

# DP Coin Change DP 01 Knapsack

Week-06, Lecture-02

**Course Code:** CSE221

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# Knapsack Problem

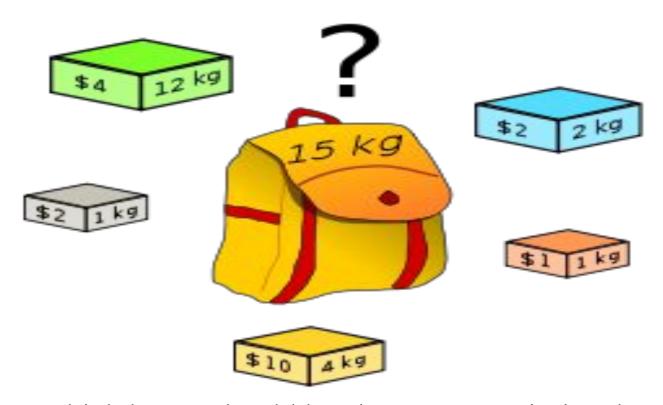
# Knapsack problem

A 1998 study of the Stony Brook University Algorithm Repository showed that, out of 75 algorithmic problems, the knapsack problem was the 18th most popular and the 4th most needed after kd-trees, suffix trees, and the bin packing problem

#### Knapsack Problem

- UGiven a set of items, each with a mass and a value, determine the number of each item to include in a collection so that the total weight is less than or equal to a given limit and the total value is as large as possible.
- It derives its name from the problem faced by someone who is constrained by a fixed-size knapsack and must fill it with the most valuable items.

#### Knapsack Problem



which boxes should be chosen to maximize the amount of money while still keeping the overall weight under or equal to 15 kg?

Answer: 3 yellow boxes and 3 grey boxes

#### **Dynamic Programming**

Recursive formula for sub problems:

$$V[k,s] = \begin{cases} V[k-1,s] & if s_k > \\ \max\{V[k-1,s], V[k-1,s-s_k] + v_k\} & els \end{cases}$$

#### **Dynamic Programming**

for 
$$s = 0$$
 to  $S$   
 $V[0, s] = 0$   
for  $i = 1$  to  $n$   
 $V[i, 0] = 0$ 

#### **Dynamic Programming**

```
for i = 1 to n
if S_i \leq S
      if v_i + V[i-1, s-s_i] > V[i-1, s]
            V[i, s] = v_i + V[i - 1, s - s_i]
      else
             V[i, s] = V[i - 1, s]
else V[i, s] = V[i - 1, s]
```

```
Given:

n = 4 (# of elements)

S = 5 (maximum size)

Elements (size, value) =

{ (2, 3), (3, 4), (4, 5), (5, 6) }
```

i\s 0 0

for 
$$s = 0$$
 to  $S$   
V[0,  $s$ ] = 0

```
i\s 0 1 2 3 4 5
0 0 0 0 0 0 0 0
1
2
3
```

for 
$$i = 0$$
 to  $n$   

$$V[i, 0] = 0$$

```
i\s 0 1 2 3 4 5
0 0 0 0 0 0 0
1 0
2 0
3 0
4 0
```

```
\begin{aligned} &\text{if } s_i \leq s \\ &\text{if } v_i + V[i-1,s-s_i] > V[i-1,s] \\ & & \textbf{V[i,s]} = \boldsymbol{v_i} + \textbf{V[i-1,s]} \end{aligned} else & V[i,s] = V[i-1,s] else V[i,s] = V[i-1,s]
```

```
\begin{aligned} \text{if } s_i &\leq s \\ \text{if } v_i + V[i-1,s-s_i] > V[i-1,s] \\ & \textbf{V[i,s]} = \textbf{v_i} + \textbf{V[i-1,s]} \end{aligned} else V[i,s] = V[i-1,s]
```

else V[i, s] = V[i - 1, s]

Items  $(s_i, v_i)$ 1: (2, 3) if  $v_i + V[i-1, s-s_i] > V[i-1, s]$ 2: (3, 4)  $V[i,s] = v_i + V[i-1,s-s_i]$ 3: (4, 5)

4: (5, 6)

else V[i,s] = V[i-1,s]else V[i, s] = V[i - 1, s]

i\s 5 3 3

if  $s_i \leq s$ 

 $\begin{array}{l} \text{if } s_i \leq s \\ \text{if } v_i + V[i-1,s-s_i] > V[i-1,s] \\ \textbf{V[i,s]} = \textbf{\textit{v}}_i + \textbf{\textit{V}}[i-1,s-s_i] \end{array}$  else  $\begin{array}{l} \text{v[i,s]} = V[i-1,s] \end{array}$ 

Items  $(s_i, v_i)$ 

i∖s	0	1	2	3	4	5
0	0	0	0	0 _	0	0
1	0	0	3	3	3	<b>→</b> 3

**2** 0

else V[i, s] = V[i - 1, s]

3

4 (

```
if s_i \leq s
 if v_i + V[i - 1, s - s_i] > V[i - 1, s]
         V[i, s] = v_i + V[i - 1, s - s_i]
 else
         V[i,s] = V[i-1,s]
else V[i, s] = V[i - 1, s]
   i\s
                                        3
                                                     3
```

- 1: (2, 3)
- 2: (3, 4)
- 3: (4, 5)
- 4: (5, 6)
  - 5
- 0 (
- 3

if 
$$s_i \le s$$
  
if  $v_i + V[i - 1, s - s_i] > V[i - 1, s]$   
 $V[i, s] = v_i + V[i - 1, s - s_i]$   
else  
 $V[i, s] = V[i - 1, s]$   
else  $V[i, s] = V[i - 1, s]$   
i\s 0 1 2 3  
0 0 0 0 0  
1 0 0 3 3

## Items $(s_i, v_i)$ 1: (2, 3) 2: (3, 4) 3: (4, 5) 4: (5, 6) 5

$$\begin{array}{l} \text{if } s_i \leq s \\ \text{if } v_i + V[i-1,s-s_i] > V[i-1,s] \\ \textbf{V[i,s]} = \textbf{v_i} + \textbf{V[i-1,s-s_i]} \\ \end{array}$$
 else

V[i,s] = V[i-1,s]

else V[i, s] = V[i - 1, s]

i\s	0	1	2	3
0	0	0	0	0
1	0	0	3	3
2	0	0	3	<b>4</b>
3	0			

Items  $(s_i, v_i)$ 

1: (2, 3)

2: (3, 4)

3: (4, 5)

4: (5, 6)

5

(

Items  $(s_i, v_i)$ 

if 
$$s_i \le s$$
  
if  $v_i + V[i-1, s-s_i] > V[i-1, s]$   

$$\mathbf{V[i, s]} = v_i + V[i-1, s-s_i]$$

else

$$V[i, s] = V[i - 1, s]$$
  
else  $V[i, s] = V[i - 1, s]$ 

i\s 0

)

 $\mathbf{C}$ 

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Items  $(s_i, v_i)$ 

if 
$$s_i \le s$$
  
if  $v_i + V[i-1, s-s_i] > V[i-1, s]$   

$$\mathbf{V[i, s]} = v_i + V[i-1, s-s_i]$$

2: (3, 4)

1: (2, 3)

else

$$V[i,s] = V[i-1,s]$$

4: (5, 6)

else V[i, s] = V[i - 1, s]

i\s

5

3

$$\begin{aligned} &\text{if } s_i \leq s \\ &\text{if } v_i + V[i-1,s-s_i] > V[i-1,s] \\ &V[i,s] = v_i + V[i-1,s-s_i] \\ &\text{else} \end{aligned}$$

$$V[i,s] = V[i-1,s]$$
else  $V[i,s] = V[i-1,s]$ 

#### Items $(s_i, v_i)$

4
 0
 3
 4
 7

if 
$$s_i \le s$$
  
if  $v_i + V[i-1, s-s_i] > V[i-1, s]$   

$$\mathbf{V[i, s]} = v_i + V[i-1, s-s_i]$$

else

$$V[i,s] = V[i-1,s]$$

else V[i, s] = V[i - 1, s]

i\s

3

Items  $(s_i, v_i)$ 

1: (2, 3)

2: (3, 4)

4: (5, 6)

$$\begin{split} \text{if } s_i &\leq s \\ \text{if } v_i + V[i-1,s-s_i] > V[i-1,s] \\ V[i,s] &= v_i + V[i-1,s-s_i] \end{split}$$

else

$$V[i,s] = V[i-1,s]$$

else V[i, s] = V[i - 1, s]

i\s

3

Items  $(s_i, v_i)$ 

1: (2, 3)

2: (3, 4)

4: (5, 6)

5

Items  $(s_i, v_i)$ 

$$\begin{aligned} &\text{if } s_i \leq s \\ &\text{if } v_i + V[i-1,s-s_i] > V[i-1,s] \\ &V[i,s] = v_i + V[i-1,s-s_i] \\ &\text{else} \end{aligned}$$

$$V[i,s] = V[i-1,s]$$

#### else V[i, s] = V[i - 1, s]

i\s

Items  $(s_i, v_i)$ 

5

$$\begin{aligned} &\text{if } s_i \leq s \\ &\text{if } v_i + V[i-1,s-s_i] > V[i-1,s] \\ &V[i,s] = v_i + V[i-1,s-s_i] \end{aligned}$$
 else

$$V[i,s] = V[i-1,s]$$

#### else V[i, s] = V[i - 1, s]

i\s	0	1	2	3	4
0	0	0	0	0	0
1	0	0	3	3	3
2	0	0	3	4	4
3	0	0	3	4	5

#### Items $(s_i, v_i)$

$$i_{\bullet}=n, k=S$$
 while  $i, k>0$  if  $V[i,k] \neq V[1-i,k]$  
$$i=i-1, k=k-s_i$$
 else

3: (4, 5)

$$i = i - 1$$

i\s	0	1	2	3	4	5
0	0	0	0	0	0	0
1	0	0	3	3	3	3
2	0	0	3	4	4	7
3	0	0	3	4	5	7
4	0	0	3	4	5	7

Items  $(s_i, v_i)$ 

$$i_{\bullet}=n, k=S$$
 while  $i, k>0$  if  $V[i,k] \neq V[1-i,k]$  
$$i=i-1, k=k-s_i$$

else

$$i = i - 1$$

i\s	0	1	2	3	4	5
0	0	0	0	0	0	0
1	0	0	3	3	3	3
				4		_
3	0	0	3	4	5	7

Items  $(s_i, v_i)$ 

$$i=n, k=S$$
 while  $i,k>0$  if  $V[i,k]\neq V[1-i,k]$  
$$i=i-1, k=k-s_i$$
 else

#### i = i - 1

i\s	0	1	2	3	4	5
0	0	0	0	0	0	0
				3		
2	0	0	3	4	4	$\begin{pmatrix} 7 \\ 7 \end{pmatrix}$
3	0	0	3	4	5	7

Items  $(s_i, v_i)$ 

$$i_{\bullet}=n, k=S$$
 while  $i,k>0$  if  $V[i,k]\neq V[1-i,k]$  
$$i=i-1, k=k-s_i$$
 else

i = i - 1

1:	(2, 3)	
2:	(3, 4)	
	(4, 5)	
4:	(5, 6)	

i\s	0	1	2	3	4	5
0	0	0	0	0	0	0
1	0	0	3 —	3	3	3
2	0	0	3	4	4	7
3	0	0	3	4	5	7
4	0	0	3	4	5	7

Items  $(s_i, v_i)$ 

$$i_{\bullet}=n, k=S$$
 while  $i,k>0$  if  $V[i,k]\neq V[1-i,k]$  
$$i=i-1, k=k-s_i$$

i = i - 1

else

i\s	0	1	2	3	4	5
0	0	0	0	0	0	0
1	0	0	0 3 3	3	3	3
2	0	0	3	4	4	7
			3		5	

#### Items $(s_i, v_i)$

$$i_{\bullet}=n, k=S$$
  
while  $i, k>0$   
if  $V[i,k] \neq V[1-i,k]$   
 $i=i-1, k=k-s_i$ 

3: (4, 5)

4: (5, 6)

$$i = i - 1$$

else

i\s	0	1	2	3	4	5
0	0	0	0	0	0	0
1	0	0	3	3	3	3
2	0	0	3	4	4	7
3	0	0	3	4	5	7
4	0	0	3	4	5	7

The optimal knapsack should contain  $\{1,2\} = 7$ 

Items  $(s_i, v_i)$ 

1: (2, 3)

2: (3, 4)

3: (4, 5)

4: (5, 6)

i∖s	0	1	2	3	4	5
0	0	0	0	0	0	0
1	0	0	$\begin{pmatrix} 0 \\ 3 \end{pmatrix}$	0	3	3
2	0	0	3	4	4	7
3	0	0	3	4	5	7
4	0	0	3	4	5	7

#### **Exercise**

Solve the Unbounded Knapsack Problem using dynamic programming for the following instance:

- Weights: [1, 3, 4, 5]
- Values: [1, 4, 5, 7]
- Capacity: 7

#### Textbooks & Web References

- Reference book iii (Chapter 19)
- www.geeksforgeeks.org
- www.codeforces.com
- www.topcoder.com

# Thank you & Any question?