

June-2020 version for STM's STM32CubeIDE toolchain (newlib 3.0).

[See 2020 version for NXP MCUXpresso with newlib 2.5-3.1](#)

Background

For reliability, smaller embedded systems typically don't use dynamic memory, thus avoiding associated problems of leakage, fragmentation, and resulting out-of-memory crashes. In Nadler & Associates smaller embedded projects (and large avionics products), we've historically used ***no*** runtime dynamic memory.

Some tool chains (and especially some features of C++, the topic of a separate article) **require dynamic memory internally, even if your application uses no free storage calls**. The newlib C-runtime library (used in many embedded tool chains) internally uses it's own malloc-family routines. newlib maintains some internal buffers and requires some support for thread-safety.

This article explains how to use newlib safely in a FreeRTOS project with GNU toolchain.

Warning: This article discusses newlib version 3.0. Later versions **may be** different...

STM CubeMX Users - Beware

Newlib 3.0 is the only runtime library distributed in STM's STM32CubeIDE development environment. You can select *standard* or *reduced* for each of C and C++ (4 possible combinations). As of July-2019 (and still in June 2020! Unbelievable!), **Cube-generated projects using FreeRTOS** do not properly support malloc/free/etc and family, nor general newlib RTL reentrancy. **Your application will corrupt memory if it calls malloc/free/etc:**

- **directly**
- **via a newlib C-RTL function called by your application (for example `sprintf %f`), or**
- **via STM-provided HAL-LL code (for example STM's USB stack)**

This document explains how to fix this (follow the instructions at the bottom of this page).

newlib requirements

[GNU ARM Embedded Toolchain](#) distributions include a non-polluting reduced-size runtime library called newlib (or newlib-nano for the smallest variant). Unfortunately, newlib **internally** uses free storage (malloc/free) in startling places within the C runtime library.

Thus, newlib free storage routines get dragged in and used unexpectedly.

Ooops: Not MISRA-compliant, if you believe in that stuff.

The most common functions that bite unsuspecting embedded developers are `sprintf` using `%f`, `dtoa`, `ftoa`, `rand`, and `strtok`. [Here's a list of newlib functions with reentrancy support \(and hence using malloc\).](#)

newlib requires platform-specific support to:

- provide thread-safety,
- switch thread contexts, and
- obtain memory to be doled out and managed by malloc/free/etc (`sbrk`).

[Here's the detailed list of support functions required by newlib.](#)

If this support is properly implemented, newlib works well in a threaded environment like FreeRTOS. CubeMX does not properly set up these support functions for a FreeRTOS project.

See: [Bug Report to STM regarding CubeMX FreeRTOS projects](#) and the STM bug discussion below.

Example newlib use of malloc/malloc_r:

Placing a breakpoint within malloc_r (innermost malloc routine) shows the following:

- newlib 3.0 does **not** use malloc before main() is called (unlike earlier NXP/newlib 2.5 startup, which malloc'd 5k). Verified for all four newlib variants.
- With simple decimal output, sprintf does not call malloc
- For %f output, sprintf does 4 malloc totalling ~200 bytes (only the first time it is called per task/thread).
%f also requires proper linker arguments for float support.
- printf or similar function expected to do IO (as opposed to string operation) allocates an IO control structure of 428 bytes.

Tip: How can I see who is calling malloc ?

The basic steps are:

- create a wrapper function malloc_r,
- add the wrapper options to the linker, and
- place a breakpoint in malloc_r.

A complete example with instructions is included in the code below.

newlib under the hood - concurrency

newlib maintains information it needs to support each separate context (thread/task/ISR) in a **reentrancy structure**. This includes things like a thread-specific errno, thread-specific pointers to allocated buffers, etc. The active reentrancy structure is pointed at by global pointer `_impure_ptr`, which initially points to a statically allocated structure instance. If you use no threading or tasks, and you don't use malloc/free/etc or any of the reentrancy-dependent functions in multiple execution contexts, nothing more is required to use newlib in your application (newlib will maintain its required info in the single static structure). If your application or any library you use requires malloc/free/etc or any of the reentrancy-dependent functions in multiple contexts, newlib requires:

- concurrency protection for malloc/free/etc. The free storage pool will be corrupted if multiple threads call into these functions concurrently! To prevent reentrant execution of malloc/free/etc routines [newlib requires hook procedures __malloc_lock/unlock](#). If an RTOS-based application does not replace newlib's **complete** internal malloc family (FreeRTOS does not), `_malloc_lock/unlock` must be provided for thread safety.
- multiple reentrancy structures (one per context), and a mechanism to create, initialize, and cleanup these structures, plus switching `_impure_ptr` to point at the correct reentrancy structure each time the context changes.

WTF? Why are newlib's dtoa, ftoa, and sprintf calling malloc???

To quote Steele and White, *Isn't it a pain when you ask a computer to divide 1.0 by 10.0 and it prints 0.0999999?* To summarize and over-simplify a very long story, in 1990 my friend Will Clinger published a seminal paper giving a solution to this problem, but Will's solution required arbitrary precision and hence storage. See [Will's retrospective on his seminal 1990 paper and subsequent developments](#). David Gay published an improved implementation (based on Will's paper), which was extremely widely adopted. newlib uses David Gay's code for dtoa and ftoa (using malloc), and sprintf uses dtoa and ftoa. [Keith Packard's picolibc implements the no-malloc precise Ryū algorithm](#). There has been some recent discussion on the newlib mailing list including [details of Gay's algorithm and maybe incorporating no-malloc Ryū algorithm from picolibc into newlib](#).

Really now, aren't you sorry you asked me why?

FreeRTOS support for newlib

FreeRTOS provides support for newlib's context management. In `FreeRTOSconfig.h`, add:

```
#define configUSE_NEWLIB_REENTRANT 1 // Required for thread-safety of newlib sprintf, strtok, etc...
```

With this option FreeRTOS does the following (in `task.c`):

- For each task, allocate and initialize a newlib reentrancy structure in the task control block (TCB). This adds 96* bytes overhead per task (* size hugely dependant on newlib build options), plus anything else newlib allocates as needed.
- Each task switch, set `_impure_ptr` to point to the newly active task's reentrancy structure.
- On task destruction, clean up the reentrancy structure (help newlib free any associated memory).

Careful: Depending on how your supplier built newlib, and whether you are using nano, `configUSE_NEWLIB_REENTRANT` can chew up considerable memory! Check carefully if this is OK for your application!

FreeRTOS memory management

FreeRTOS internally uses its own memory management scheme and API (implemented by your choice of `heapxx.c` routines). Some FreeRTOS applications only use FreeRTOS-provided memory management. Unfortunately many libraries use the standard "C" `malloc`-family routines internally, including newlib itself (for example STM's USB stack LL and HAL routines). For these cases I've implemented the FreeRTOS memory management API on top of the "C" standard (using newlib) in the module `heap_useNewlib.c` below.

Bugs in projects generated by STM's CubeMX

Note as of spring 2020: ST has attempted to solve some of problems discussed below, but rather than just redistributing `heap_useNewlib` (as NXP and others have done), ST has tried to redo the implementation - and failed to do it correctly! Best you delete their unfortunate failed attempts and use proven working code provided here.

From earlier versions: There are a number of problems in projects generated with STM's CubeMX tool, which lead to memory corruption and other unfortunate happenings:

1. STM does not provide `__malloc_lock/unlock`. Hence memory management is not thread-safe.
2. *In a stunning bit of coding malpractice, STM's USB stack calls `malloc` from within an ISR* (see for example function `USBD_CDC_Init`). Normally thread-safety is provided by suspending task-switching during memory management. Because of this incompetent use of `malloc` inside an ISR, interrupts need to be suspended during memory management. Unfortunately this can greatly increase interrupt latency. *ToDo: Verify priority set by `taskENTER_CRITICAL_FROM_ISR` actually blocks USB interrupt.*
3. newlib's free storage requires an external implementation of `sbrk` to provide the heap storage then doled out by the `malloc`-family (see [newlib sbrk requirements](#)).
The `sbrk` implementation provided by STM cannot work with FreeRTOS.
Any `malloc` request after the scheduler starts will fail.
4. STM did not enable (in `FreeRTOSconfig.h`):

```
#define configUSE_NEWLIB_REENTRANT 1 // Required for thread-safety of newlib sprintf, strtok, etc...
```

Using newlib safely with FreeRTOS - Possible Approaches

If you are certain your application does not use any newlib functions or any library that internally use the `malloc`-family and/or depend on thread-specific reentrant context, you could do nothing. But, are you *really* sure??? What about the libraries you use (for example STM's USB stack LL and HAL routines)? For some hints, see this [list of newlib functions with reentrancy support](#). The only really safe way to preclude accidental use is to provide HCF stubs for every one of these functions. Alternatively you can provide a linker -wrap command for each forbidden function but don't implement the wrappers, so link fails if forbidden routines are ever referenced.

You really need to be sure!

To avoid using newlib's `printf` and dragging in newlib reentrancy components, there are many cut-down `printf` implementations available (that do not use `malloc`). Won't help if you're using `dtoa` or `strtok` or others! Example light-weight `printf` implementations:

- [Mario Viara's light-weight printf](#) (optional floating point and reentrancy), used in *MCU on Eclipse Processor Expert*.
- [mpaland printf](#) (optional floating point).

- printf-stdarg.c distributed in [the FreeRTOS Lab TCP/IP example](#); this one only uses stack storage but does not implement floating point.
- https://github.com/ksstech/support_common.

You must ensure you don't accidentally use any newlib facilities requiring reentrancy support and/or malloc-family in multiple tasks.

Another option is wrap newlib's malloc-family to use FreeRTOS free storage (ie heap_4.c), and specify newlib support for FreeRTOS. Tell the linker to wrap all newlib's malloc-family functions (using -Xlinker --wrap=malloc etc.), and provide a wrapper function that calls the FreeRTOS functions. I tried that, but newlib's printf family uses realloc, which is not supported in FreeRTOS heap implementations.

In the end (thanks to Richard Damon for encouraging this approach), I implemented the FreeRTOS memory API on top of newlib's malloc family, and provided all the hooks newlib's malloc family requires.

Using newlib safely with FreeRTOS - Recommended Solution Details

If your application needs a complete malloc family implementation, or you are using *any* newlib functions that require malloc (for example printf family or strtok), do the following (I've provided an implementation below):

- Implement the hooks required by newlib (sbrk, __malloc_lock/unlock). Make sure your linker file matches the sbrk implementation!
- Provide a heap implementation that implements the FreeRTOS memory API using the malloc family of newlib.

To use the implementation I've provided in your project:

- Exclude from all builds any current FreeRTOS heap implementation, typically something like: Middlewares\Third_Party\FreeRTOS\Source\portable\MemMang\heap_4.c
- Exclude from all builds any current sbrk implementation. In older versions of RubeMX, ST generated a stand-alone sbrk.c module. Later versions hide it in syscalls.c to obfuscate their messes. Get rid of it!
- Add the module heap_useNewlib_ST.c I've provided.
- If in multiple tasks your application needs a complete sprintf implementation, strtok, dtoa>, or other newlib functions requiring reentrancy support, or you're not *really* sure... Configure FreeRTOS for newlib support. In FreeRTOSconfig.h, add the line:
#define configUSE_NEWLIB_REENTRANT 1
In RubeMX FreeRTOS options, select configUSE_NEWLIB_REENTRANT + deselect "FW pack heap file".
- If are sure you're only going to access newlib's reentrant routines from a single FreeRTOS task and want to skip configUSE_NEWLIB_REENTRANT (with its attendant overhead)? Ensure malloc is never called except during startup and from the designated single task. Verify with a check in malloc_lock (sorry, I did not provide example code).

[Note: Latest code and further documentation is on Github](#)

Hide Code...

heap_useNewlib.c

```

1  /**
2   * \file heap_useNewlib_ST.c
3   * \brief Wrappers required to use newlib malloc-family within FreeRTOS.
4   *
5   * \par Overview
6   * Route FreeRTOS memory management functions to newlib's malloc family.
7   * Thus newlib and FreeRTOS share memory-management routines and memory pool,
8   * and all newlib's internal memory-management requirements are supported.
9   *

```

```

10  * \author Dave Nadler
11  * \date 20-August-2019
12  * \version 27-Jun-2020 Correct "FreeRTOS.h" capitalization, commentary
13  * \version 24-Jun-2020 commentary only
14  * \version 11-Sep-2019 malloc accounting, comments, newlib version check
15  *
16  * \see http://www.nadler.com/embedded/newlibAndFreeRTOS.html
17  * \see https://sourceware.org/newlib/libc.html#Reentrancy
18  * \see https://sourceware.org/newlib/libc.html#malloc
19  * \see https://sourceware.org/newlib/libc.html#index-\_005f\_005fenv\_005flock
20  * \see https://sourceware.org/newlib/libc.html#index-\_005f\_005fmalloc\_005flock
21  * \see https://sourceforge.net/p/freertos/feature-requests/72/
22  * \see http://www.billgatliff.com/newlib.html
23  * \see http://wiki.osdev.org/Porting\_Newlib
24  * \see http://www.embecosm.com/appnotes/ean9/ean9-howto-newlib-1.0.html
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55  * ANY THEORY OF LIABILITY, WHETHER IN CONTRACT, STRICT LIABILITY, OR TORT
56  * (INCLUDING NEGLIGENCE OR OTHERWISE) ARISING IN ANY WAY OUT OF THE USE OF THIS
57  * SOFTWARE, EVEN IF ADVISED OF THE POSSIBILITY OF SUCH DAMAGE.
58  */
59
60 // =====
61 // ===== Configuration =====
62 // These configuration symbols could be provided by from build...
63 #define STM_VERSION // Replace sane LD symbols with STM CubeMX's poor standard e
64 #define ISR_STACK_LENGTH_BYTES (configISR_STACK_SIZE_WORDS*4) // bytes to reserve
65 // ===== Configuration =====
66 // =====
67
68
69 #include <stdlib.h> // maps to newlib...
70 #include <malloc.h> // mallinfo...
71 #include <errno.h> // ENOMEM
72 #include <stdbool.h>
73 #include <stddef.h>
74

```

```

75 #include "newlib.h"
76 #if ((__NEWLIB__ == 2) && (__NEWLIB_MINOR__ < 5)) || ((__NEWLIB__ == 3) && (__NEWLIB_MINOR__ < 1))
77     #warning "This wrapper was verified for newlib versions 2.5 - 3.1; please ensure you are using the correct version"
78 #endif
79
80 #include "FreeRTOS.h" // defines public interface we're implementing here
81 #if !defined(configUSE_NEWLIB_REENTRANT) || (configUSE_NEWLIB_REENTRANT!=1)
82     #warning "#define configUSE_NEWLIB_REENTRANT 1 // Required for thread-safety of newlib"
83     // If you're *REALLY* sure you don't need FreeRTOS's newlib reentrancy support,
84 #endif
85 #include "task.h"
86
87 // =====
88 // External routines required by newlib's malloc (sbrk/_sbrk, __malloc_lock/unlock)
89 // =====
90
91 // Simplistic sbrk implementations assume stack grows downwards from top of memory
92 // and heap grows upwards starting just after BSS.
93 // FreeRTOS normally allocates task stacks from a pool placed within BSS or DATA
94 // Thus within a FreeRTOS task, stack pointer is always below end of BSS.
95 // When using this module, stacks are allocated from malloc pool, still always pointing to
96 // current unused heap area...
97
98 // Doesn't work with FreeRTOS: STM CubeMX 2018-2019 Incorrect Implementation
99 #if 0
100     caddr_t _sbrk(int incr)
101     {
102         extern char end asm("end"); // From linker: lowest unused RAM address, just after BSS
103         static char *heap_end;
104         char *prev_heap_end;
105         if (heap_end == 0) heap_end = &end;
106         prev_heap_end = heap_end;
107         if (heap_end + incr > stack_ptr) // Fails here: always true for FreeRTOS
108         {
109             errno = ENOMEM; // ...so first call inside a FreeRTOS task lands here
110             return (caddr_t) -1;
111         }
112         heap_end += incr;
113         return (caddr_t) prev_heap_end;
114     }
115 #endif
116
117 register char * stack_ptr asm("sp");
118
119 #ifdef STM_VERSION // Use STM CubeMX LD symbols for heap+stack area
120     // To avoid modifying STM LD file (and then having CubeMX trash it), use available symbols
121     // Unfortunately STM does not provide standardized markers for RAM suitable for heap/stack
122     // STM CubeMX-generated LD files provide the following symbols:
123     //     end /* aligned first word beyond BSS */
124     //     _estack /* one word beyond end of "RAM" Ram type memory, for STM32F429
125     // Kludge below uses CubeMX-generated symbols instead of sane LD definitions
126     #define __HeapBase end
127     #define __HeapLimit _estack // In K64F LD this is already adjusted for ISR stack
128     static int heapBytesRemaining;
129     // no DRN HEAP_SIZE symbol from LD... // that's (&__HeapLimit)-(&__HeapBase)
130     uint32_t TotalHeapSize; // publish for diagnostic routines; filled in first task
131 #else
132     // Note: DRN's K64F LD provided: __StackTop (byte beyond end of memory), __StackLimit
133     // __HeapLimit was already adjusted to be below reserved stack area.
134     extern char HEAP_SIZE; // make sure to define this symbol in linker LD command
135     static int heapBytesRemaining = (int)&HEAP_SIZE; // that's (&__HeapLimit)-(&__HeapBase)
136 #endif
137
138
139 #ifdef MALLOCS_INSIDE_ISR // STM code to avoid malloc within ISR (USB CDC stack)

```



```

140 // We can't use vTaskSuspendAll() within an ISR.
141 // STM's stunningly bad coding malpractice calls malloc within ISRs (for exam
142 // So, we must just suspend/resume interrupts, lengthening max interrupt res
143 #define DRN_ENTER_CRITICAL_SECTION(_usis) { _usis = taskENTER_CRITICAL_FROM_ISR(_usis);
144 #define DRN_EXIT_CRITICAL_SECTION(_usis) { taskEXIT_CRITICAL_FROM_ISR(_usis);
145 #else
146 #define DRN_ENTER_CRITICAL_SECTION(_usis) vTaskSuspendAll(); // Note: safe to
147 #define DRN_EXIT_CRITICAL_SECTION(_usis) xTaskResumeAll(); // Note: safe to
148 #endif
149
150 #ifndef NDEBUG
151     static int totalBytesProvidedBySBRK = 0;
152 #endif
153 extern char __HeapBase, __HeapLimit; // symbols from linker LD command file
154
155 // Use of vTaskSuspendAll() in _sbrk_r() is normally redundant, as newlib malloc
156 // __malloc_lock before calling _sbrk_r(). Note vTaskSuspendAll/xTaskResumeAll s
157
158 ///! _sbrk_r version supporting reentrant newlib (depends upon above symbols defini
159 void * _sbrk_r(struct _reent *pReent, int incr) {
160     #ifndef MALLOCES_INSIDE_ISRS // block interrupts during free-storage use
161         UBaseType_t usis; // saved interrupt status
162     #endif
163     static char *currentHeapEnd = &__HeapBase;
164     #ifndef STM_VERSION // Use STM CubeMX LD symbols for heap
165         if(TotalHeapSize==0) {
166             TotalHeapSize = heapBytesRemaining = (int)((&__HeapLimit)-(&__HeapBase))-
167         };
168     #endif
169     char* limit = (xTaskGetSchedulerState()==taskSCHEDULER_NOT_STARTED) ?
170         stack_ptr : // Before scheduler is started, limit is stack pointer
171         &__HeapLimit-ISR_STACK_LENGTH_BYTES; // Once running, OK to reuse al
172     DRN_ENTER_CRITICAL_SECTION(usis);
173     if (currentHeapEnd + incr > limit) {
174         // Ooops, no more memory available...
175         #if( configUSE_MALLOC_FAILED_HOOK == 1 )
176             {
177                 extern void vApplicationMallocFailedHook( void );
178                 DRN_EXIT_CRITICAL_SECTION(usis);
179                 vApplicationMallocFailedHook();
180             }
181             #elif defined(configHARD_STOP_ON_MALLOC_FAILURE)
182                 // If you want to alert debugger or halt...
183                 // WARNING: brkpt instruction may prevent watchdog operation...
184                 while(1) { __asm("bkpt #0"); }; // Stop in GUI as if at a breakpoint
185             #else
186                 // Default, if you prefer to believe your application will gracefully
187                 pReent->errno = ENOMEM; // newlib's thread-specific errno
188                 DRN_EXIT_CRITICAL_SECTION(usis);
189             #endif
190             return (char *)-1; // the malloc-family routine that called sbrk will ret
191         }
192         // 'incr' of memory is available: update accounting and return it.
193         char *previousHeapEnd = currentHeapEnd;
194         currentHeapEnd += incr;
195         heapBytesRemaining -= incr;
196         #ifndef NDEBUG
197             totalBytesProvidedBySBRK += incr;
198         #endif
199         DRN_EXIT_CRITICAL_SECTION(usis);
200         return (char *) previousHeapEnd;
201     }
202     ///! non-reentrant sbrk uses is actually reentrant by using current context
203     // ... because the current _reent structure is pointed to by global _impure_ptr
204     char * sbrk(int incr) { return _sbrk_r(_impure_ptr, incr); }

```

```

205  //! _sbrk is a synonym for sbrk.
206  char * _sbrk(int incr) { return sbrk(incr); };
207
208  #ifndef MALLOCES_INSIDE_ISR // block interrupts during free-storage use
209      static UBaseType_t malLock_uxSavedInterruptStatus;
210  #endif
211  void __malloc_lock(struct _reent *r)    {
212      #if defined(MALLOCES_INSIDE_ISR)
213          DRN_ENTER_CRITICAL_SECTION(malLock_uxSavedInterruptStatus);
214      #else
215          bool insideAnISR = xPortIsInsideInterrupt();
216          configASSERT( !insideAnISR ); // Make damn sure no more mallocs inside ISRs!
217          vTaskSuspendAll();
218      #endif
219  };
220  void __malloc_unlock(struct _reent *r) {
221      #if defined(MALLOCES_INSIDE_ISR)
222          DRN_EXIT_CRITICAL_SECTION(malLock_uxSavedInterruptStatus);
223      #else
224          (void)xTaskResumeAll();
225      #endif
226  };
227
228  // newlib also requires implementing locks for the application's environment mem
229  // accessed by newlib's setenv() and getenv() functions.
230  // As these are trivial functions, momentarily suspend task switching (rather tha
231  // Not required (and trimmed by linker) in applications not using environment var
232  // ToDo: Move __env_lock/unlock to a separate newlib helper file.
233  void __env_lock()      { vTaskSuspendAll(); };
234  void __env_unlock()    { (void)xTaskResumeAll(); };
235
236  #if 1 // Provide malloc debug and accounting wrappers
237      /// /brief Wrap malloc/malloc_r to help debug who requests memory and why.
238      /// To use these, add linker options: -Xlinker --wrap=malloc -Xlinker --wrap=_
239      // Note: These functions are normally unused and stripped by linker.
240      size_t TotalMallocdBytes;
241      int MallocCallCnt;
242      static bool inside_malloc;
243      void * __wrap_malloc(size_t nbytes) {
244          extern void * __real_malloc(size_t nbytes);
245          MallocCallCnt++;
246          TotalMallocdBytes += nbytes;
247          inside_malloc = true;
248          void *p = __real_malloc(nbytes); // will call malloc_r...
249          inside_malloc = false;
250          return p;
251      };
252      void * __wrap__malloc_r(void *reent, size_t nbytes) {
253          extern void * __real__malloc_r(size_t nbytes);
254          if (!inside_malloc) {
255              MallocCallCnt++;
256              TotalMallocdBytes += nbytes;
257          };
258          void *p = __real__malloc_r(nbytes);
259          return p;
260      };
261  #endif
262
263  // =====
264  // Implement FreeRTOS's memory API using newlib-provided malloc family.
265  // =====
266
267  void *pvPortMalloc( size_t xSize ) PRIVILEGED_FUNCTION {
268      void *p = malloc(xSize);
269      return p;

```



```
270 }
271 void vPortFree( void *pv ) PRIVILEGED_FUNCTION {
272     free(pv);
273 };
274
275 size_t xPortGetFreeHeapSize( void ) PRIVILEGED_FUNCTION {
276     struct mallinfo mi = mallinfo(); // available space now managed by newlib
277     return mi.fordblks + heapBytesRemaining; // plus space not yet handed to newlib
278 }
279
280 // GetMinimumEverFree is not available in newlib's malloc implementation.
281 // So, no implementation is provided: size_t xPortGetMinimumEverFreeHeapSize( void )
282
283 //! No implementation needed, but stub provided in case application already calls it
284 void vPortInitialiseBlocks( void ) PRIVILEGED_FUNCTION {};
```

Summary

Unfortunately some vendors distribute tool chains with incorrect examples of FreeRTOS/newlib. However, with a bit of care newlib works well and safely with FreeRTOS.

Enjoy!

Best Regards, Dave

PS: newlib provides facilities for wrapping stdio functions, not covered in this article. You will want to use these, for example, if your application uses posix IO functions to read and write to a USB stick using a local FAT implementation.

PS: Hope Richard Barry and the FreeRTOS team will add heap_useNewlib.c to FreeRTOS ;-)

Additional References

<http://www.billgatliff.com/newlib.html>

http://wiki.osdev.org/Porting_Newlib

<http://www.embecosm.com/appnotes/ean9/ean9-howto-newlib-1.0.html>

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