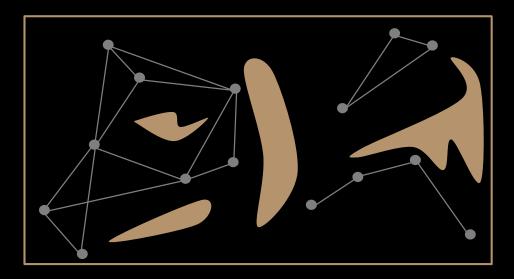
# Sampling-based Planning 2

A lot of Material from Howie Choset, Nancy Amato, Sujay Bhattacharjee, G.D. Hager, S. LaValle, J. Kuffner

## Last time...

We discussed PRMs



- Two issues with the PRM:
  - 1. Uniform random sampling misses narrow passages
  - 2. Exploring whole space, but all we want is a path

# Outline

- Sampling strategies
- RRTs

## Sampling Strategies

- Most common is uniform random sampling
  - The bigger the area, the more likely it will be sampled
  - Problem: Narrow passages



- Are narrow passages inherently bad?
  - Does A\* running on a 2D grid have problems with narrow passages?

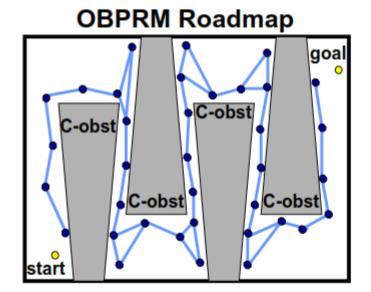
## OBPRM: An Obstacle-Based PRM

[IEEE ICRA'96, IEEE ICRA'98, WAFR'98]

#### To Navigate Narrow Passages we must sample in them

most PRM nodes are where planning is easy (not needed)

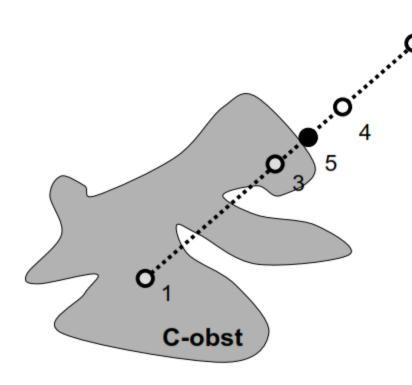
# C-obst C-obst C-obst



Idea: Can we sample nodes near C-obstacle surfaces?

we cannot explicitly construct the C-obstacles...

## OBPRM: Finding Points on C-obstacles

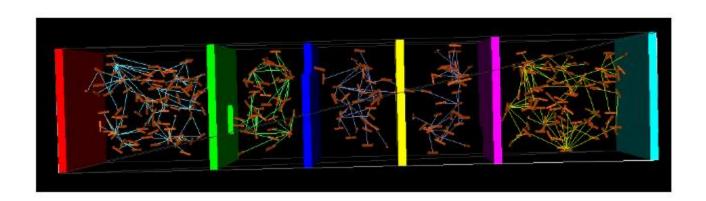


## Basic Idea (for workspace obstacle S)

- Find a point in S's C-obstacle (robot placement colliding with S)
- 2. Select a random direction in C-space
- 3. Find a free point in that direction
- Find boundary point between them using binary search (collision checks)

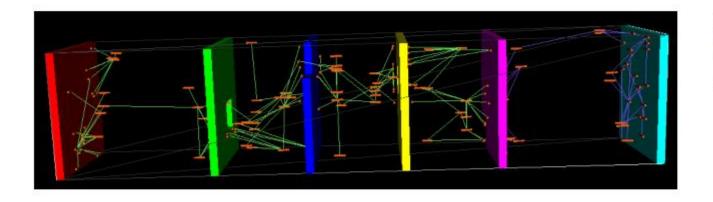
Note: we can use more sophisticated heuristics to try to cover C-obstacle

## PRM vs OBPRM Roadmaps



## **PRM**

- 328 nodes
- 4 major CCs

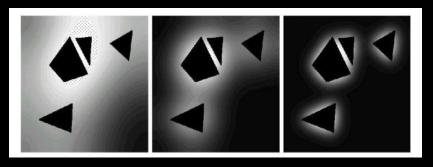


## **OBPRM**

- 161 nodes
- 2 major CCs

## Sampling strategies: Gaussian

- Gaussian sampler
  - Find a q<sub>1</sub> at random
  - Pick a q<sub>2</sub> from a Gaussian distribution centered at q<sub>1</sub>
  - If both are in collision or collision-free, discard them, if one free, keep it

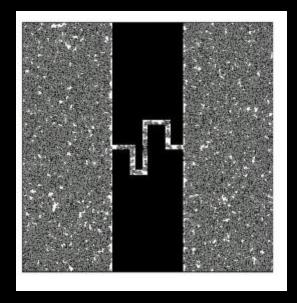


Sampling distribution for varying Guassian width (width decreasing from left to right)

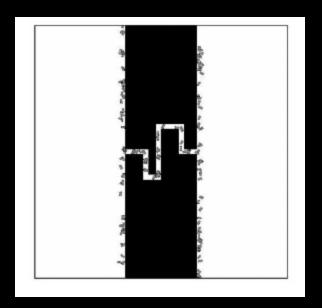
Boor, Valérie, Mark H. Overmars, and A. Frank van der Stappen. "The gaussian sampling strategy for probabilistic roadmap planners." *Robotics and Automation*, 1999. *Proceedings*. 1999 IEEE International Conference on. Vol. 2. IEEE, 1999.

## Sampling Strategies: Gaussian

Performs well in narrow passages



**Uniform Random Sampling** 



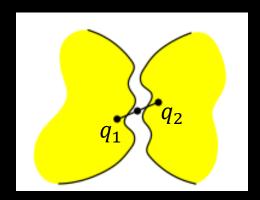
Gaussian Sampling

## Sampling Strategies

 Can we come up with a case where obstacle-biased sampling is worse than uniform random sampling?

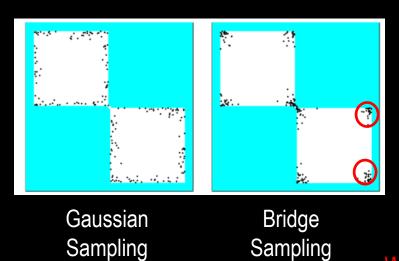
## Sampling Strategies: Bridge

- Sample a q<sub>1</sub> that is in collision
- Sample a q<sub>2</sub> in neighborhood of q<sub>1</sub> using some probability distribution (e.g. gaussian)
- If q<sub>2</sub> in collision, get the midpoint of (q<sub>1</sub>, q<sub>2</sub>)
- Check if midpoint is in collision, if not, add it as a node



Hsu, David, et al. "The bridge test for sampling narrow passages with probabilistic roadmap planners." *Robotics and Automation, 2003. Proceedings. ICRA'03. IEEE International Conference on.* Vol. 3. IEEE, 2003.

## Sampling Strategies: Bridge



What's going on at the corners?

Bridge Sampling performs well in narrow passages

- The problem: Random sampling (biased or not) can be unpredictable and irregular
  - Each time your run your algorithm you get a different sequence of samples, so performance varies
  - In the limit, space will be sampled well, but in finite time result may be irregular

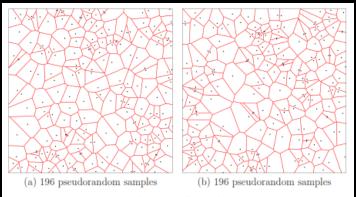


Figure 5.3: Irregularity in a collection of (pseudo)random samples can be nicely observed with Voronoi diagrams.

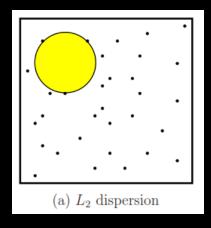
Can we do better?

- What do we care about?
  - Dispersion

$$\delta(P) = \sup_{x \in X} \big\{ \min_{p \in P} \big\{ \rho(x, p) \big\} \big\}.$$

*P* is a finite set of points,  $(X, \rho)$  is a metric space  $(\rho)$  is a distance metric)

In English: the radius of the largest empty ball



- What do we care about?
  - Discrepancy

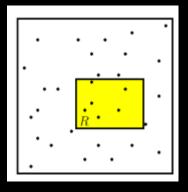
$$D(P, \mathcal{R}) = \sup_{R \in \mathcal{R}} \left\{ \left| \frac{|P \cap R|}{k} - \frac{\mu(R)}{\mu(X)} \right| \right\}$$

P is a finite set of points, k = |P|

R is the range space (set of axis-aligned boxes)

X is a metric space,  $\mu$  is a measure of volume

- In English: The largest volume estimation error that can be obtained over all sets in  $\mathcal R$
- Each term captures how well P can be used to estimate the volume of R.
  - Example: If  $\mu(R)$  is 1/2 of  $\mu(X)$ , then we want 1/2 of the points in P to be in R.



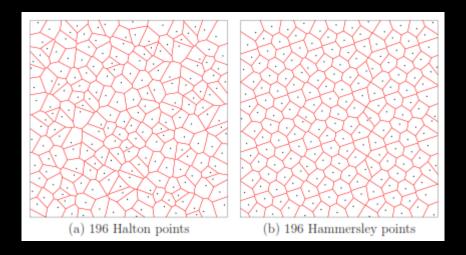
Discrepancy

- Deterministic Sampling
  - Similar to discretization we saw in Discrete Motion Planning, but order of samples matters
  - Example: van der Corput sequence

	Naive		Reverse	Van der		
i	Sequence	Binary	Binary	Corput	Points in $[0,1]/\sim$	
1	0	.0000	.0000	0	•	
2	1/16	.0001	.1000	1/2	0 0	
3	1/8	.0010	.0100	1/4	$\circ$	
4	3/16	.0011	.1100	3/4	$\circ$	
5	1/4	.0100	.0010	1/8	$\circ \hspace{-1pt} \bullet \hspace{-1pt} \circ \hspace{-1pt} \bullet \hspace{-1pt} \circ \hspace{-1pt} \circ \hspace{-1pt} \circ$	
6	5/16	.0101	.1010	5/8	0-0-0-0-0	
7	3/8	.0110	.0110	3/8	0-0-0-0-0	
8	7/16	.0111	.1110	7/8	0-0-0-0-0-0	
9	1/2	.1000	.0001	1/16	000-0-0-0-0-0	
10	9/16	.1001	.1001	9/16	000-0-0-0-0-0	
11	5/8	.1010	.0101	5/16	000-0•0-000-0-0	
12	11/16	.1011	.1101	13/16	000-000-000-0	
13	3/4	.1100	.0011	3/16	000000000000000000000000000000000000000	
14	13/16	.1101	.1011	11/16	0000000-000•000-0	
15	7/8	.1110	.0111	7/16	000000000000000000000000000000000000000	
16	15/16	.1111	.1111	15/16	000000000000000000	
Figure	Figure 5.2: The van der Corput sequence is obtained by reversing the bits in the					

Figure 5.2: The van der Corput sequence is obtained by reversing the bits in the binary decimal representation of the naive sequence.

- Halton sequence: n-dimensional generalization of van der Corput sequence
- Hammersley sequence: Adaptation of Halton sequence that yields a better distribution. BUT need to know number of samples in advance.



See LaValle's book for formulas.

# Break

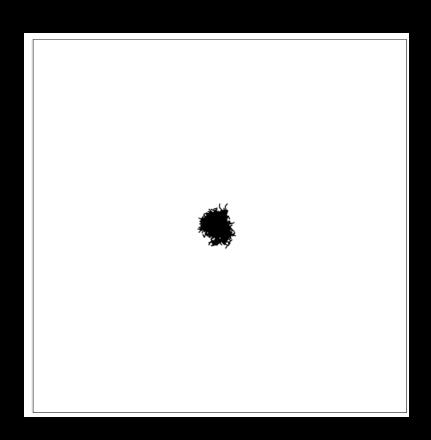
# Rapidly-exploring Random Trees (RRTs)

## Single-query methods

- Motivation: Why try to capture the connectivity of the whole space when all you need is one path?
- Algorithms:
  - Single-Query BiDirectional Lazy PRM (SBL-PRM)
  - Expansive Space Trees (EST)
  - Rapidly-exploring Random Tree (RRT)
- Key idea: Build a *tree* instead of a general graph.
- The tree "grows" in  $C_{free}$ 
  - Like PRM, captures some connectivity
  - Unlike PRM, only explores what is connected to  $q_{start}$

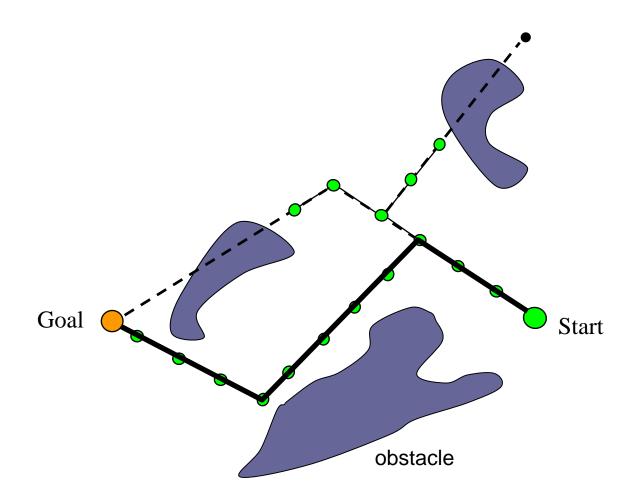
## Naïve Tree Algorithm

```
\begin{array}{l} q_{\text{node}} = q_{\text{start}} \\ \\ \text{For i = 1 to NumberSamples} \\ \\ q_{\text{rand}} = \text{Sample near } q_{\text{node}} \\ \\ \text{Add edge e = } (q_{\text{rand}} \text{ , } q_{\text{node}}) \text{ if } \\ \\ \text{collision-free} \\ \\ q_{\text{node}} = \text{Pick random node of tree} \end{array}
```



# RRT Growing in Empty Space

# RRT with obstacles and goal bias



# Path Planning with Rapidly-Exploring Random Trees (RRTs)

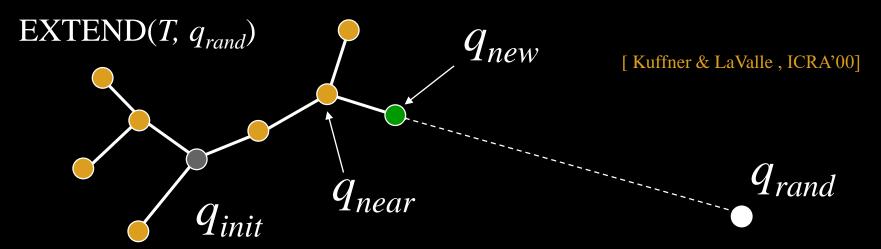
```
BUILD_RRT (q_{init}) {

T.init(q_{init});

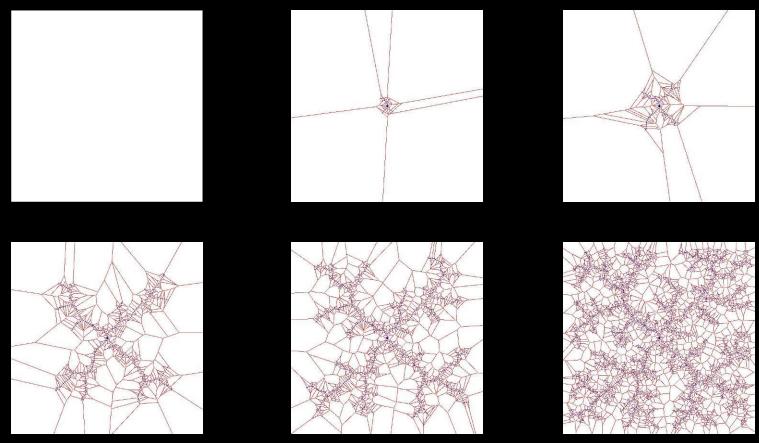
for k = 1 to K do

q_{rand} = RANDOM\_CONFIG();

EXTEND(T, q_{rand})
}
```



## RRTs bias exploration toward large Voronoi regions



http://msl.cs.uiuc.edu/rrt/gallery.html

What about obstacles?

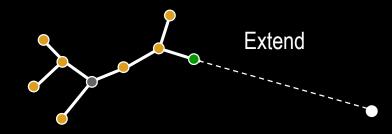
## RRT Goal Biasing

- In "pure" form RRTs are great at filling space, but we need a path!
- Need to bias RRTs toward goal to produce a path
  - When generating a random sample, with some probability pick the goal instead of a random node
  - This introduces another parameter
  - James Kuffner's experience is that 5-10% is the right choice

What happens if you set probability of sampling goal to 100%?

## **RRT Extension Types**

- RRT-Extend
  - Take one step toward a random sample



- RRT-Connect
  - Step toward random sample until it is either
    - Reached
    - You hit an obstacle

- Connect

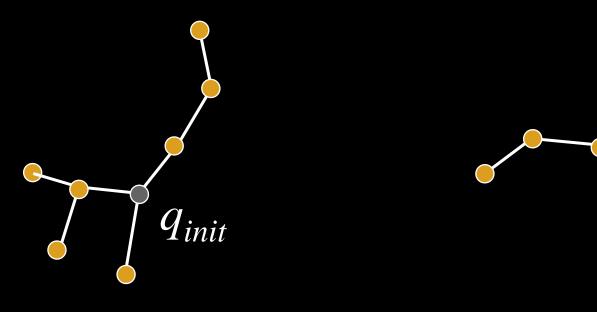
  Fixed step sizes except for last one
- Note that RRT-Connect is defined differently in some places online!

## **BiDirectional RRTs**

- BiDirectional RRT
  - Grow trees from both start and goal
  - Try to get trees to connect to each other
  - Trees can both use Extend or both use Connect or one use Extend and one Connect

- BiDirectional RRT with Connect for both trees is my favorite, I always try this first
  - This variant has only one parameter; the step size

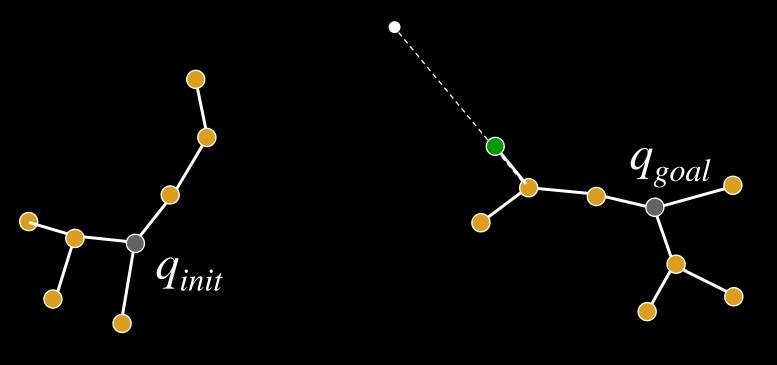
# Example of BiDirectional RRT



Connect Extend

 $q_{goal}$ 

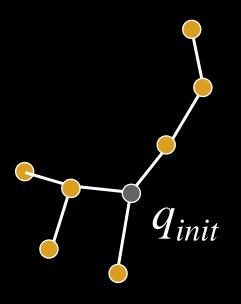
# 1) One tree grown using random target



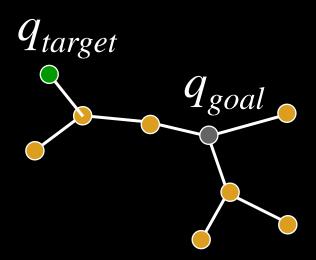
Connect

Extend

# 2) New node becomes target for other tree

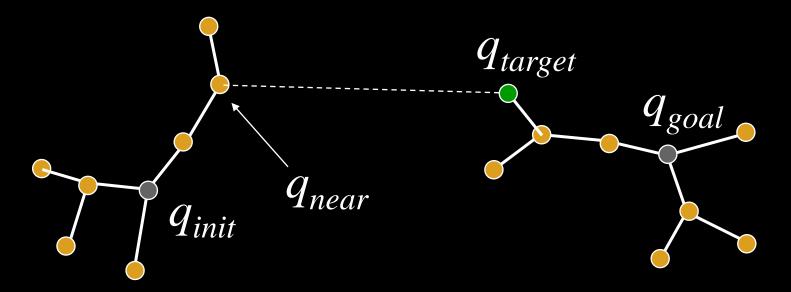


Connect



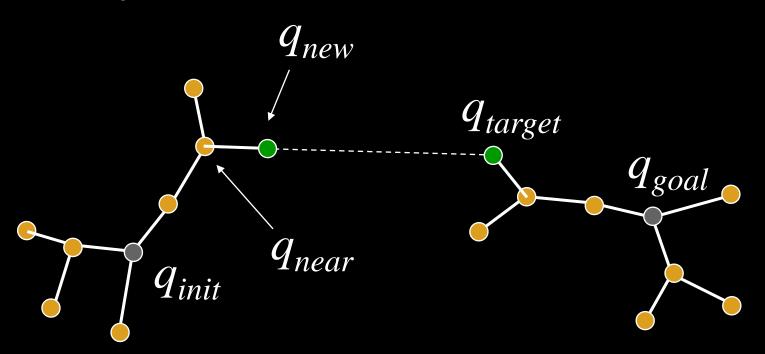
Extend

# 3) Calculate node "nearest" to target



Connect Extend

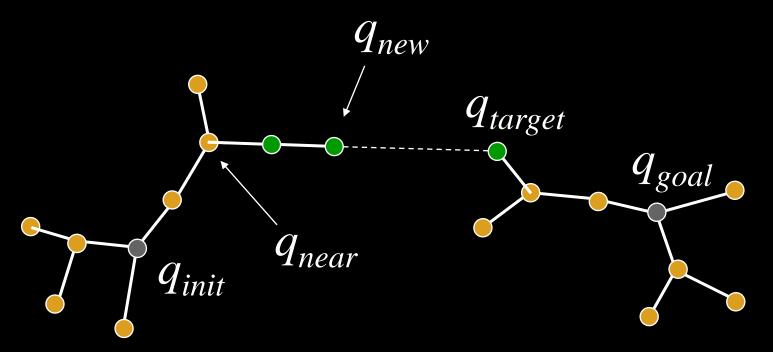
# 4) Try to add new collision-free branch



Connect

Extend

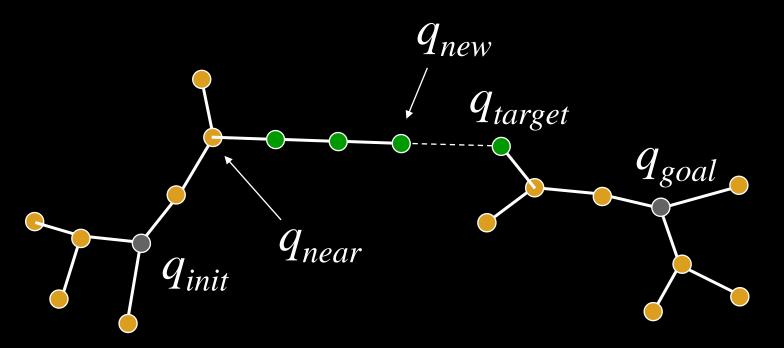
# 5) If successful, keep extending branch



Connect

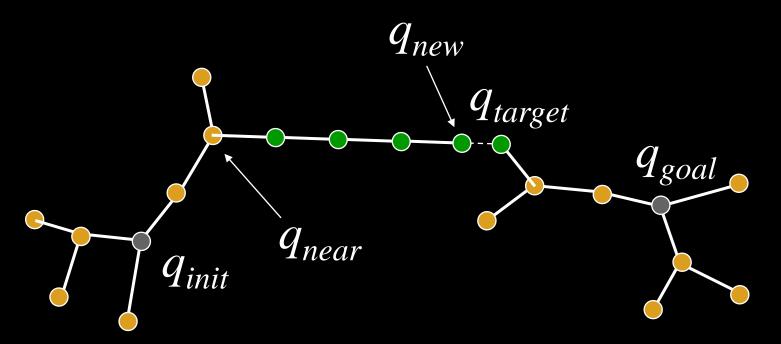
Extend

# 5) If successful, keep extending branch



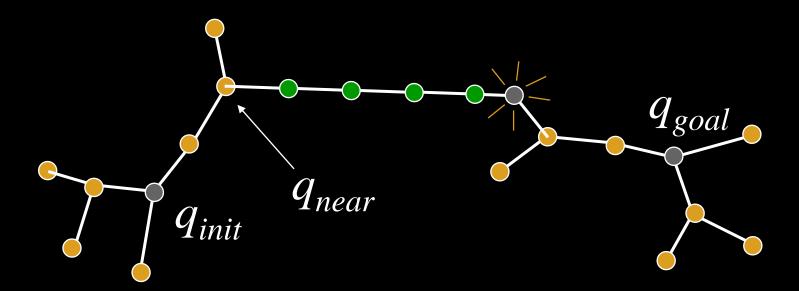
Connect Extend

# 5) If successful, keep extending branch



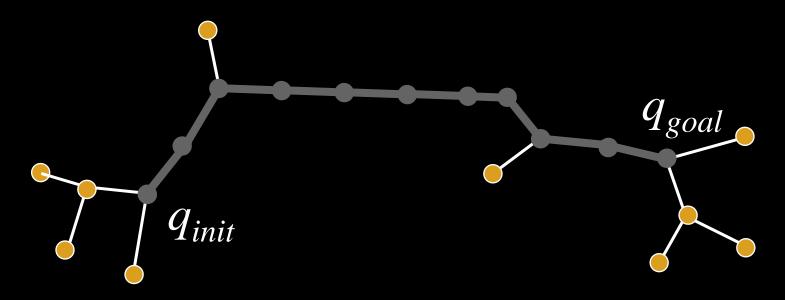
Connect Extend

### 6) Path found if branch reaches target



Connect Extend

# 7) Return path connecting start and goal



Connect Extend

### Tree Swapping and Balancing

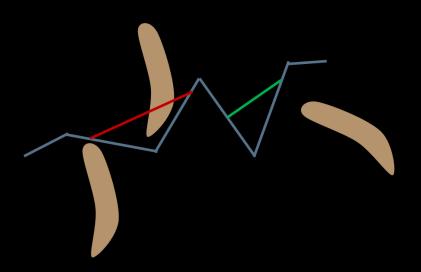
- Usually, swap trees after each iteration
- Some use Tree Balancing:

```
RDT BALANCED BIDIRECTIONAL (q_I, q_G)
       T_a.init(q_I); T_b.init(q_G);
       for i = 1 to K do
          q_n \leftarrow \text{NEAREST}(S_a, \alpha(i));
           q_s \leftarrow \text{STOPPING-CONFIGURATION}(q_n, \alpha(i));
          if q_s \neq q_n then
               T_a.add_vertex(q_s);
               T_a.add_edge(q_n, q_s);
               q'_n \leftarrow \text{NEAREST}(S_b, q_s);
               q'_s \leftarrow \text{STOPPING-CONFIGURATION}(q'_n, q_s);
               if q'_r \neq q'_r then
                   T_b.add_vertex(q'_s);
                   T_b.add_edge(q'_n, q'_s);
 12
 13
               if q'_s = q_s then return SOLUTION;
           if |T_b| > |T_a| then SWAP(T_a, T_b);
 14
 15 return FAILURE
```

- What is a situation where this would help performance?
- What is a situation where this would hurt performance?

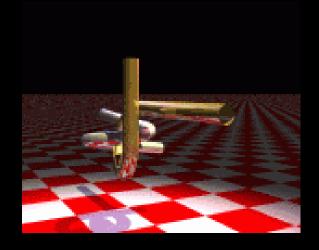
### Path Smoothing/Optimization

- RRTs produce notoriously bad paths
  - Not surprising since no consideration of path quality
- ALWAYS smooth/optimize the returned path
  - Many methods exists, e.g. shortcut smoothing (from previous lecture)



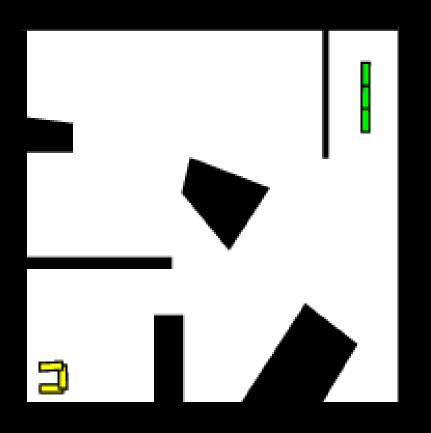
#### RRT Examples: The Alpha Puzzle

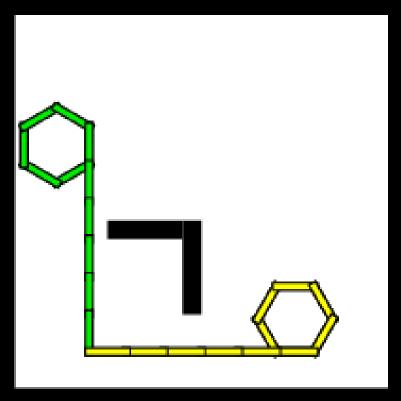
VERY hard 6DOF motion planning problem (long, winding narrow passage)



- "In 2001, it was solved by using a balanced bidirectional RRT, developed by James Kuffner and Steve LaValle. There are no special heuristics or parameters that were tuned specifically for this problem. On a current PC (circa 2003), it consistently takes a few minutes to solve" –RRT website
- RRT became famous in large part because it was able to solve this puzzle

### RRT Examples: Articulated Objects





### RRT Analysis

The limiting distribution of vertices:

THEOREM: X<sub>k</sub> converges to X with probability 1 as time goes to infinity

 $X_k$ : The RRT vertex distribution at iteration k

**X**: The distribution used for generating samples

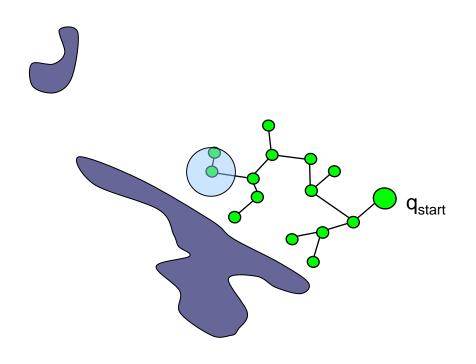
If using uniform distribution, tree nodes converge to the free space

#### Probabilistic Completeness

• Definition: A path planner is *probabilistically complete* if, given a solvable problem, the probability that the planner solves the problem goes to 1 as time goes to infinity.

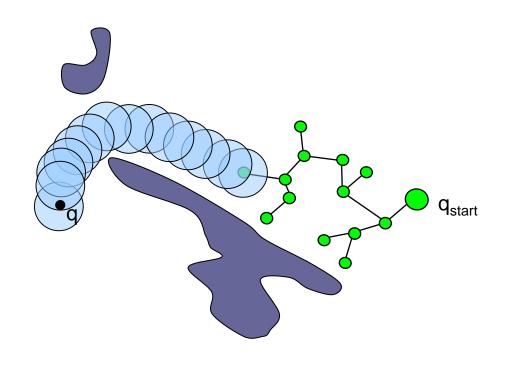
Will RRT explore the whole space?

# RRT P.C. Proof



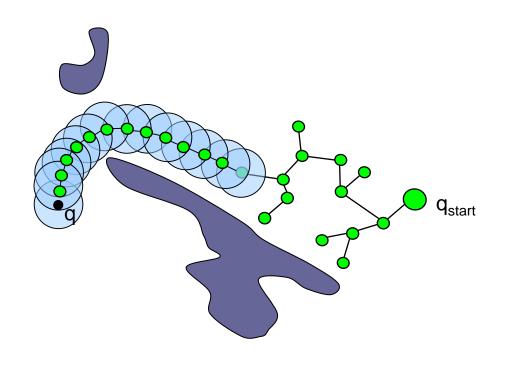
Kuffner and LaValle, ICRA, 2000

### RRT P.C. Proof



Kuffner and LaValle, ICRA, 2000

# RRT P.C. Proof



Kuffner and LaValle, ICRA, 2000

#### Probabilistic Completeness

• As the RRT reaches all of  $C_{free}$  to  $q_{start}$ , the probability that  $q_{rand}$  immediately becomes a new vertex approaches 1.

So, is RRT probabilistically complete?

### Sampling-Based Planning

- The good:
  - Provides fast feasible solution
  - Popular methods have few parameters
  - Works on practical problems
  - Works in high-dimensions
  - Works even with the wrong distance metric

### Sampling-Based Planning

- The bad:
  - No quality guarantees on paths\*
    - In practice: smooth/optimize path afterwards
  - No termination when there is no solution
    - In practice: set an arbitrary timeout
  - Probabilistic completeness is a weak property
    - Completeness in high-dimensions is impractical

<sup>\*</sup>Asymptotically-optimal sampling-based planners can make some guarantees

#### Homework

- Read LaValle Ch. 14-14.5
- Read "How to Read a Paper" guide (link on website)
- Make sure you've started HW2!