

## 13.1 Generic chaining and Dudley's entropy integral

Let  $\{x(t), t \in T\}$  be a stochastic process, i.e. a family of random variables indexed by  $T$ .

**Example 1.** Brownian motion.

**Example 2.**  $\mathcal{F} = \{f : S \rightarrow \{\pm 1\}\}$  a collection of binary classifiers.  $(X_1, Y_1), \dots, (X_n, Y_n) \in S \times \{\pm 1\}$  iid from  $P$ .

Define a function of  $f$  by

$$Z_n(f) = \frac{1}{n} \sum_{j=1}^n I\{Y_j \neq f(X_j)\} - Pr(Y \neq f(X)).$$

For fixed  $f$  the terms are small by LLN, and Hoeffding's inequality bounds the total difference.

**Problem 1.** What is the size of  $\sup_{t \in T} |x(t) - \mathbb{E}x(t)|$ ?

To ensure that this quantity is measurable, we will *define*

$$\sup_{t \in T} x(t) := \sup_{S \subset T \text{ finite}} \max_{t \in S} x(t).$$

Assume now that  $(T, d)$  is a metric space and that  $T$  is finite (we lose no generality here by our definition of sup above). For  $S \subset T$ ,  $t \in T$ , define

$$d(t, S) = \inf_{s \in S} d(t, s).$$

We also now make the assumption that the process  $\{x(t) : t \in T\}$  has sub-Gaussian increments with respect to  $d$ , meaning for all  $t_1, t_2 \in T$ ,

$$x(t_1) - x(t_2) = SG(d(t_1, t_2)^2).$$

**Example 3.** Brownian motion:  $\{w(t), t \in (0, 1)\}$ ,  $w(t) - w(s) \sim \mathcal{N}(0, |t - s|)$ ,  $d(s, t) = \sqrt{|t - s|}$ .

Let  $T_0 \subset T_1 \subset T_2 \subset \dots \subset T_n \subset \dots = T$ , and assume that the cardinality of  $T_k$  is  $2^{2^{k-1}}$  (note that this chain terminates because  $T$  is finite).

**Theorem 1.** For all  $u \geq 0$  and any approximating sequence  $T_n$ ,

$$Pr \left( \sup_{t \in T} |x(t) - x(t_0)| \geq 2u \sup_{t \in T} \sum_{n=0}^{\infty} 2^{n/2} d(t, T_n) \right) \leq ce^{-u^2/4}$$

for some constant  $c$ .

*Proof:*

$$\gamma_2(T) := \inf_{T_n} \sup_t \sum_{n \geq 0} 2^{n/2} d(t, T_n)$$

is called the *generic chaining complexity* and depends only on the metric space  $(T, d)$ .

First let  $u > 0$  be given. Define  $\pi_n : T \rightarrow T_n$  to be the nearest-neighbor function which sends an element of  $T$  to its nearest neighbor in  $T_n$ , so that  $d(t, \pi_n t) = d(t, T_n)$ . Then

$$x(t) - x(t_0) = \sum_{j=0}^{\infty} (x(\pi_{j+1} t) - x(\pi_j t)).$$

Define the event

$$E = \{|x(s_1) - x(s_2)| \leq u 2^{j/2} d(s_1, s_2) \forall s_1 \in T_j, s_2 \in T_{j+1}, j \geq 0\}.$$

We want to control  $Pr(E^c)$ ; we have

$$\begin{aligned} Pr(E^c) &= Pr \left( \bigcup_{\substack{j \geq 0 \\ s_1 \in T_j \\ s_2 \in T_{j+1}}} \{|x(s_1) - x(s_2)| \geq u 2^{j/2} d(s_1, s_2)\} \right) \\ &\leq \sum_{j=0}^{\infty} \sum_{(s_1, s_2) \in T_j \times T_{j+1}} Pr(|x(s_1) - x(s_2)| > u 2^{j/2} d(s_1, s_2)) \\ &\leq 2 \exp \left( -\frac{u^2 2^j d(s_1, s_2)^2}{2 d(s_1, s_2)^2} \right) = 2e^{-u^2 2^{j-1}} \\ &\leq \sum_{j \geq 0} \text{card}(T_j) \text{card}(T_{j+1}) 2e^{-u^2 2^{j-1}} \\ &= 2 \sum_{j \geq 0} 2^{2j+2j+1} e^{-u^2 2^{j-1}} \\ &= 2 \sum_{j \geq 0} e^{-u^2 2^{j-1} + 2^{j-1}(2+4) \log 2} \\ &\leq 2 \sum_{j \geq 0} \exp(-u^2 2^{j-1} + 6 \cdot 2^{j-1} \log 2). \end{aligned}$$

If  $u^2 > 24 \log 2$ , then the sum will decay quickly, and this sum is  $\leq 2e^{-u^2 2^{j-3}} \leq ce^{-u^2/8}$  for some  $c$ . On  $E$ ,

$$|x(t) - x(t_0)| \leq \sum_{j \geq 0} |x(\pi_{j+1} t) - x(\pi_j t)| \leq u \sum_{j \geq 0} 2^{j/2} d(\pi_j t, \pi_{j+1} t) \leq u \sum_{j \geq 0} 2^{j/2} d(t, T_j) + \frac{1}{\sqrt{2}} \sum_{j \geq 0} 2^{(j+1)/2} d(t, T_{j+1})$$

which proves the result.