

Balanced Search Tree



Outline:

- Introduction
- Rebalancing
- Single Rotation
- Double Rotation
- Insertion
- Deletion

Introduction

- AVL Tree has been proposed by Adel'son-Vel'skii and Landis
- an AVL tree is a binary search tree such that for **any node** in the tree, the height of the left and right subtrees can differ by at most 1.

Introduction

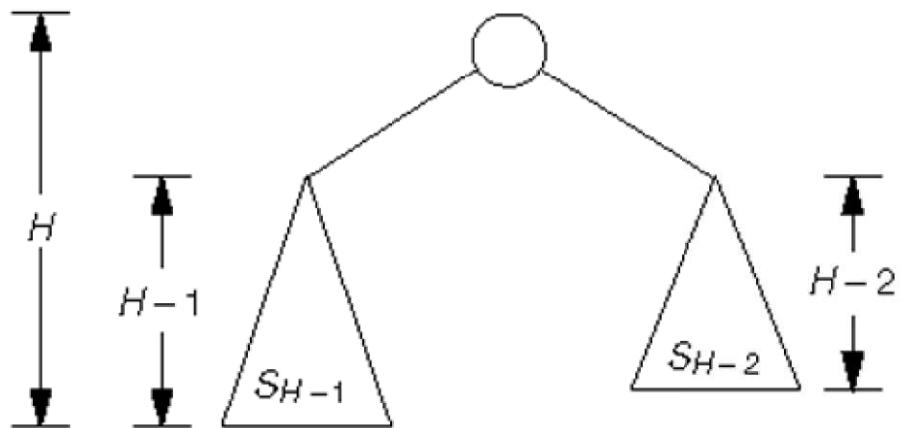
- **AVL Tree** is any tree with these conditions :
 - ✓ must be a binary tree
 - ✓ must be a binary search tree
 - ✓ must be a balanced tree

Introduction

- *Balance factors* are the differences between the heights of the left and right subtrees. For an AVL Tree, all balance factors should be +1, 0, or -1.
- AVL tree *approximates* the ideal tree (completely balanced tree)

Introduction

Minimum tree of height H



Introduction

▪ Properties:

- ✓ An update (insert or delete) in an AVL tree could destroy the balance. It must then be rebalanced before the operation can be considered complete.
- ✓ After an insertion, only nodes that are on the path from the insertion point to the root can have their balances altered.

Rebalancing

- Suppose the node to be rebalanced is **X**. There are four cases that we might have to fix :
 - ✓ Case 1: An insertion in the left subtree of the left child of **X**
 - ✓ Case 2: An insertion in the right subtree of the left child of **X**
 - ✓ Case 3: An insertion in the left subtree of the right child of **X**
 - ✓ Case 4: An insertion in the right subtree of the right child of **X**

Rebalancing

▪ Balancing Operations: Rotations

- ✓ **Case 1** and **case 4** are symmetric and requires the same operation for balance.
 - Cases 1 and 4 are handled by *single rotation*
- ✓ **Case 2** and **case 3** are symmetric and requires the same operation for balance.
 - Cases 2 and 3 are handled by *double rotation*

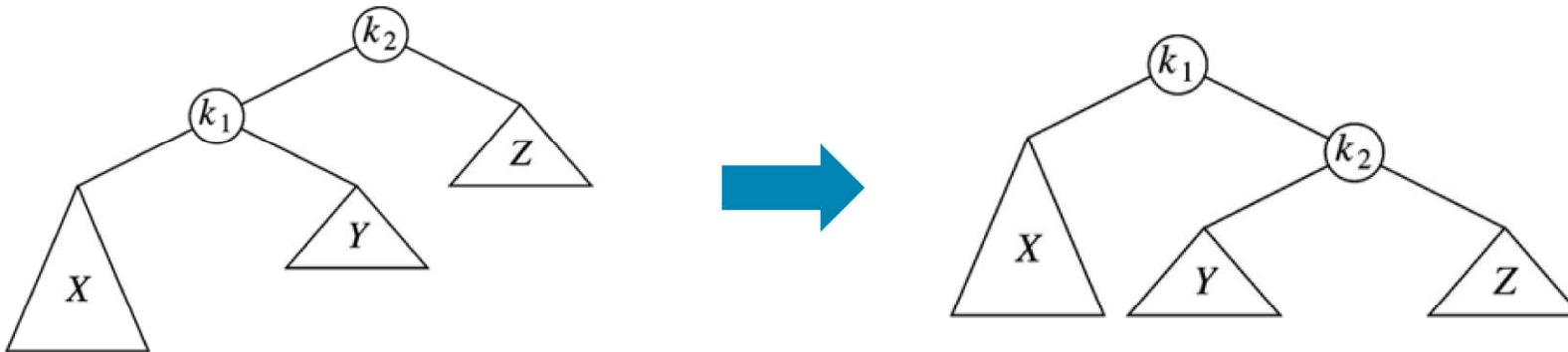
Single Rotation

- A single rotation switches the roles of the parent and child while maintaining the search order
- Rotate between a node and its child
 - ✓ Child becomes parent
 - ✓ Parent becomes right child in case 1, left child in case 4
- The result is a binary search tree that satisfies the AVL property

Single Rotation

Case 1:

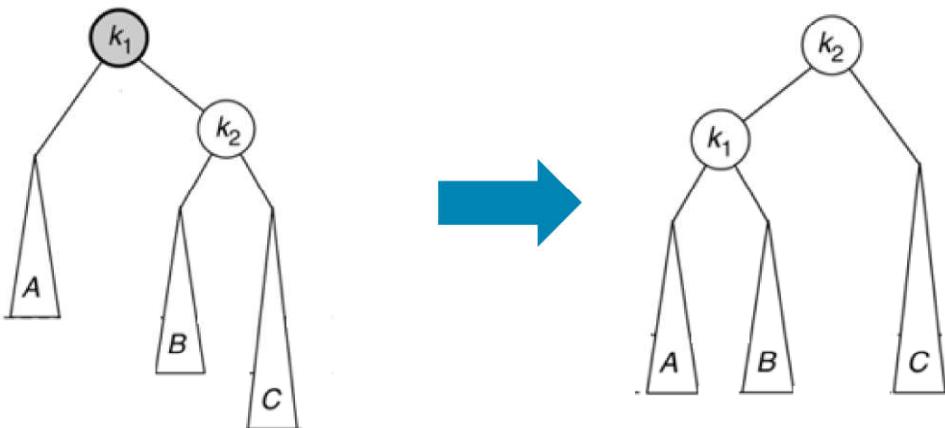
- ✓ Replace node k_2 by node k_1
- ✓ Set node k_2 to be right child of node k_1
- ✓ Set subtree Y to be left child of node k_2



Single Rotation

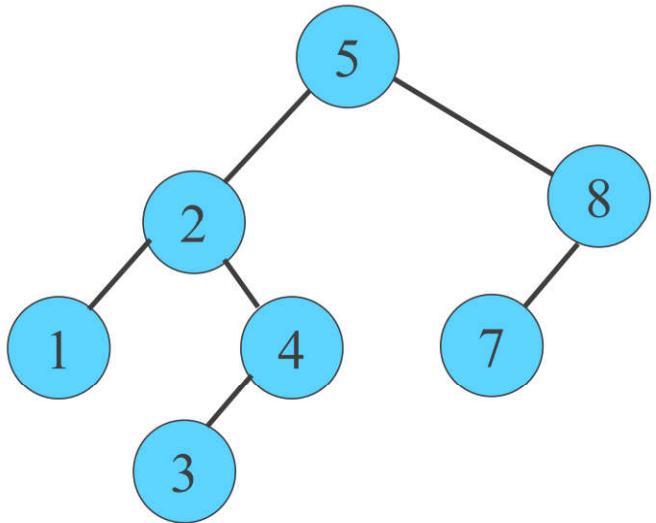
Case 4:

- ✓ Replace node k_1 by node k_2
- ✓ Set node k_1 to be left child of node k_2
- ✓ Set subtree B to be right child of node k_1



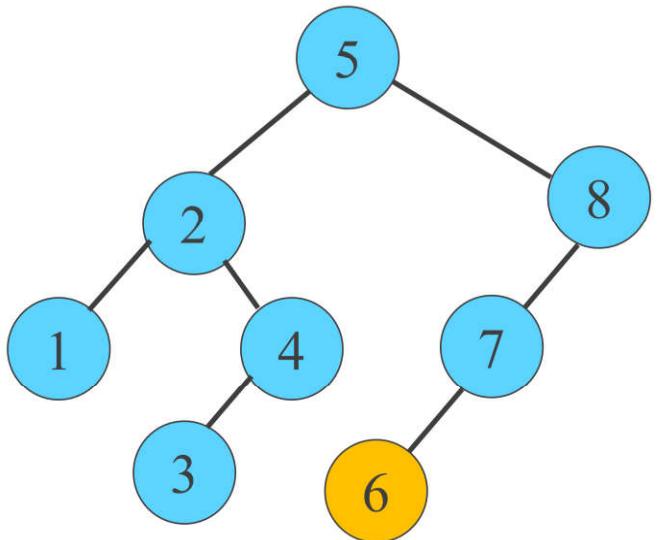
Single Rotation

- Example: Inserting node 6 in the AVL tree



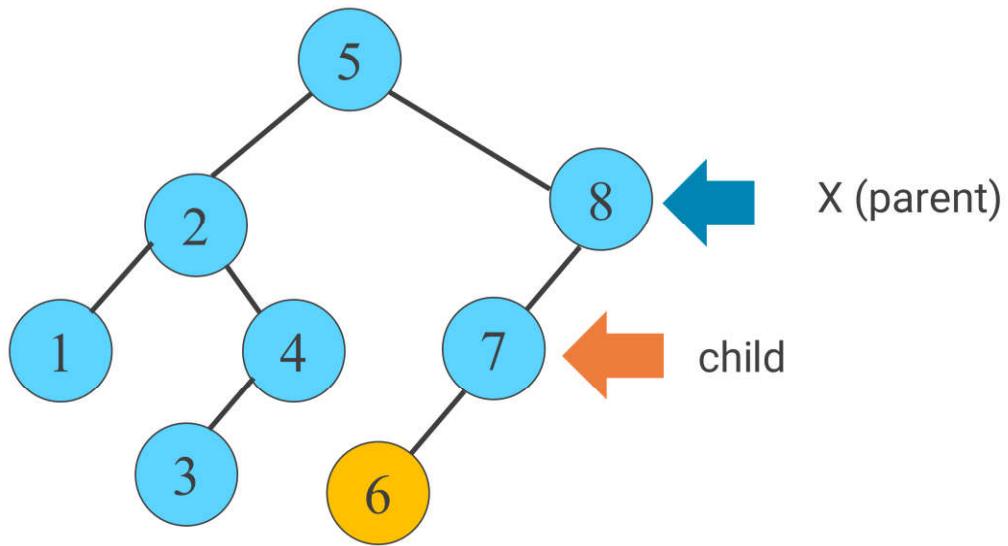
Single Rotation

- Solution:



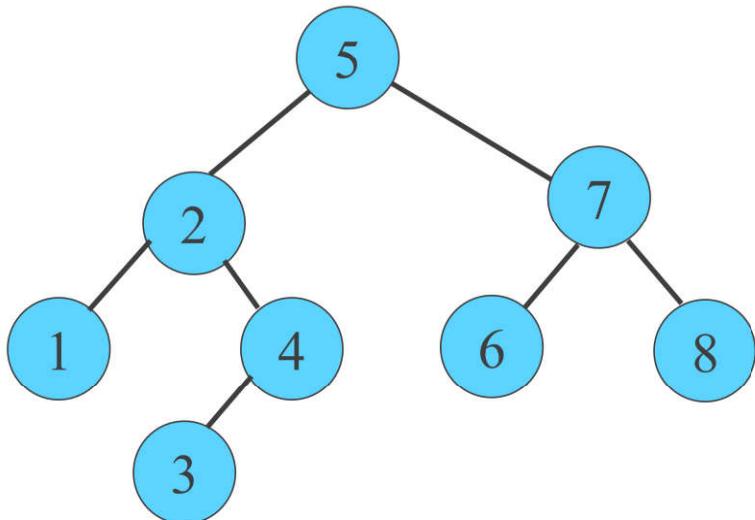
Single Rotation

- Solution:



Single Rotation

- Solution:

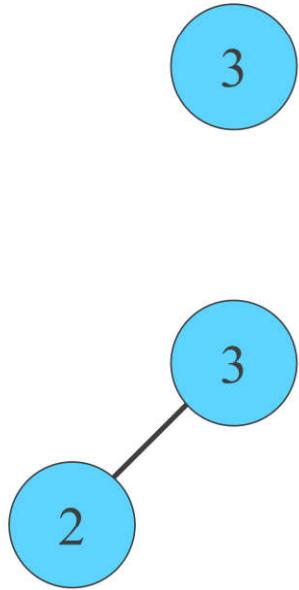


Single Rotation

- Example: Inserting 3, 2, 1, 4, 5, 6 and 7 in an empty AVL tree

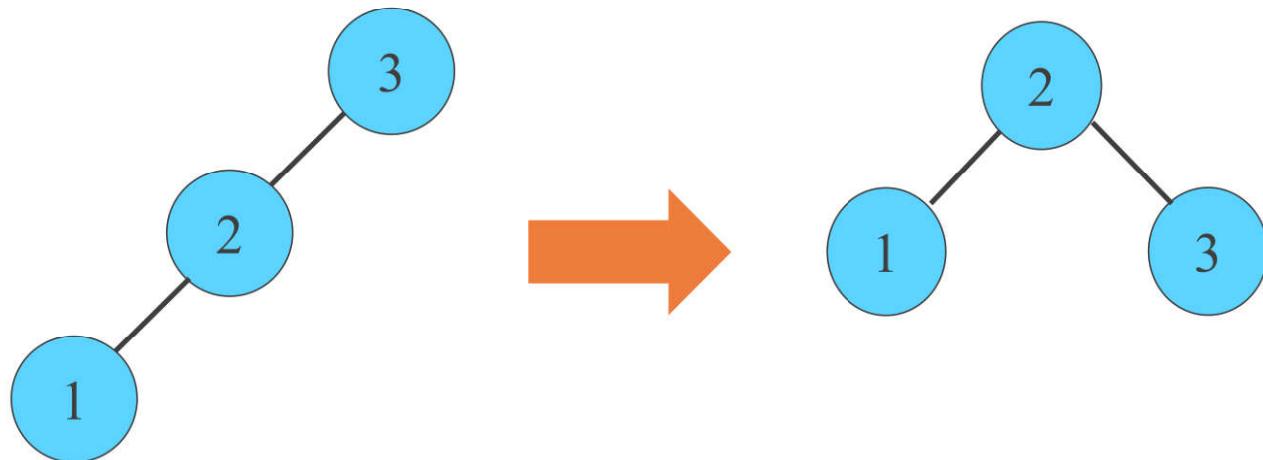
Single Rotation

- Solution:
 - ✓ inserting 3
 - ✓ inserting 2



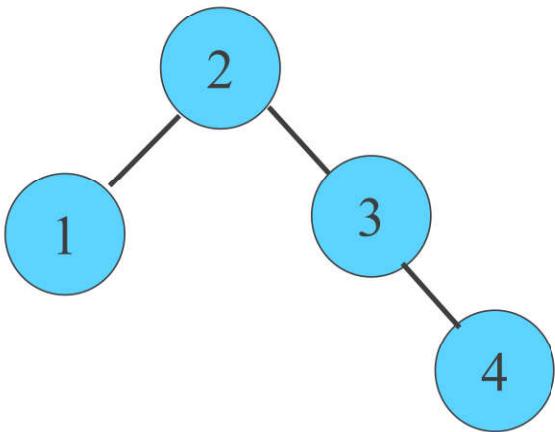
Single Rotation

- Solution:
 - ✓ inserting 1



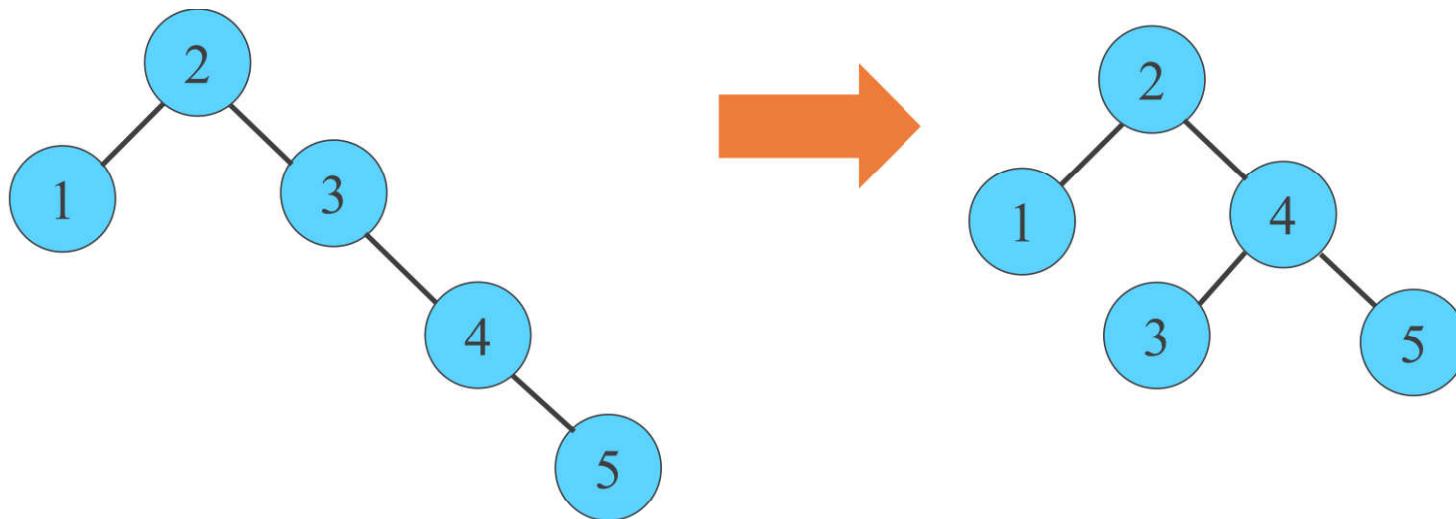
Single Rotation

- Solution:
 - ✓ inserting 4



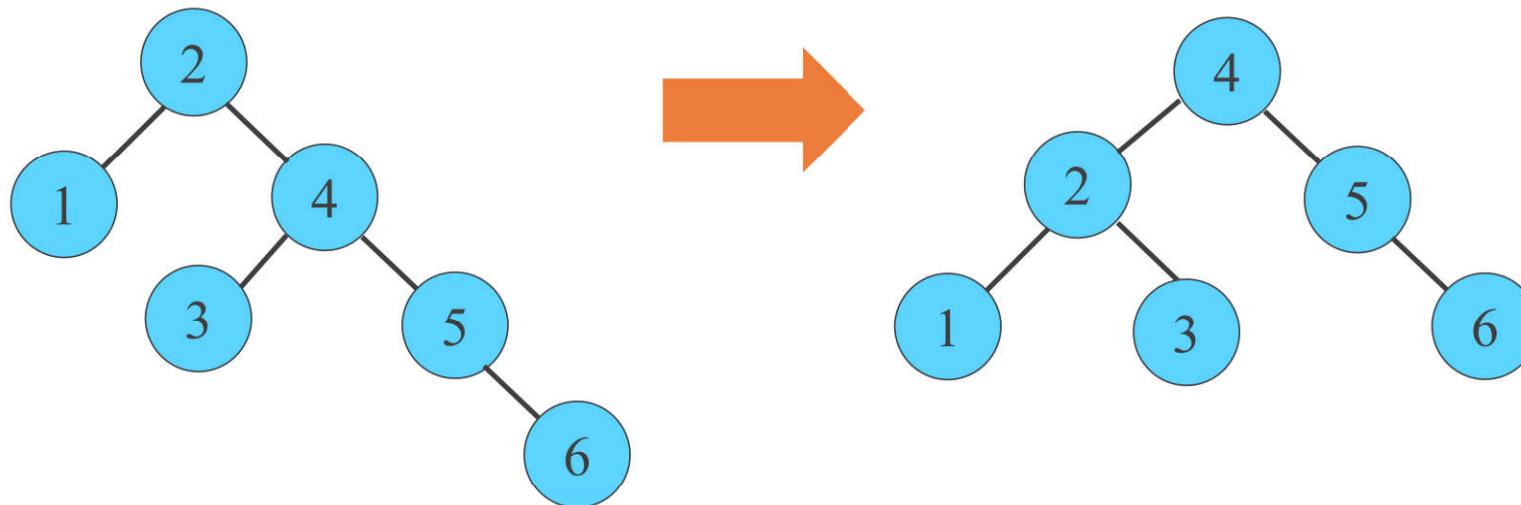
Single Rotation

- Solution:
 - ✓ inserting 5



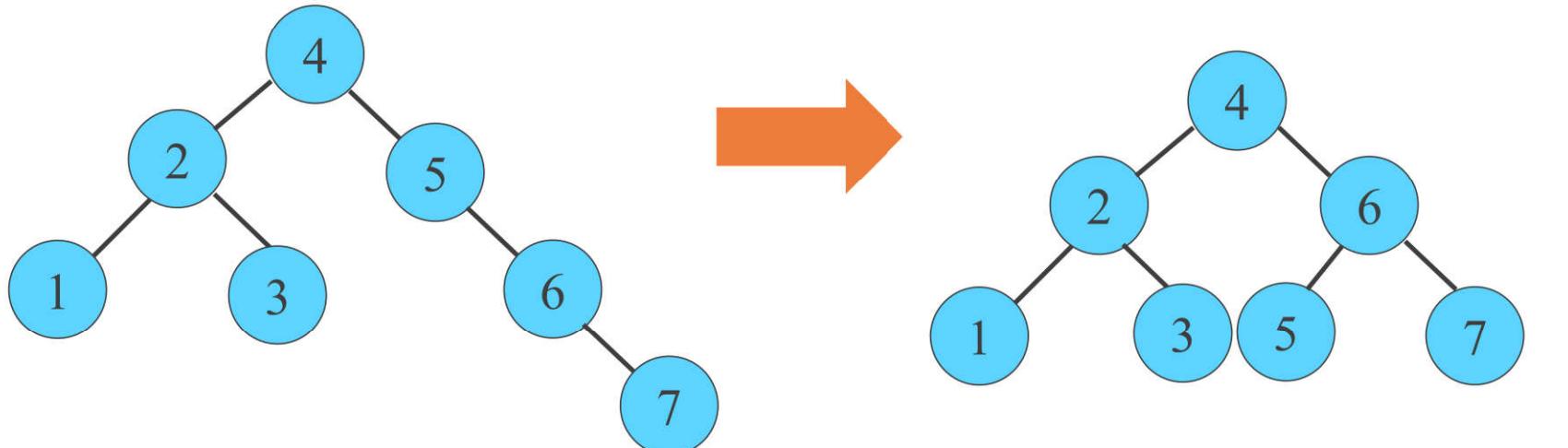
Single Rotation

- Solution:
 - ✓ inserting 6



Single Rotation

- Solution:
 - ✓ inserting 7



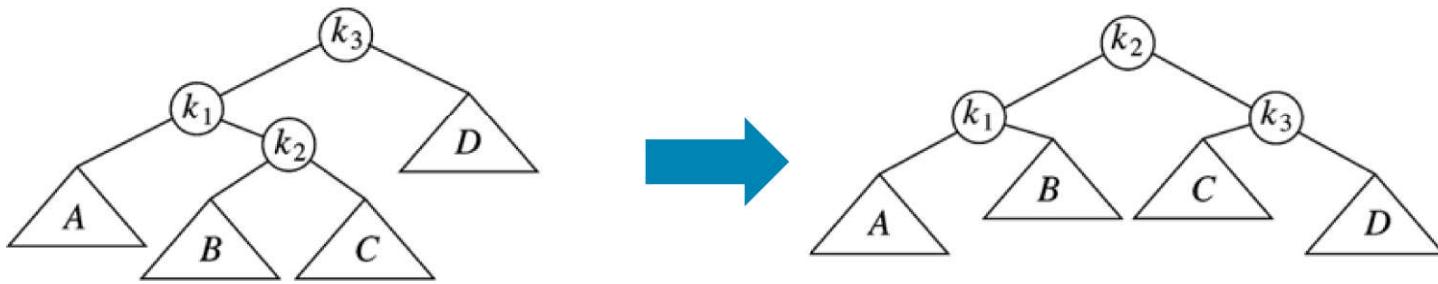
Double Rotation

- A Single rotation does not fix the case 2 and case 3
- These cases require a *double* rotation, involving three nodes and four subtrees

Double Rotation

▪ Case 2:

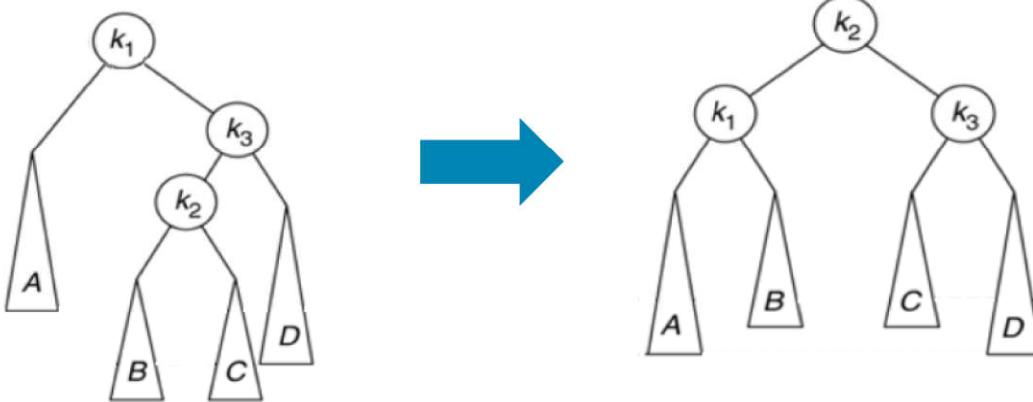
- ✓ Left-right double rotation to fix case 2
- ✓ First rotate between k_1 and k_2
- ✓ Then rotate between k_2 and k_3



Double Rotation

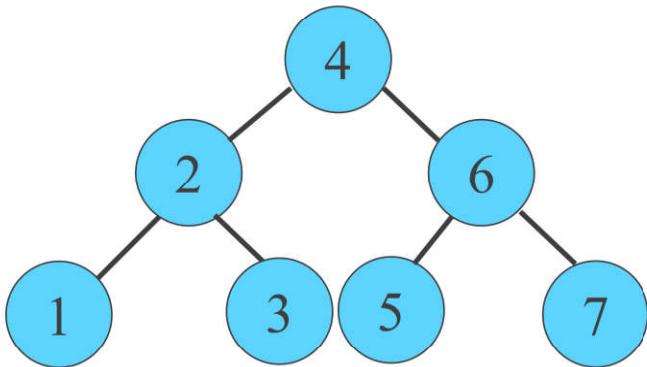
Case 3:

- ✓ Right-left double rotation to fix case 3
- ✓ First rotate between k_2 and k_3
- ✓ Then rotate between k_2 and k_1



Double Rotation

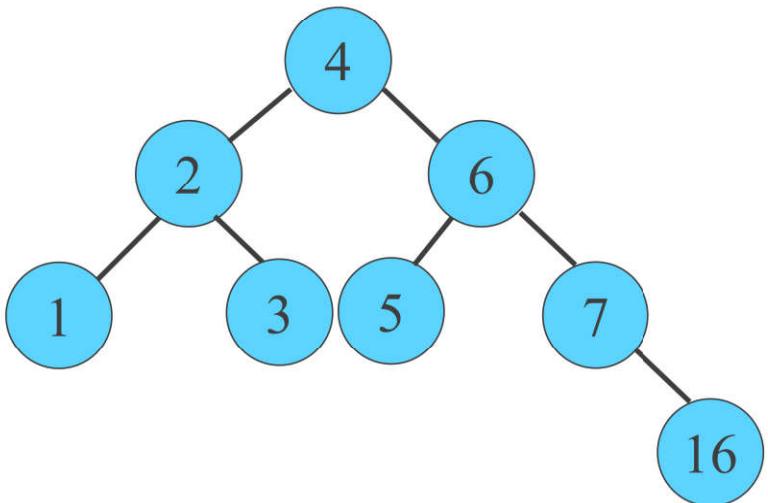
- Example: Inserting 16 and 15 in the AVL tree



Double Rotation

- Solution:

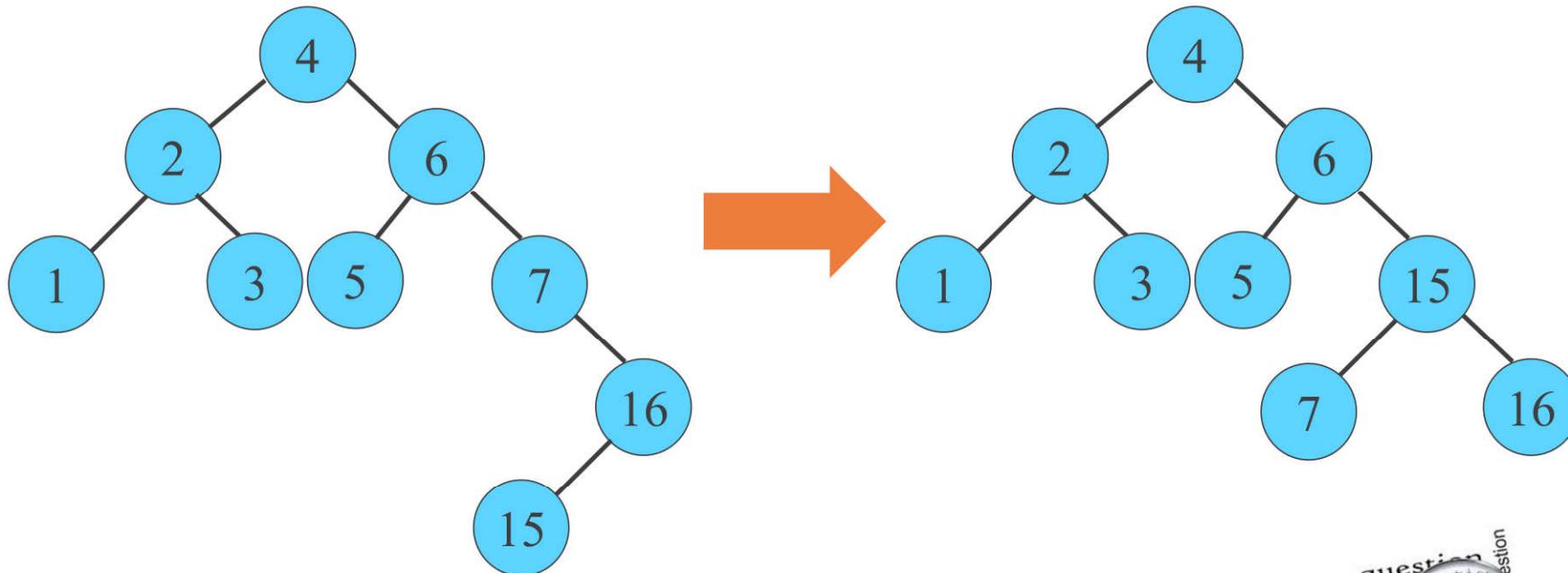
- ✓ Inserting 16



Double Rotation

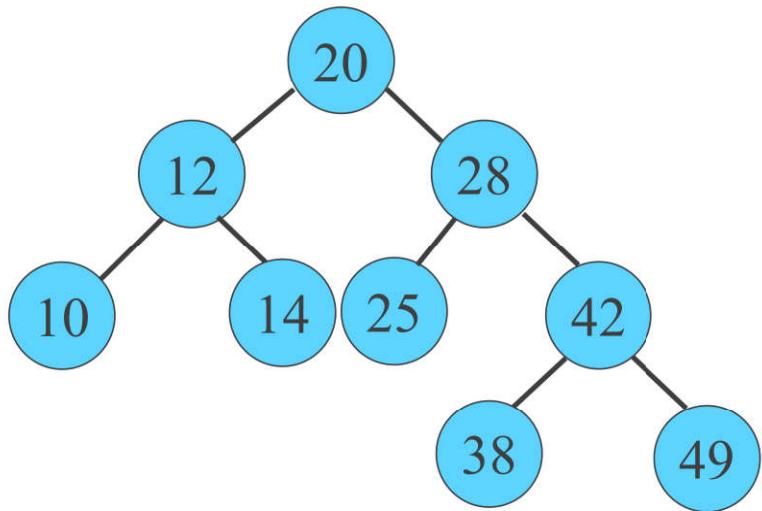
- Solution:

- ✓ Inserting 15



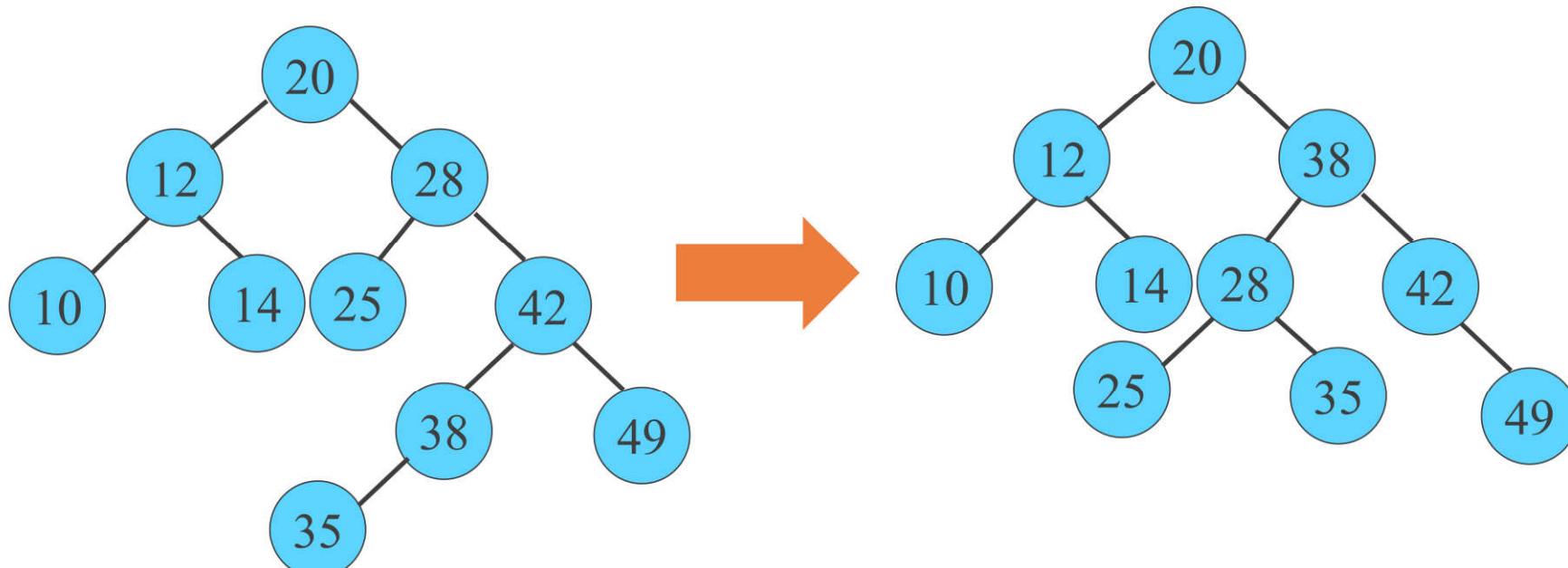
Double Rotation

- Example: Inserting 35 in the AVL tree



Double Rotation

- Solution:



Insertion

- Algorithm:

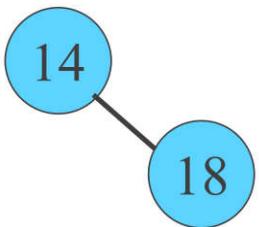
- ✓ insert a node **X** in the AVL tree using the Binary search tree insertion
- ✓ adjusts the balancing factor of all nodes in the path from the parent node of **X** to the root node
- ✓ If the balancing factor of all nodes is +1 or 0 or -1 then end the operation
else adjust the balancing factor

Insertion

- **Example:** Inserting 14, 18, 12, 9, 55, 5 and 13 in an empty AVL tree

Insertion

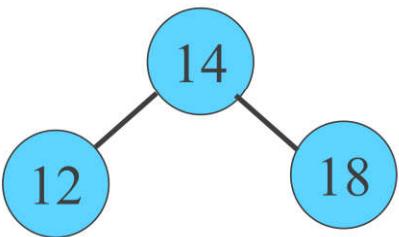
- Solution:
 - ✓ inserting 14
 - ✓ inserting 18



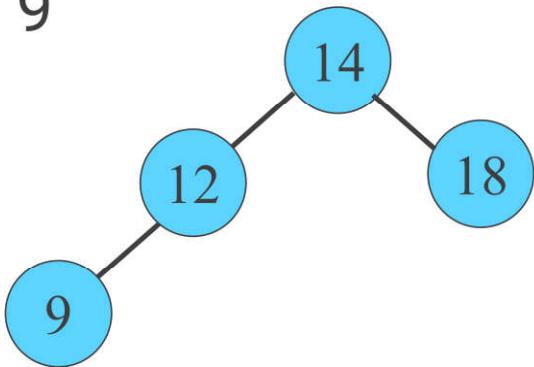
Insertion

- Solution:

✓ inserting 12

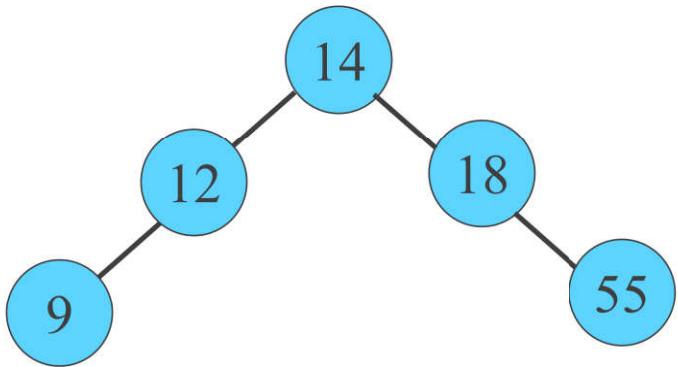


✓ inserting 9



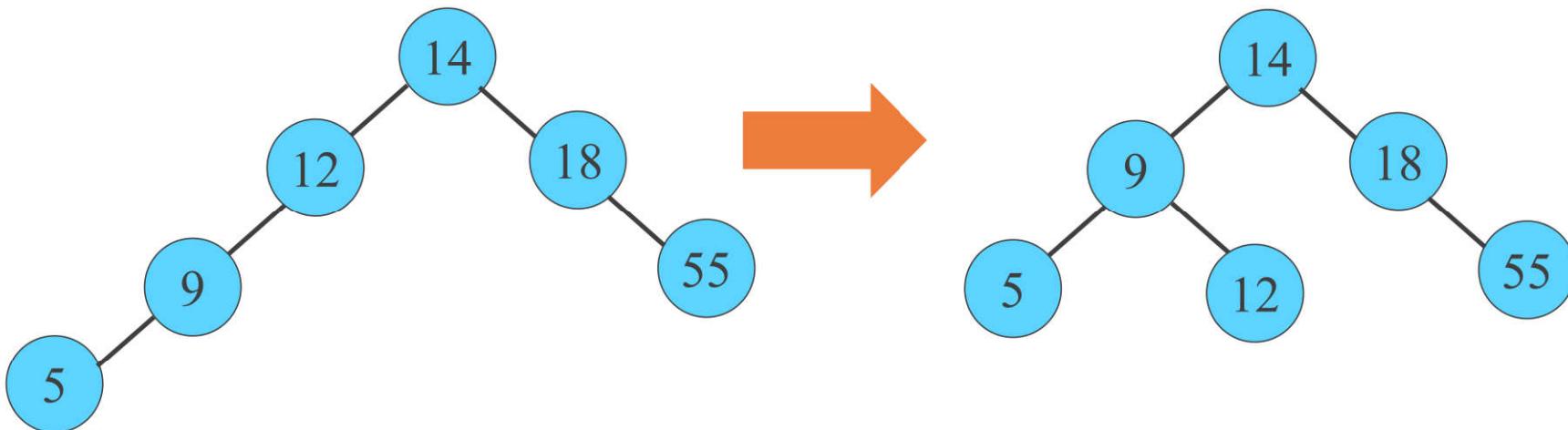
Insertion

- Solution:
 - ✓ inserting 55



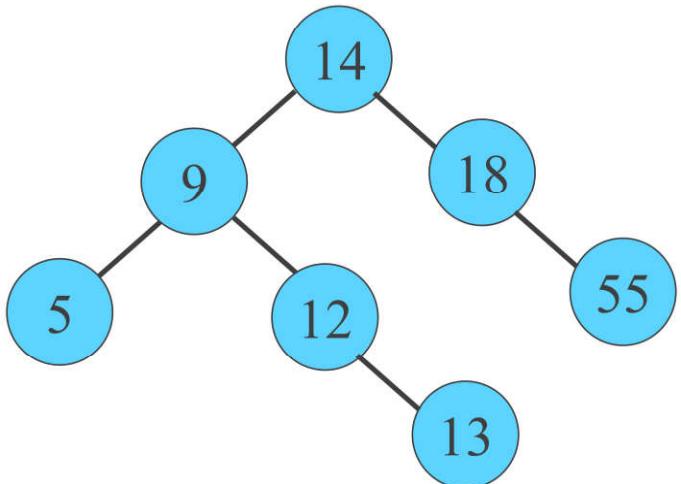
Insertion

- Solution:
 - ✓ inserting 5



Insertion

- Solution:
 - ✓ inserting 13



Deletion

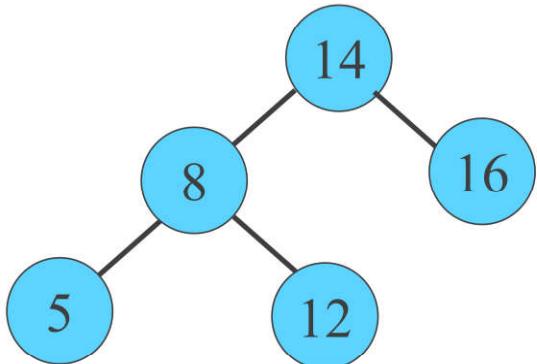
- Deletion is more complicated
- May need more than one rebalance on the path from deleted node to root
- Deletion of a node X from an AVL tree requires the same basic ideas, including single and double rotations, that are used for insertion

Deletion

- Algorithm:
 - ✓ delete a node **X** in the AVL tree using the Binary search tree deletion
 - ✓ adjusts the balancing factor of all nodes in the path from the parent node of **X** to the root node
 - ✓ If the balancing factor of all nodes is +1 or 0 or -1, end the operation
 - ✓ otherwise adjust the balancing factor

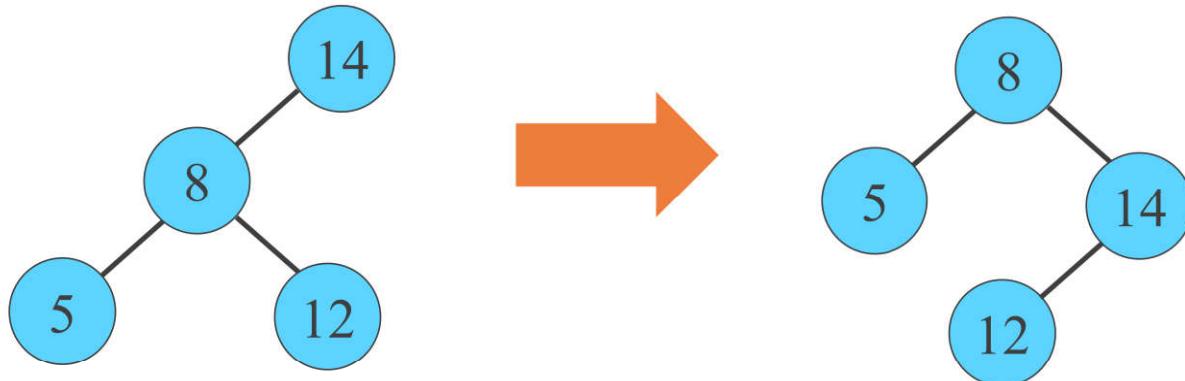
Deletion

- Example: Deleting 16 from the AVL tree



Deletion

- Solution:





Class Activity

Q & A



Formative Assessment