

COM 410 – Slides 4

Representing and Manipulating
Information

Some Vocabulary

- Bits
- Bytes
- Overflow
- Memory and virtual memory
- Memory address and virtual address space
- Program objects
- Information type
- Word size
- Big endian vs. little endian
- MSB, LSB

Three types of number representation:

- Unsigned notation
(binary and hexadecimal notations)
- Two's Complement
(positive and negative numbers)
- Floating Point

Digit, Hex, and Binary notations

Hex digit	0	1	2	3	4	5	6	7
Decimal value	0	1	2	3	4	5	6	7
Binary value	0000	0001	0010	0011	0100	0101	0110	0111
Hex digit	8	9	A	B	C	D	E	F
Decimal value	8	9	10	11	12	13	14	15
Binary value	1000	1001	1010	1011	1100	1101	1110	1111

Type sizes are not guaranteed!

C declaration		Bytes	
Signed	Unsigned	32-bit	64-bit
<code>[signed] char</code>	<code>unsigned char</code>	1	1
<code>short</code>	<code>unsigned short</code>	2	2
<code>int</code>	<code>unsigned</code>	4	4
<code>long</code>	<code>unsigned long</code>	4	8
<code>int32_t</code>	<code>uint32_t</code>	4	4
<code>int64_t</code>	<code>uint64_t</code>	8	8
<code>char *</code>		4	8
<code>float</code>		4	4
<code>double</code>		8	8

```
typedef long                SLong;
typedef unsigned long       ULong;
typedef short               SWord;
typedef unsigned short      UWord;
typedef char                SByte;
typedef unsigned char       UByte;
//typedef unsigned char      Boolean;
typedef unsigned short      Boolean;
//typedef long double        GFloat;
```

Ox01234567 stored as:

Big endian

	0x100	0x101	0x102	0x103	
...	01	23	45	67	...

Little endian

	0x100	0x101	0x102	0x103	
...	67	45	23	01	...

Representations on different machines:

Machine	Value	Type	Bytes (hex)
Linux 32	12,345	int	39 30 00 00
Windows	12,345	int	39 30 00 00
Sun	12,345	int	00 00 30 39
Linux 64	12,345	int	39 30 00 00
Linux 32	12,345.0	float	00 e4 40 46
Windows	12,345.0	float	00 e4 40 46
Sun	12,345.0	float	46 40 e4 00
Linux 64	12,345.0	float	00 e4 40 46
Linux 32	&ival	int *	e4 f9 ff bf
Windows	&ival	int *	b4 cc 22 00
Sun	&ival	int *	ef ff fa 0c
Linux 64	&ival	int *	b8 11 e5 ff ff 7f 00 00

Boolean Algebra Overview

$$\begin{array}{r} 0110 \\ \& \underline{1100} \\ 0100 \end{array}$$

$$\begin{array}{r} 0110 \\ | \underline{1100} \\ 1110 \end{array}$$

$$\begin{array}{r} 0110 \\ \sim \underline{1100} \\ 1010 \end{array}$$

$$\begin{array}{r} \sim \underline{1100} \\ 0011 \end{array}$$

Shifting Operations

Operation	Value 1	Value 2
Argument x	[01100011]	[10010101]
x << 4	[00110000]	[01010000]
x >> 4 (logical)	[00000110]	[00001001]
x >> 4 (arithmetic)	[00000110]	[11111001]

Masking Operation

e.g. $0x25AF3255 \& 0x00FF0000 = ?$

Shifting Operation

e.g. $0x25AF3255 \ll 8 = ?$

C data type	Minimum	Maximum
[signed] char	−127	127
unsigned char	0	255
short	−32,767	32,767
unsigned short	0	65,535
int	−32,767	32,767
unsigned	0	65,535
long	−2,147,483,647	2,147,483,647
unsigned long	0	4,294,967,295
int32_t	−2,147,483,648	2,147,483,647
uint32_t	0	4,294,967,295
int64_t	−9,223,372,036,854,775,808	9,223,372,036,854,775,807
uint64_t	0	18,446,744,073,709,551,615

Two's complement
(signed numbers in binary)

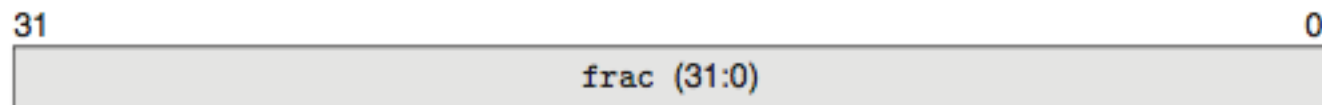
Many beginning programmers are surprised to find that adding two positive numbers can yield a negative result, and that the comparison $x < y$ can yield a different result than the comparison $x - y < 0$. These properties are artifacts of the finite nature of computer arithmetic. Understanding the nuances of computer arithmetic can help programmers write more reliable code.

Floating point

Single precision



Double precision



One useful exercise for understanding floating-point representations is to convert sample integer values into floating-point form. For example: 12,345 has binary representation [11000000111001]. We create a normalized representation of this by shifting 13 positions to the right of a binary point, giving $12,345 = 1.10000001110012 \times 2^{13}$. To encode this in IEEE single-precision format, we construct the fraction field by dropping the leading 1 and adding 10 zeros to the end, giving binary representation [10000001110010000000000]. To construct the exponent field, we add bias (127 for single precision floats) to 13, giving 140, which has binary representation [10001100]. We combine this with a sign bit of 0 to get the floating-point representation in binary of [01000110010000001110010000000000].

0	0	0	0	3	0	3	9
0000000000000000000011000000111001							

4	6	4	0	E	4	0	0
01000110010000001110010000000000							