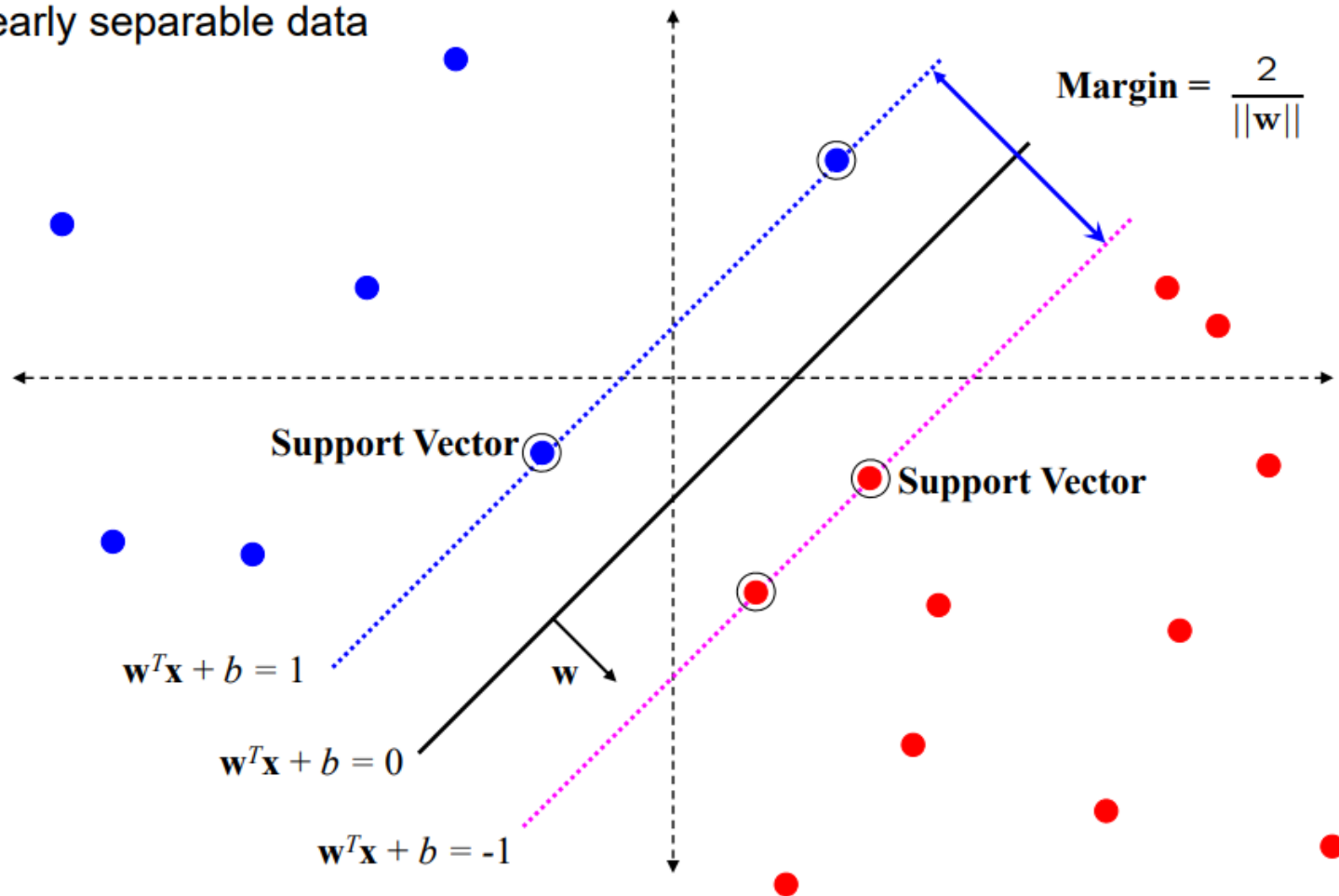


- Choose hypothesis space
- Construct loss function (ideally convex)
- Minimize loss to “learn” correct parameters

Recap: SVM



linearly separable data



- Recall: The SVM optimization problem:

$$\min_{w,b} \|w\|^2$$

such that

$$y^{(i)}(w^T x^{(i)} + b) \geq 1, \text{ for all } i$$

- This is a standard quadratic programming problem
 - Falls into the class of **convex optimization problems**
 - Can be solved with many specialized optimization tools (e.g., `quadprog()` in MATLAB)

Constrained Optimization



A mathematical detour, we'll come back to SVMs soon!

$$\begin{aligned} & \min_{x \in \mathbb{R}^n} f_0(x) \quad \leftarrow \frac{1}{2} \|w\|^2 \\ \text{subject to:} & \quad (w, b) \\ & \rightarrow f_i(x) \leq 0, \quad i = 1, \dots, m \quad \leftarrow 1 - y^{(i)}(w^T x^{(i)} + b) \\ & \quad h_i(x) = 0, \quad i = 1, \dots, p \quad \leq 0 \end{aligned}$$

$$\min_{x \in \mathbb{R}^n} f_0(x)$$

f_0 is not necessarily convex

subject to:

$$\begin{aligned} f_i(x) &\leq 0, & i &= 1, \dots, m \\ h_i(x) &= 0, & i &= 1, \dots, p \end{aligned}$$

$$\min_{x \in \mathbb{R}^n} f_0(x)$$

subject to:

$$\begin{aligned} f_i(x) &\leq 0, & i &= 1, \dots, m \\ h_i(x) &= 0, & i &= 1, \dots, p \end{aligned}$$

Constraints do not need to
be linear

Example



$$\min_{x \in \mathbb{R}^2} x_1 \log x_1 + x_2 \log x_2$$

subject to:

$$x_1 + x_2 = 1$$

$$x_1 \geq 0$$

$$x_2 \geq 0$$

Example



$$\begin{array}{c} f_0(n) \\ \min_{x \in \mathbb{R}^3} x_1 \log x_1 + x_2 \log x_2 \end{array}$$

subject to:

$$\begin{array}{lll} h_1(n) = 0 & \longrightarrow & 1 - x_1 - x_2 = 0 \quad h_1(n) \\ f_1(n) \leq 0 & \longrightarrow & -x_1 \leq 0 \\ f_2(n) \leq 0 & \longrightarrow & -x_2 \leq 0 \end{array} \quad \left. \begin{array}{l} \\ \\ \end{array} \right\} f_i(n)$$

Lagrangian



$$\begin{array}{ccccc}
 \text{Obj} & & \text{Ineq Const-} & & \text{Eq Const-} \\
 \downarrow & & \downarrow & & \downarrow \\
 L(x, \lambda, \nu) = f_0(x) + \sum_{i=1}^m \lambda_i f_i(x) + \sum_{i=1}^p \nu_i h_i(x) \\
 & & \downarrow & & \uparrow \\
 & & f_i(x) \leq 0 & & h_i(x) = 0
 \end{array}$$

- Incorporate constraints into a new objective function
- $\lambda \geq 0$ and ν are vectors of **Lagrange multipliers**
- The Lagrange multipliers can be thought of as enforcing soft constraints

Example



$$\min_{x \in \mathbb{R}^3} x_1 \log x_1 + x_2 \log x_2$$

subject to:

$$\begin{aligned} 1 - x_1 - x_2 &= 0 \leftarrow v_1 \rightarrow \text{eq.} \\ -x_1 &\leq 0 \leftarrow \lambda_1 \\ -x_2 &\leq 0 \leftarrow \lambda_2 \end{aligned} \quad \text{Ineq.}$$

$$\begin{aligned} L(x_1, x_2, v_1, \lambda_1, \lambda_2) \\ = x_1 \log x_1 + x_2 \log x_2 + v_1 \cdot (1 - x_1 - x_2) - \lambda_1 x_1 - \lambda_2 x_2 \end{aligned}$$

- Construct a **dual function** by minimizing the Lagrangian over the primal variables

$$\underbrace{g(\lambda, \nu)}_{\min_x} = \inf_x L(x, \lambda, \nu)$$

- $g(\lambda, \nu) = -\infty$ whenever the Lagrangian is not bounded from below for a fixed λ and ν

Example



$$\min_{x \in \mathbb{R}^2} x_1 \log x_1 + x_2 \log x_2$$

subject to:

$$1 - x_1 - x_2 = 0$$

$$-x_1 \leq 0$$

$$-x_2 \leq 0$$

Example



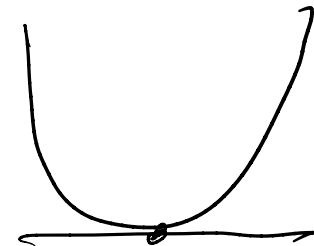
$$\min_{x \in \mathbb{R}^3} x_1 \log x_1 + x_2 \log x_2$$

subject to:

$$1 - x_1 - x_2 = 0$$

$$-x_1 \leq 0$$

$$-x_2 \leq 0$$



$$L(x_1, x_2, v_1, \lambda_1, \lambda_2)$$

$$= x_1 \log x_1 + x_2 \log x_2 + v_1 \cdot (1 - x_1 - x_2) - \lambda_1 x_1 - \lambda_2 x_2$$

$$\frac{\partial L}{\partial x_1} = 1 + \log x_1 - v_1 - \lambda_1 = 0 \Rightarrow x_1 = \exp(v_1 + \lambda_1 - 1)$$

$$\frac{\partial L}{\partial x_2} = 1 + \log x_2 - v_1 - \lambda_2 = 0 \Rightarrow x_2 = \exp(v_1 + \lambda_2 - 1)$$

$$\frac{\partial L}{\partial v_1} = 1 - x_1 - x_2 = 0$$

$$g(\lambda, v) = \text{Subst. } x_1 \text{ \& } x_2 \text{ into } L(x_1, x_2, v_1, \lambda_1, \lambda_2)$$

Example



$$\min_{x \in \mathbb{R}^3} x_1 \log x_1 + x_2 \log x_2$$

subject to:

$$1 - x_1 - x_2 = 0$$

$$-x_1 \leq 0$$

$$-x_2 \leq 0$$

The Primal Problem



$$\min_{x \in \mathbb{R}^n} f_0(x) = \max_{\lambda \geq 0, v} L(x, \lambda, v)$$

subject to:

$$\begin{aligned} f_i(x) &\leq 0, & i &= 1, \dots, m \\ h_i(x) &= 0, & i &= 1, \dots, p \end{aligned}$$

Equivalently,

$$\min_x \max_{\lambda \geq 0, v} L(x, \lambda, v)$$

Why are these equivalent?

The Primal Problem



subject to:

$$\min_{x \in \mathbb{R}^n} f_0(x)$$

$$\begin{aligned} f_i(x) &\leq 0, & i &= 1, \dots, m \\ h_i(x) &= 0, & i &= 1, \dots, p \end{aligned}$$

Original
Problem

Equivalently,

$$\min_x \max_{\lambda \geq 0, v} L(x, \lambda, v)$$

$$\sup_{\lambda \geq 0, v} \left[f_0(x) + \sum_{i=1}^m \lambda_i f_i(x) + \sum_{i=1}^p v_i h_i(x) \right] = \infty$$

whenever x violates the constraints

The Dual Problem



Equivalently,

$$\begin{array}{c} \max_{\lambda \geq 0, \nu} g(\lambda, \nu) \\ \downarrow \\ \max_{\lambda \geq 0, \nu} \min_x L(x, \lambda, \nu) \end{array}$$

- The dual problem is always concave, even if the primal problem is not convex
 - For each x , $L(x, \lambda, \nu)$ is a linear function in λ and ν
 - Minimum (or infimum) of linear functions is concave!

Primal vs. Dual



$$\underbrace{\sup_{\lambda \geq 0, v} \inf_x L(x, \lambda, v)}_{\text{Dual.}} \leq \underbrace{\inf_x \sup_{\lambda \geq 0, v} L(x, \lambda, v)}_{\text{Primal.}}$$

• Why?

- $g(\lambda, v) \leq L(x, \lambda, v)$ for all x $g(\lambda, v) = \min_x L(x, \lambda, v)$
- $L(x', \lambda, v) \leq f_0(x')$ for any feasible $x', \lambda \geq 0$
 - x is **feasible** if it satisfies all of the constraints
- Let x^* be the optimal solution to the primal problem and $\lambda \geq 0$

$$g(\lambda, v) \leq L(x^*, \lambda, v) \leq f_0(x^*)$$

$$L(x, \lambda, v) = f_0(x) + \sum_{i=1}^m \lambda_i f_i(x) + \sum_{i=1}^p v_i h_i(x) = 0$$

$\lambda_i \geq 0$

- Under certain conditions, the two optimization problems are equivalent

$$\sup_{\lambda \geq 0, \nu} \inf_x L(x, \lambda, \nu) = \inf_x \sup_{\lambda \geq 0, \nu} L(x, \lambda, \nu)$$

- This is called **strong duality**
- If the inequality is strict, then we say that there is a **duality gap**
 - Size of gap measured by the difference between the two sides of the inequality

Slater's Condition



For any optimization problem of the form

subject to:

$$\begin{array}{l} \min_{x \in \mathbb{R}^n} \underline{f_0(x)} \leftarrow \text{Convex} \\ \text{Convex.} \\ \downarrow \\ \underline{f_i(x)} \leq 0, \quad i = 1, \dots, m \\ Ax = b \leftarrow \text{Linear Equality.} \end{array}$$

where f_0, \dots, f_m are **convex functions**, strong duality holds if there exists an x such that

$$\begin{array}{l} f_i(x) < 0, \quad i = 1, \dots, m \\ Ax = b \end{array}$$

$$\min_w \frac{1}{2} \|w\|^2 \leftarrow \text{Objective}$$

such that

$$y_i(w^T x^{(i)} + b) \geq 1, \text{ for all } i \leftarrow \text{Ineq Constraints}$$

- Note that Slater's condition holds as long as the data is linearly separable

$$L(w, b, \lambda) = \frac{1}{2} w^T w + \sum_i \lambda_i (1 - y_i (w^T x^{(i)} + b))$$

Convex in w , so take derivatives to form the dual

$$\frac{\partial L}{\partial w_k} = w_k + \sum_i -\lambda_i y_i x_k^{(i)} = 0$$

$$\frac{\partial L}{\partial b} = \sum_i -\lambda_i y_i = 0$$

$$L(w, b, \lambda) = \frac{1}{2} w^T w + \sum_i \lambda_i (1 - y_i (w^T x^{(i)} + b))$$

Convex in w , so take derivatives to form the dual

$$w = \sum_i \lambda_i y_i x^{(i)}$$

$$\sum_i \lambda_i y_i = 0$$

$$\max_{\lambda \geq 0} -\frac{1}{2} \sum_i \sum_j \lambda_i \lambda_j y_i y_j \underbrace{x^{(i)T} x^{(j)}}_{\substack{\uparrow \\ K(x^{(i)}, x^{(j)})}} + \sum_i \lambda_i$$

such that

$$\sum_i \lambda_i y_i = 0$$

- By strong duality, solving this problem is equivalent to solving the primal problem
 - Given the optimal λ , we can easily construct w (b can be found by **complementary slackness...**)

$$x_{\text{test}}, \quad y_{\text{pred}} = \text{Sign}(w^T x_{\text{test}} + b)$$

- the dual
- $$f_0(x^*) = g(\lambda^*, v^*) \leftarrow \text{Primal} = \text{Dual.}$$

$$\leq f_0(x^*) + \sum_{i=1}^m \lambda_i^* f_i(x^*) + \sum_{i=1}^p \nu_i^* h_i(x^*) \quad \color{red}{= 0}$$

$$p_i^* \geq 0, \quad f_i(x^*) \leq 0$$

Complementary Slackness



- This means that

$$\sum_{i=1}^m \lambda_i^* f_i(x^*) = 0$$

Handwritten red annotations: a bracket above the sum with ≤ 0 written above it.

- As $\lambda \geq 0$ and $f_i(x_i^*) \leq 0$, this can only happen if $\lambda_i^* f_i(x^*) = 0$ for all i
- Put another way,
 - If $f_i(x^*) < 0$ (i.e., the constraint is not tight), then $\lambda_i^* = 0$
 - If $\lambda_i^* > 0$, then $f_i(x^*) = 0$ *← i is a support vector*
- ONLY applies when there is no duality gap

Dual SVM (Obtaining b)



$$\max_{\lambda \geq 0} -\frac{1}{2} \sum_i \sum_j \lambda_i \lambda_j y_i y_j x^{(i)T} x^{(j)} + \sum_i \lambda_i$$

such that

$$\sum_i \lambda_i y_i = 0$$

- By complementary slackness, $\lambda_i^* > 0$ means that $x^{(i)}$ is a support vector (can then solve for b using w)
- In particular,

$$b = y_i - w \cdot x_i$$

for any i where $\lambda_i > 0$ [support vectors]

$$\max_{\lambda \geq 0} -\frac{1}{2} \sum_{i=1}^M \sum_{j=1}^M \lambda_i \lambda_j y_i y_j x^{(i)T} x^{(j)} + \sum_i \lambda_i$$

such that

$$\sum_i \lambda_i y_i = 0$$

$O(M^2)$

- Takes ~~$O(n^2)$~~ time just to evaluate the objective function
 - Active area of research to try to speed this up