# UNIX

# Sockets

- Interprocess Communication channel:
  - descriptor based
  - two way communication
  - can connect processes on different machines
- Three most typical socket types (colloquial names):
  - unix (local connection)
  - tcp (reliable, byte stream connection over network)
  - **udp** (unreliable, datagram connection over network)

#### Communication styles

- SOCK\_STREAM
  - binary stream
  - reliable
  - context
  - communication with one peer only
  - no records boundaries
  - connection oriented
  - fifo-like connection type
  - bidirectional
  - out of the band data may be (and usually is) supported

- Communication styles
  - SOCK\_DGRAM
    - packets
    - data loss is possible
    - data duplicates are possible
    - order of delivery may be different than the order of sending
    - connectionless
    - communication with many peers
    - each packet has to be individually addressed

- Communication styles (cont.)
  - SOCK\_RAW
    - low level network access
    - e.g. icmp ping request
    - requires proper capability or root privileges
    - SOCK\_DGRAM similar type of connection

- Addresses, namespace, domain
  - usually internet AF\_INET or local (unix) AF\_LOCAL,
     AF\_UNIX
  - new socket can can be bound to an address
    - obligatory for server address
    - client socket is usually bound automatically on first send or connect operation

#### Protocol

- tcp for byte stream contextual reliable connection
- udp for packet contextless unreliable connection
- usually programs use default (zero) protocol for given communication style and name space
- Connection requires both ends to use the same protocol in the same namespace and communication style

• Structure sockaddr is used for all namespaces:

```
sa_family_t sa_family
char sa_data[14]
```

- this type is used only to cast addresses used as parameters to bind, getsockname and other functions
- sa\_family can be one of AF\_INET, AF\_LOCAL, AF\_UNIX
- each namespace defines its own structure, field
   sa\_family is common to all those structures and can be used to recognize socket type

- If other end of stream socket connection is closed and the connection buffer is empty:
  - reading from this channel will return end of file status (read returns zero as number of bytes read)
  - writing to the channel results in delivery of SIGPIPE signal. If this signal is handled, ignored or blocked write/send/sendto will return the EPIPE error
  - writing and reading may result in ECONNRESET error
- EPIPE and ECONNRESET are not critical errors and can be properly recognized and handled in students' apps without terminating the application

- Socket connection typical use case (very rarely modified)
  - server (connection listener)
    - create socket
    - bind it with the address
    - set listen queue (SOCK\_STREAM only)
    - accept connections (SOCK\_STREAM only)
    - do work
    - close socket

- client (connection initiator)
  - create socket
  - optionally bind with an address
  - connect
  - do work
  - close

#### • Out of Band data

- priority messages can be send over the socket
- usually used to indicate exceptional conditions
- data is sent independently of the ordinary data
- process must use send/recv functions with
   MSG\_OOB flag, without this flag only ordinary data will be received
- select and poll function can wait for OOB data
- special marker is set to indicate the position of band data in the ordinary data
- more information: glibc manual chapter "Out-of-Band

- Sockets can be used to communicate processes working on different architectures, this affects the way data types are treated
- Failure to recognize the problem is the most common source of errors in students' works
- Byte Order
  - does not affect 1 byte data types
    - char (not unicode wchar\_t)
    - [u]int8 t
    - the safest method to send data is to format it as text!

- Byte Order (cont.)
  - 2- and more- byte integers
    - little-endian less significant byte goes first
      - the higher memory address the more significant byte (255=FF00)
      - x86, amd64
    - big-endian more significant byte goes first
      - the lower memory address the more significant byte (255=00FF)
      - SPARC, PPC
    - network byte order (big-endian), all data should be converted to network byte order before being sent and back to host byte order after being received

#### • Byte Order (cont.)

- 2 byte data type ([u]int16\_t) can be converted to network order with htons function and back to host order with ntohs function
- 4 byte data type ([u]int32\_t) can be converted to network order with htonl function and back to host order with ntohl function
- 8 byte data types ([u]int64\_t) must be converted manually, no standard function or macro exists, it is easy to write one though

- Byte Order (cont.)
  - double and float
    - do not have network format defined!
    - the floating point format may be different!
    - it should be sent as human readable string and parsed at destination
  - arrays
    - all members of array should be converted separately
  - structures
    - all members of a structure should be converted separately
    - structure must be packed (see next slides)

#### Data types sizes

- char is a one-byte data type (always)
- do not use traditional data types int and long as they may differ in size
- use UNIX standardized integer types:
  - [u]int[8|16|32|64]\_t
- Structures have different field alignment
  - some architectures prefer 2,4 or 8 byte alignment
  - if communicating architectures use different alignment the structure will be malformed after transport

- Structures have different field alignment (cont.)
  - compiler can enforce the smallest alignment possible that is the same on all architectures so called *packing*
  - unfortunately packing options are specific to given compilers
    - some use #pragma pack;
    - gcc uses \_\_attribute\_\_((\_\_packed\_\_)) on all structures to pack or -fpack-struct option to pack all structures in the code
  - there is no portable way of packing, structures should not be sent over socket in portable programs

#### • Function socket:

- creates unbound and unconnected socket descriptor
- domain selects namespace, usually one of PF\_INET,
   PF\_LOCAL, PF\_UNIX (notice that constant has PF\_not AF\_prefix)
- type indicates communication style, usually
   SOCK STREAM or SOCK DGRAM
- protocol usually defaults to zero as only one protocol exists for each combination of domain and type

#### • Function socketpair:

- creates a pair of unnamed and connected socket
- only related processes can use it
- can be used the same way as pipe, except that it offers twoway communication
- domain can be only PF\_LOCAL, PF\_UNIX on Linux and most of other systems
- other parameters are passed similarly as to socket function
- socket\_vector array of 2 undistinguishable socket descriptors (return parameter)

#### • Function bind:

- assigns name to the socket
  - not obligatory, process can use unbound socket, and the name will be assigned automatically, process can use f.

    getsockname to learn this name
  - name of local end of the connection is assigned
  - name of remote peer can be read with f. getpeername
- address must be cast to struct sockaddr \*
- address structure real length must be calculated accordingly to namespace type (see examples) as it is not always matter of a simple sizeof.

- Functions setsockopt and getsockopt:
  - control various socket options (see examples)
  - useful POSIX compliant options for AF\_INET:
    - SO REUSEADDR(SOL SOCKET level)
      - application fails to bind to an address some time after previous run (no mater how it was terminated). It takes even a few minutes for the system to clear the address, this option will speed up reuse, it should be instant but makes tep less reliable as time-out for lost packets is removed
    - SO\_KEEPALIVE(SOL\_SOCKET level)
      - starts transparent exchange of test message, if connection is broken it can speed up detection of the problem. Waiting processes are sent SIGPIPE if test message fails to go through.

- Functions setsockopt and getsockopt(cont.):
  - SO\_BROADCAST (SOL\_SOCKET level)
    - permits sending broadcasts form the sockets
    - broadcast can be blocked on firewall, usually no broadcasts form the external network are allowed in LAN
  - IP\_MTU (IPPROTO\_IP level)
    - reads current known MTU (Maximum Transport Unit) the maximum reliable size of datagram process can send
    - can be obtained (never set) for connected socket only
    - depends on network connection
    - can change
    - cannot be less than 576
  - SO\_ERROR (SOL\_SOCKET level)
    - check for network pending errors

#### • Function connect:

- connects socket with the other end
- usually client connects to the server
- socket being connected does not need to be bound to any name
- server socket must have a name

- Function connect (cont.):
  - packet (datagram)
    - does not require permanent connection as each packet can be addressed individually
    - connecting is still possible and useful when many packets are to be sent to the same destination
    - process can connect the same packet socket many times to change destination address
    - connect will never block as the real connection is not established

- Function connect (cont.):
  - SOCK\_STREAM connection
    - connection is obligatory and possible only once per given socket
    - connect will block (provided that O\_NONBLOCK flag was not set on socket descriptor) until connection is made
    - if socket is in non-blocking mode, or connect was interrupted by signal handling routine, connecting will **continue asynchronously**!!! In such a case:
      - process should wait for socket write readiness (see examples)
      - or assume socket to be connected, sleep 1-2 sec. If assumption is wrong the first operation on socket will fail (rather untidy approach)
      - repeated connect returns EALREADY error, it must be recognised if
         TEMP FAILURE RETRY macro is used

#### • Function listen:

- applies only to byte stream communication
- enables connection requests on the socket
- marks socket as server socket that should be used only to accept connections
- backlog argument
  - merely a hint of size of queue of pending connections (i.e. connections already initiated by client and not yet accepted)
  - not the maximum number of simultaneously connected clients
  - POSIX suggests that passing zero as backlog parameter <u>may</u> mean implementation-defined minimal value
  - if pending connections queue is exhausted peer may be rejected

#### • Function accept:

- applies only to byte stream communication
- creates a new socket for communication with connecting peer, returns new descriptor
- new socket communication type and namespace are inherited from listening socket
- newly created socket <u>may</u> not inherit flags set on listening socket (on Linux flags are not inherited), portable application should set flags explicitly
- listening socket remains unconnected and can accept next connection
- connections are accepted in the order they were queued

- Function accept (cont.):
  - blocks unless non-blocking flag is set for descriptor
  - listening socket indicates readiness for a new connection by readable event (f. select), even when select returns, such a connection may be already lost due to network errors, thus accept will block. To avoid this situation, non-blocking mode should be set on the socket.

- Function accept and getpeername (cont.):
  - can obtain peer socket address
    - via output parameters
    - if address buffer is too small, the address will be **silently** truncated (the function cannot overflow the buffer)
    - if connection is made by unbound peer the value of the address is unspecified (udp and tcp use bounded clients only)
    - there is a very interesting discussion on the type of last argument to accept (see NOTES in Linux manual on the accept function)
    - On HP-UX does not support AF\_UNIX

- Functions socket, socketpair, listen, accept and bind (cont.):
  - from linux manual:

POSIX.1-2001 does not require the inclusion of <sys/types.h>, and this header file is not required on Linux. However, some historical (BSD) implementations required this header file, and portable applications are probably wise to include it.

#### Operations on socket

- blocks if operation buffer is full (and flag O\_NONBLOCK is not set)
- success of operation does not mean the delivery, only successful sending
- if datagram is received it must be read at once, otherwise the remaining part will be truncated!
- can report errors from previous operations (pending network errors)
- using read/write functions on connected socket

- Operations on socket (cont.)
  - with recv/send functions on connected socket
    - works as read/write
    - process can apply socket specific flags (POSIX list)
      - MSG\_OOB out of band data
      - MSG\_EOR end of record marker (not supported on byte stream)
      - MSG\_PEEK do not remove data from the buffer
      - MSG\_WAITALL the read shall not be terminated in the middle of a message (but still can be interrupted before any data is actually read)
      - flag MSG DONTWAIT is Linux specific!
      - <u>flag MSG\_NOSIGNAL is Linux specific!</u>

- Operations on socket (cont.)
  - with recvfrom/sendto functions
    - works as recv/send
    - receiving process can learn peer address
    - sending process can address the message (datagrams only)
  - with recvmsg/sendmsg functions
    - works as recvfrom/sendto
    - sends messages with control information

#### Closing a socket

- with close function
- with shutdown function a process can shut only a part of a connection
  - SHUT\_RD process stops reading from the socket any incoming data should be rejected, process can still send data. If process tries to read from such a socket it will get EOF marker. Peer may get EPIPE or data may be silently discarded, in practice data is still received and queued, can be read in some systems!
  - SHUT\_WR process stops writing to socket any outgoing data will be rejected, but process still can read from the socket. If process tries to write to such a socket it will be send SIGPIPE (or EPIPE error)
  - SHUT\_RDWR closes both ways, still requires f.close.

#### Unix Domain Sockets

- Communication between processes on the same computer
  - socket name is simply a name in local file system
  - processes must have write (w) and search (x) permissions to the directory where socket is created
  - process must have read and write permissions to the directory containing the socket it is connecting to
  - permissions on local sockets can be ignored by some systems and should not be used as security measures
  - usually local sockets are put in /tmp

#### Unix Domain Sockets

- both communication styles (*SOCK\_STREAM*, *SOCK\_DGRAM*) are supported in this namespace
- name in file system is necessary only to establish connection
- socket file persists and can/should be deleted by the process
- local sockets can not be used to connect remote process (even through nfs or afs) at the moment but it is possible that it will change in the future.

### Unix Domain Sockets

- for better portability process should not depend on the fact that the other process resides on the same machine
  - pid numbers
  - byte order
  - sending structures
- for sake of simplicity students can assume that local sockets are truly local in laboratory applications
- Linux implementation does not support out of bound data transfer on local sockets

### Unix Domain Sockets

• Address in unix domain sockets are stored in struct sockaddr un:

```
sa_family_t sun_family
char sun path[108]
```

• The length of this structure (passed to bind) should be calculated as:

```
sizeof(sa family t)+strlen(a.sun path)
```

- The SUN\_LEN macro returns proper size
- Do not use sizeof (sockaddr\_un)!!!

- Communication between processes over the tcp/ip network
  - IPv4 vs. IPv6, at the moment only IPv4 is common and will be discussed here
  - socket name consists of IP address (x.x.x.x) and the port number both represented in network byte order (see byte order)
  - both communication styles (SOCK\_STREAM,
     SOCK\_DGRAM) are supported in this namespace
  - no permission control built into AF\_INET sockets !!!

• Address in internet domain sockets are stored in struct sockaddr in:

```
sa_family_t sin_family
struct in_addr sin_addr
unsigned short int sin port
```

• Where in addr is a one-filed structure:

```
u_int32_t s_addr;
```

• The length of this structure (passed to bind) should be calculated as

sizeof(sockaddr in)

- Special addresses (<u>host byte order</u>):
  - uint32 t INADDR LOOPBACK (127.0.0.1)
  - -uint32 t INADDR ANY (0.0.0.0)
    - very useful as server address
    - binds to all local interfaces
    - if socket is unbound prior to connect or send it will be automatically bound to INADDR ANY and random free port
  - uint32 t INADDR BROADCAST (255.255.255.255)
    - to send broadcast datagrams (SOCK\_DGRAM)
  - uint32 t INADDR NONE (-1=255.255.255.255)
    - error indicator

- IP address by type:
  - unicast addressing of single peer
    - only option for SOCK STREAM
  - multicast addressing a set of peers in the network
    - not available by default, popular, speeds up streaming
    - will not be discussed here
  - broadcast addressing all the peers in the network
    - useful communication method in local networks
    - can be blocked by routers and local firewalls and we experienced many problems with broadcast tasks during labs

- Socket address (cont.)
  - to convert IP string e.g. "194.29.178.1" to uint32\_t
     POSIX defines inet addr function
    - returns INADDR\_NONE as error, it may be mistaken with INADDR\_BROADCAST
    - inet addr returns an address in <a href="mailto:network byte order">network byte order</a>
    - GNU extension inet\_aton reports errors more reliably but is not standardized
  - to convert binary address in <u>network byte order</u> to string POSIX defines inet ntoa
    - returns string in static buffer not thread safe

- Socket address (cont.)
  - as socket requires IP number to bind, domain name must be resolved (DNS) with gethostbyname
    - on success returns hostent structure, where address in question is available as field char \*h\_addr\_list[0] (must be cast to struct in addr)
    - errors are reported in global variable h\_errno not errno!!!
    - can be used with domain names and IP numbers as strings on Linux (POSIX does not define it)
    - returns addresses in network byte order
    - returned address may be stored in static buffer
    - not thread safe

#### Datagrams (packets)

- if size of datagram is larger than MTU on given socket sending will return with EMSGSIZE error
- process can test MTU with getsockopt or experiment with the sizes and EMSGSIZE or assume small datagrams (less than 576)
- read/write/send/recv/sendto and recvfrom will perform atomically on datagrams
- if datagram is received it must be read at once, otherwise the remaining part will be truncated!

- Datagrams (packets) cont.
  - if datagram is lost, application must retransmit it, usually if response is not received after some time-out packet can be considered lost
  - usually messages are sent in single packets
  - application logic must be ready for lost packets and duplications
  - on Linux datagrams are reliable, duplications and mixed order or delivery are not possible, but this is not POSIX behaviour
  - Linux also has a reliable datagram SOCK RDM socket type

- Binary stream
  - reliable
  - usually only one process operates on each end of connection
  - message boundaries are not preserved
  - there is no atomic message size !!!
  - read/write operation can be interrupted at any stage (EINTR)

- Binary stream (cont.)
  - usually connection is bidirectional (read/write), process can modify file flags or use shutdown function to limit access
  - file position is fixed, reading from beginning, writing at the end (FIFO order)
  - read buffer is independent from write buffer

• How to make local SOCK\_STREAM socket:

 How to bind and start listening on local SOCK\_STREAM socket:

```
#define BACKLOG 3
int bind socket(char *name) {
    struct sockaddr un addr;
    int socketfd;
    if (unlink(name) < 0 && errno != ENOENT)</pre>
          ERR("unlink");
    socketfd = make socket(name, &addr);
    if (bind(socketfd, (struct sockaddr*) &addr, \
          SUN LEN(&addr)) < 0) ERR("bind");
    if (listen(socketfd, BACKLOG) < 0) ERR("listen");</pre>
    return socketfd;
```

How to connect to local SOCK\_STREAM socket:

```
int connect socket(char *name) {
  struct sockaddr un addr; int socketfd;
  socketfd = make socket(name, &addr);
  if (connect(socketfd, (struct sockaddr*) &addr, \
                          SUN LEN(&addr)) < 0) {
      if (errno != EINTR) ERR("connect");
      else {
         fd set wfds ; int status;
         socklen t size = sizeof(int);
         FD ZERO(&wfds); FD SET(socketfd, &wfds);
         if (TEMP FAILURE RETRY(select(socketfd+1, \
             NULL, &wfds, NULL, NULL)) < 0) ERR("select");
         if (getsockopt(socketfd, SOL SOCKET, SO ERROR, \
               &status, &size) < 0) ERR ("getsockopt");
         if(0!=status) ERR("connect");
```

return socketfd; }

How to accept connection on SOCK\_STREAM socket:

How to make internet server SOCK\_DGRAM socket:

```
int make socket(uint16 t port) {
    struct sockaddr in name;
    int sock, t=1;
    sock = socket(PF INET, SOCK DGRAM, 0);
    if (sock < 0) ERR("socket");</pre>
    name.sin family = AF INET;
    name.sin port = htons(port);
    name.sin addr.s addr = htonl(INADDR ANY);
    if (setsockopt(sock, SOL SOCKET, SO REUSEADDR,
        &t, sizeof(t))) ERR("setsockopt");
    if (bind(sock, (struct sockaddr*) &name, \
        sizeof(name)) < 0) ERR("bind");</pre>
    return sock;
```

How to receive datagrams:

How to send datagrams:

```
int send_datagram(int sock, struct sockaddr_in addr,
int16_t msg) {
   int status;
   int16_t buf = htons(msg);
   status = TEMP_FAILURE_RETRY(sendto(sock, &buf, \
        sizeof(int16_t), 0, (struct sockaddr*) \
        &addr, sizeof(addr)));
   if (status < 0 && errno != EPIPE && errno ==

ECONNRESET)
        ERR("sendto");
   return status;
}</pre>
```

How to lookup domain name in DNS:

```
struct sockaddr in make address (char *address,
uint16 t port) {
                               struct sockaddr in addr;
                               struct hostent *hostinfo;
                               addr.sin family = AF INET;
                               addr.sin port = <a href="https://example.com/https://example.com/https://example.com/https://example.com/https://example.com/https://example.com/https://example.com/https://example.com/https://example.com/https://example.com/https://example.com/https://example.com/https://example.com/https://example.com/https://example.com/https://example.com/https://example.com/https://example.com/https://example.com/https://example.com/https://example.com/https://example.com/https://example.com/https://example.com/https://example.com/https://example.com/https://example.com/https://example.com/https://example.com/https://example.com/https://example.com/https://example.com/https://example.com/https://example.com/https://example.com/https://example.com/https://example.com/https://example.com/https://example.com/https://example.com/https://example.com/https://example.com/https://example.com/https://example.com/https://example.com/https://example.com/https://example.com/https://example.com/https://example.com/https://example.com/https://example.com/https://example.com/https://example.com/https://example.com/https://example.com/https://example.com/https://example.com/https://example.com/https://example.com/https://example.com/https://example.com/https://example.com/https://example.com/https://example.com/https://example.com/https://example.com/https://example.com/https://example.com/https://example.com/https://example.com/https://example.com/https://example.com/https://example.com/https://example.com/https://example.com/https://example.com/https://example.com/https://example.com/https://example.com/https://example.com/https://example.com/https://example.com/https://example.com/https://example.com/https://example.com/https://example.com/https://example.com/https://example.com/https://example.com/https://example.com/https://example.com/https://example.com/https://example.com/https://example.com/https://example.com/https://example.com/https://example.com/https://example.com/https://example.com/https://example.com/https://example.com/h
                               hostinfo = gethostbyname(address);
                               if (hostinfo == NULL) ERRH("qethost:");
                              /*h errno*/
                               addr.sin addr =
                                                             *(struct in addr*)hostinfo->h addr;
                               return addr;
```