

Planning and Scheduling Summary

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Planning and Scheduling mid semester summary in preparation of exam

1 Introduction

- DEFINITION: Plan
 - A collection of actions for performing some task or achieving some objective
- Conceptual model

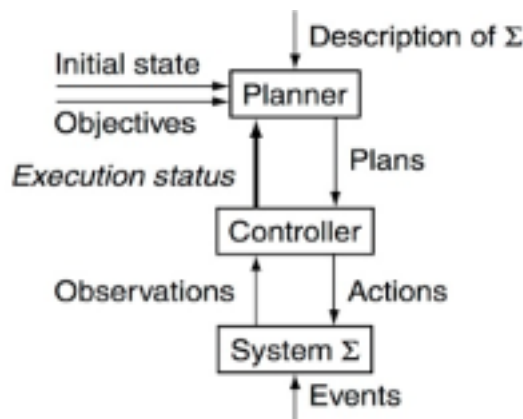


Figure 1: Conceptual model

- State-transition system
 $\Sigma = (S, A, E, \gamma)$
 - S : states
 - A : actions (controllable)
 - E : events (uncontrollable)
 - $\gamma : S \times (A \cup E) \mapsto 2^S$: state-transition function
- Observation function
 - $h : S \mapsto O$

- Planner
 - description of Σ , initial state $s_0 \in S$, some objective
- Objectives
 - S_g : set of goal states
 - condition over set of states followed by system
 - utility function
- Restrictive assumptions
 - A0: finite σ
 - A1: σ is fully observable (observation function is $\text{id}()$)
 - A2: deterministic σ , action: one possible outcome
 - A3: static σ : E is empty
 - A4: attainment goals: goal state or a set of goal states
 - A5: sequential plans
 - A6: implicit time (instant transitions)
 - A7: off-line planning
- Classical planning
 - requires all eight restrictive assumptions
 - Given (Σ, s_0, S_g)
 - find a sequence of actions $\langle a_1, a_2, \dots, a_n \rangle$
 - that produces sequence of transitions $s_1 = \gamma(s_0, a_1), \dots$
 - such that $s_n \in S_g$
- Relaxing assumptions

- A0: discrete logic (e.g. 1st order), continuous variables)
- A1: if we don't relax other restrictions, only uncertainty is about s_0
- A2: nondeterministic outcomes, seek policy/contingency plan, MDPs (probabilities)
- A1+A2: Finite POMDPs
- A0+A2: Continuous or hybrid MDPs
- A0+A1+A2: Continuous or hybrid MDPs
- A3: Other agents/dynamic environment, randomly behaving environment
- A1+A3: Imperfect-information games
- A5/A6: Temporal planning
- A0+A5+A6: Planning and resource scheduling

2 Search and Complexity

Problem Solving Agents

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Problem Types

- Deterministic, fully observable
- Non-observable
- Nondeterministic

Problem Formulation

Example Problems

Basic Search Algorithms

3 Logic and Inference

4 Representing Plans

- **Classical representation**
 - function-free FOL
 - Atom: predicate symbol and args
 $on(c1, c3), on(c1, x)$
 - Ground expression: instantiated var
 $on(c1, c3)$
 - Nonground expression: at least one var
 $on(c1, x)$
 - Substitution $\theta = x_1 \leftarrow v_1, \dots$

- Instance of expression e : result of applying θ

- **State** : set of ground atoms
- **Operator**
 - (name, precondition, effects)
 - $take(k, l, c, d, p)$
 $precond : belong(k, l)$
 $effects : holding(k, c), \neg empty(k)$
 - no need for negative effects (closed world assumption), only when they are explicitly touched (AM?)
- **Action**
 - ground **instance** of operator
- **Notation**
 - $S^+, S^-, precondition^+, \dots$
- **Result of action**
 - $\gamma(s, a) = (s \setminus effects^-(a) \cup effects^+(a))$
- **Planning problem**
 - given: domain (L, O) (Language, Operators)
 - $P = (\Sigma, s_0, S_g)$ (state-trans., initState, goal formula)
 - $\Sigma = (S, A, E, \gamma)$
- **Plan**
 - $\sigma = \langle a_1, a_2, \dots, a_n \rangle$
 - is a solution for P if it is executable and achieves g
- **Set-theoretic representation**
 - classical representation + restriction to PL
 - collection of propositions (bool) instead of ground atoms
 - effects: delete, add
- exponential blow-up
- **State-Variable representation**
 - State variables: $pos(x) = y$ (if block x is on block y)
- **Expressive power**
- linear time/space conversion between all three representations, except for expBlowup on conversion to set-theoretic

- classical rep: most popular
- set-theoretic: much more space, useful for algorithms that manipulate ground atoms directly
- state-variable: equivalent to classical, more natural for engineers, used for non-classical planning problems (incl. functions!)

5 Complexity of Classical Planning

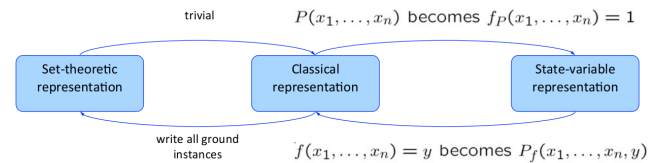
- **Complexity analysis**
- general idea: translate classical planning -i language-recognition problem and examine its complexity
- Language L , alphabet A
- recognition procedure $R(x)$: yes iff $x \in L$, else no or fail to terminate
 - **Plan-Existence** P
 - Problem has solution
 - **Plan-Length** (P, n)
 - Problem has solution $j = n$
- Runtime, space complexity
- **Restrictions of classical planning**
- Operators are fixed or input?
- Allow infinite initial states?*
- Allow function symbols?*
- Allow negative effects?
- Allow neg. preconditions?
- Allow \neg preconditions?
- Operators have conditional effects?*
- *: outside CP
- **(Un)decidability**

	Decidability	
function symbols	Plan-Existence	Plan-Length
no (CP)	dec	dec
yes	semidec	dec

- **Complexity results**
- Well I guess, Plan-Existence is less complex than Plan-Length (AM?)

• Equivalences

- Set-theoretic and classical are basically identical
- Both: exponential blowup (input)
- Class. and state-var are basically equivalent
- (state-var:some restrictions not possible)



6 State Space Planning

• State space vs plan space

- State space
 - Node: state of the world
 - Plan: path through space
- Plan space
 - State: set of partially instantiated operators and some constraints
 - Impose more constraints until Plan

• Linear search

- Work on one goal until completely solved
- Order of problem solving is linearly-related to order in which plan actions are executed

• Maintain goal stack

- Implications
 - No interleaving of goal achievement
 - Efficient search if goals do "not" interact

• Means-End analysis

- Search relevant aspects of problem, means/operators, ends/goals
 - Start from the goal
 - find difference to start state
 - find operator that reduces this difference

• General Problem Solver (GPS)

- **Forward search**

```

s ← s0
π ← empty_plan
loop
if s satisfies g → return π
E ← {a ground instance of o ∈ O, precondition(a)
true in s}
if E = ∅ → failure
nondet choose a ∈ E
s ← γ(s, a)
π ← π, a
endloop

```

- nondet choosing -i (par/seq) do all actions a
- seq, we need to use backtracking (as soon as a zero set or goal state is reached)
- = sound, complete
- breadth-first, best-first
 - sound and complete
 - memory exponential in length of solution
- depth-first, greedy search
 - Worst-case mem is linear in length of solution
 - Sound, but not complete
 - but CP has only finite states → loop-checking solves completeness
- large branching factor (need good heuristic / pruning)

- **Backward search**

- Means-end analysis

- start at goal and compute inverse state transitions
- a makes at least one of g's literals true and non false
- $g' = \gamma^{-1}(g, a) = (g \setminus effects(a)) \cup precondition(a)$
- take the goal state, remove effects of a and add the preconditions


```

π ← empty_plan
loop
if s satisfies g → return π
A ← {a | a ground instance of o in O and g' =
γ-1(g, a) defined }
if A = ∅ failure
nondet choose a ∈ A

```

```

π ← a, π
g ← γ-1(g, a)
endloop

```

- Branching factor smaller, but can be big because of more actions than needed
- Problem with (x,y): you need to instantiate y with every state

- **Lifting**

- reduce branching factor by **partially instantiating** operators (=lifting)


```

π ← empty_plan
loop
if s0 satisfies g → return π
A ← {(o, θ) | o standardization of o in O, θ is mgu for an atom of g and atom of effects+(o) and γ-1(θ(g), θ(o)) defined}
if A = ∅ failure
nondet choose (o, θ) ∈ A
π ← concatenation of θ(o) and θ(π)
g ← γ-1(θ(g), θ(o))
endloop

```

- more complicated than backward-search (keep track of substitutions)
- smaller branching factor
- if sub-problems independent, all orderings need to be tried

- **STRIPS**

- Modified backward search (instead of γ^{-1} , each set of sub-goals is precondition(a))
- a = executable? go forward as far as possible.
- (AM?) We need to write down all search algorithms as examples on the whiteboard and copy them here. Algorithms written out are just too unclear and direct comparison is difficult.
- STRIPS assumption -i solved frame problem
- introduces: difference, sub-goals
- Block stacking
- Sussman Anomaly: STRIPS cannot produce irredundant solution because it is non-interleaved.
- Register Assignment Problem (TODO!)
- **Linear Planning**

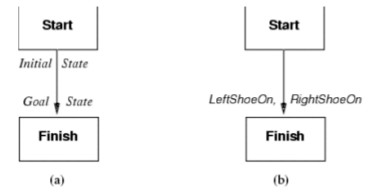
- + reduced search space (goals one at a time)
- + advantageous if goals are independent
- + sound
- - suboptimal solutions
- - incomplete

- solve -: domain-specific knowledge (fwd/bwd search)
 - ds knowledge can prune search space
 - ds specific algorithm
 - TODO solve Sussman Anomaly with ds knowledge

7 Plan Space Planning

- **Motivation**
- in SSP: try all orderings before "no solution"
- PSP: do not commit to orderings, instantiations, etc.
- **Partial Order Plan vs Total Order Plan**
- POP: "parallel" paths
- **PSP**: Backward search from g
- Each node is **partial plan**
 - { part.-instantiated operators (**steps**) }
 - Sets of constraints
- Refine until solution
- **Constraints**
 - **Precedence constraints**: a must precede b
 - **Binding constraints**: (in)equality
 - **Causal links**: use step a to establish prec p needed by c
- No more **flaws** → plan
- **Representation**
- $P = \langle \text{steps, orderings, bindings, causallinks} \rangle$

$Plan(Steps: \{S_1 : Op(ACTION:Start),$
 $S_2 : Op(ACTION:Finish,$
 $PREC: RightShoeOn \wedge LeftShoeOn)\}$
 $ORDERINGS: \{S_1 \prec S_2\}$
 $BINDINGS: \{\}$
 $CAUSALLINKS: \{\}$)



8 Graph Based Planning

9 Satisfiability Based Planning

10 Hierarchical Task Network Planning

* partial ordering enables interleaving of tasks. E.g. we need a hammer and a drill and we do not care in what order