

# Introduction to Divide and Conquer Approach Searching Binary Search Sorting, Merge Sort

Week-01, Lecture-02

**Course Code:** CSE221

**Course Title:** Algorithms

Program: B.Sc. in CSE

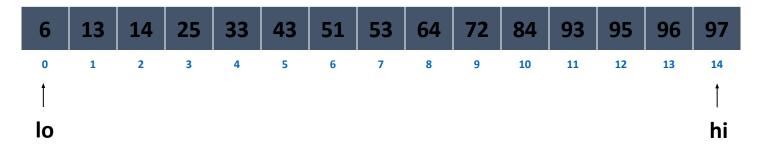
Course Teacher: Tanzina Afroz Rimi

**Designation:** Lecturer

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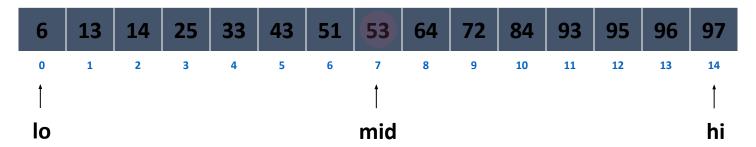
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• Invariant. Algorithm maintains a[lo] ≤ value ≤ a[hi].



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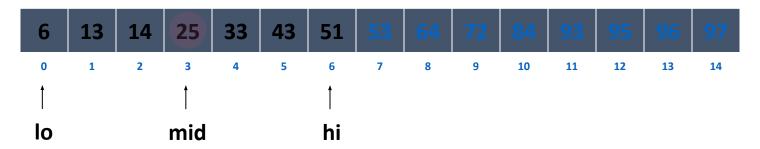
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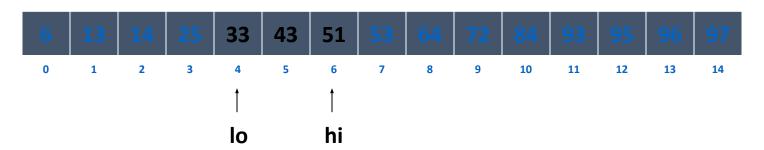
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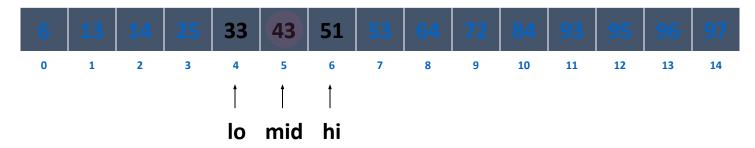
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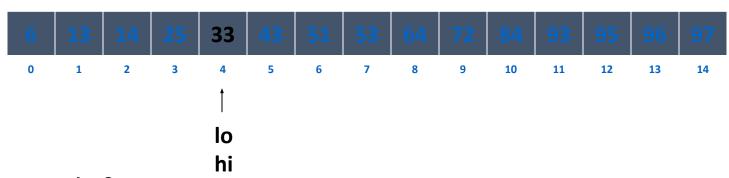
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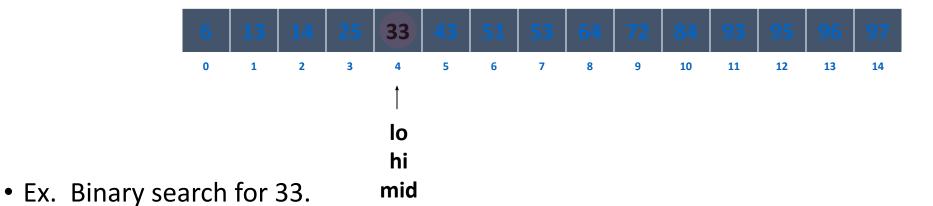


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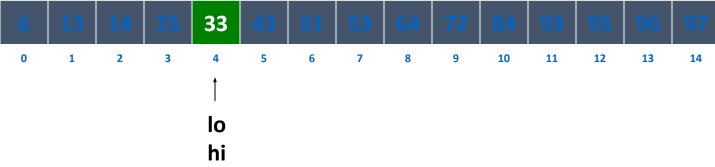
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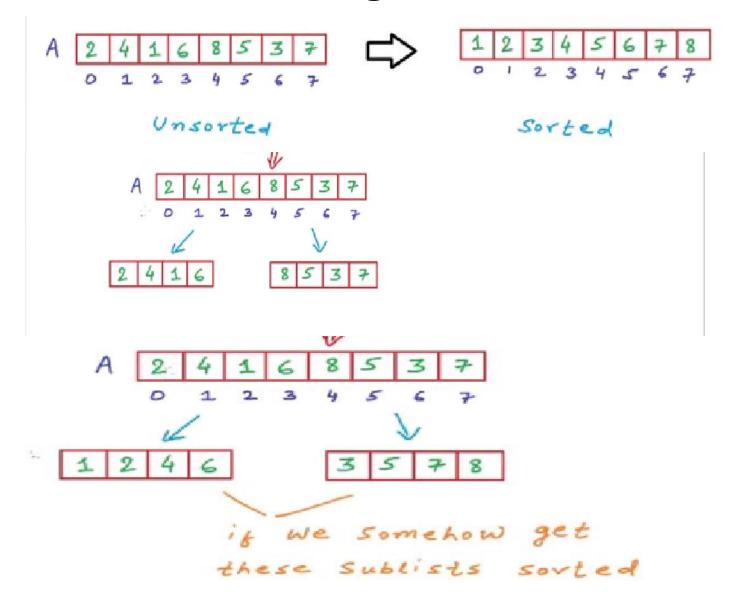


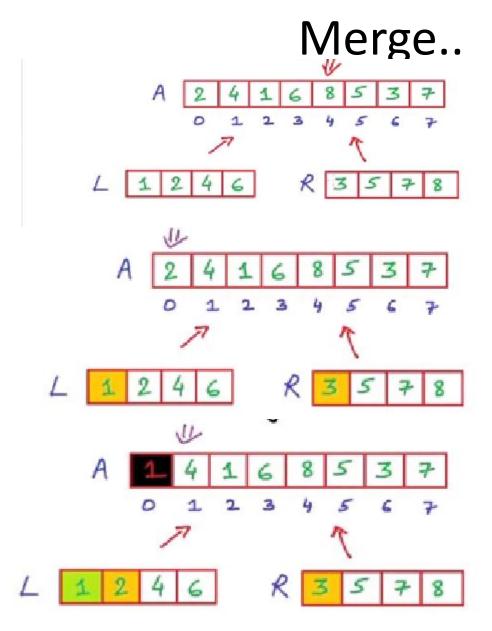
## Merge Sort

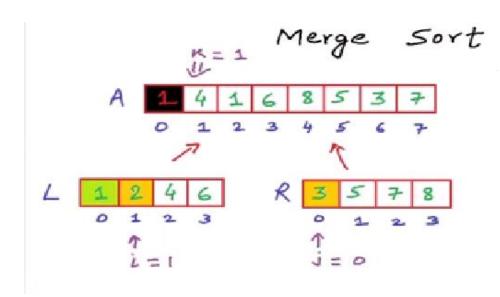
#### Divide & Conqure

- The divide-and-conquer paradigm involves three steps at each level of the recursion:
  - **Divide** the problem into a number of subproblems.
  - **Conquer** the subproblems by solving them recursively. If the subproblem sizes are small enough, however, just solve the subproblems in a straightforward manner.
  - Combine the solutions to the subproblems into the solution for the original problem.
- The merge sort algorithm closely follows the divide-and-conquer paradigm.
   Intuitively, it operates as follows.
  - **Divide:** Divide the n-element sequence to be sorted into two subsequences of n/2 elements each.
  - Conquer: Sort the two subsequences recursively using merge sort.
  - Combine: Merge the two sorted subsequences to produce the sorted answer.

#### Merge Sort







Merge (L, R, A) in Le Length (L) nRe length (R) LEJEKEO while(i< nL 44 J< nR) 1 16 ( L[i] < = R[i]) A[x] + L[i]  $i \leftarrow i + 1$ else A[K] - R[J] J = J + 1 Ke K+1

A 12368537

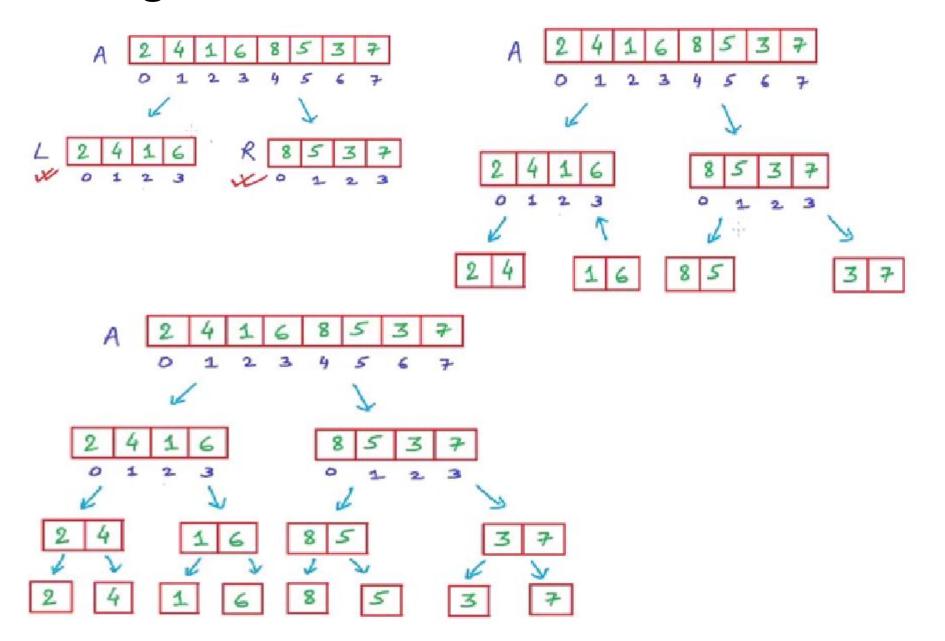
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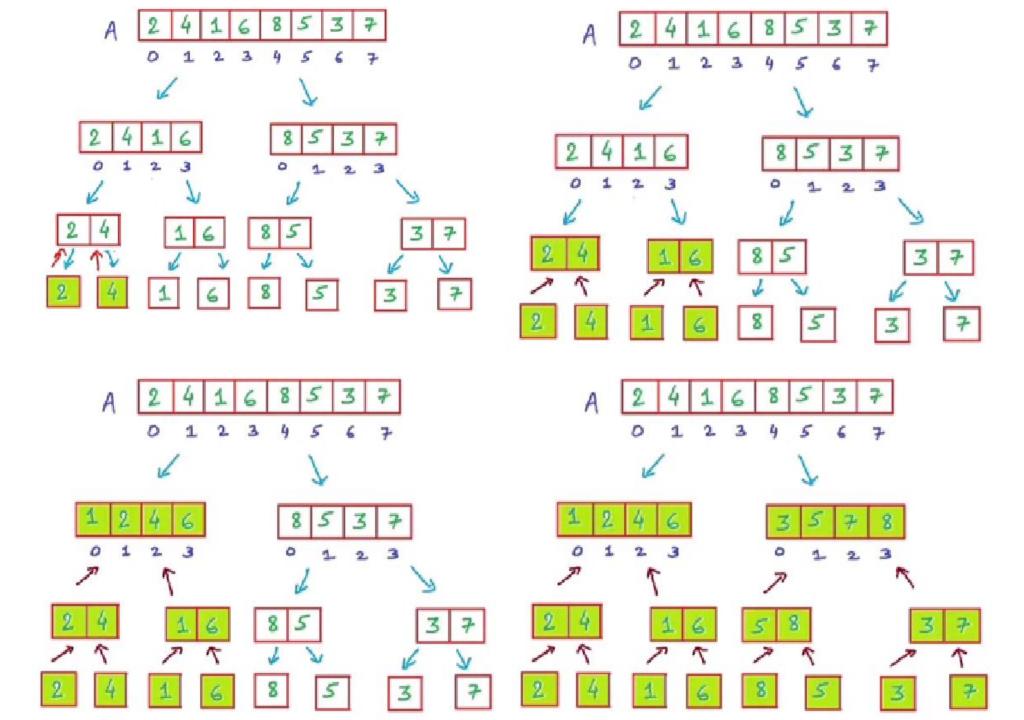
Merge (L, R, A)

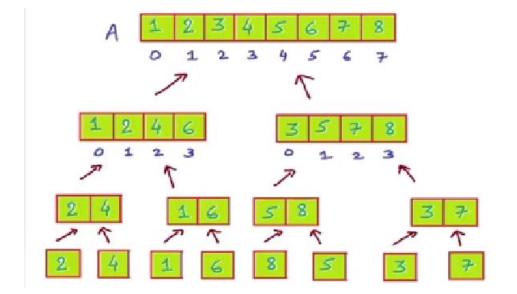
$$i = 1$$
 $i = 2$ 
 $i = 3$ 
 $i = 3$ 

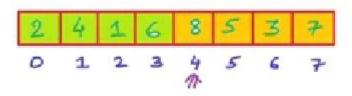
```
K=5
                    Merge (L, R, A)
                   Enle Length(L)
                      nR - length (R)
                      Lejekeo
                 while(i< ml 44 J< nR)
                   it ( L[i] < = R[i])
                      A[n] + L[i]; i + i+1
                    else
i=4
                     A[K] < R[j]; j < j+1
          K=8
                  while (i < nL)
                   A[K] < L[i]; Leiti; Kektij
                  while (VK nR)
                 1. ALKJ + R[i], JEJ+1; Ke K+1;
i = 4
```

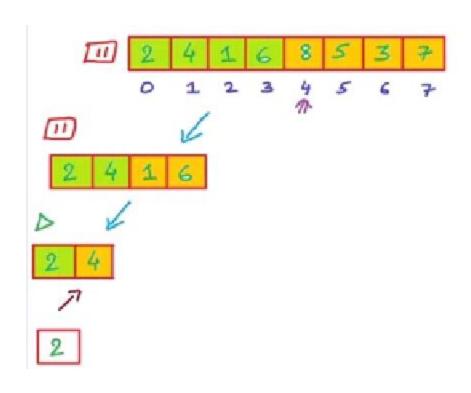
### Merge Sort..





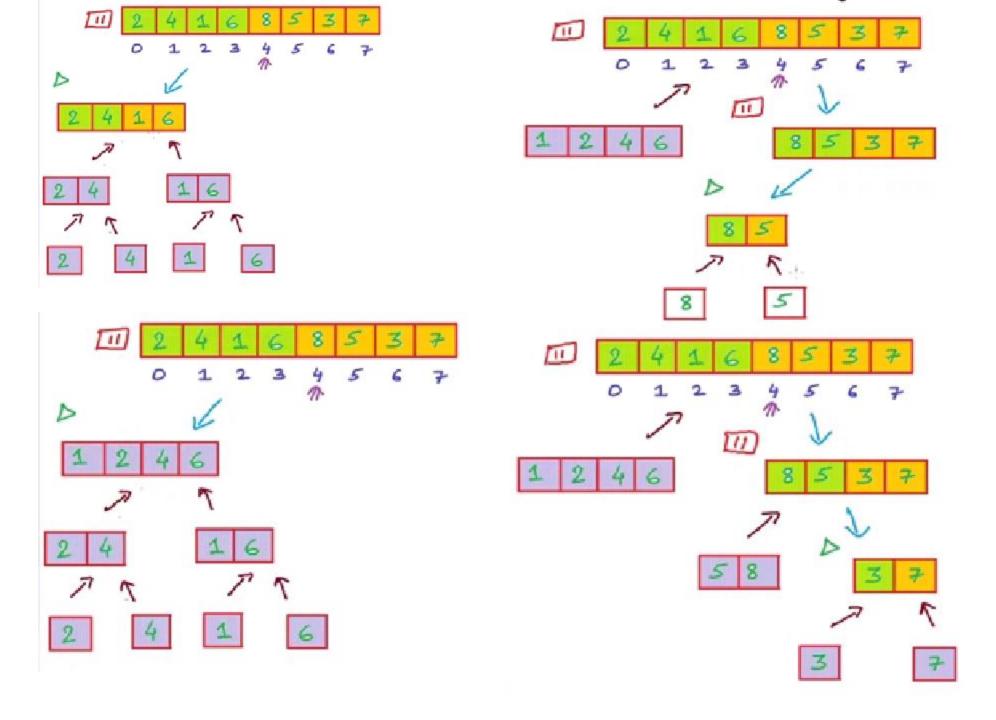


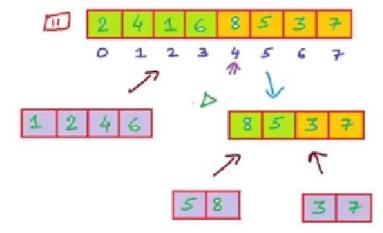


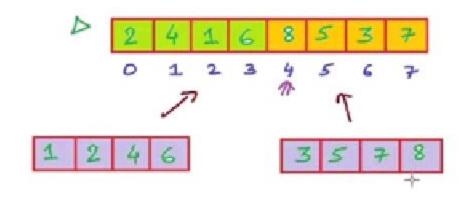


tand Length (A) if (n < 2) return Condition left array of size (mid) for it o to mid-1 reft[i] - A[i] for i < mid to n-1 right[i-mid] < A[i] Mergesort (left) Mergesort (right) , Merge (left, right, A)

Base







Mergesort (A) 1 > n - length (A) if (n < 2) return mid = n/2 left array of size (mid) right array of size (n-mid) for it o to mid-1 reft[i] = A[i] for i ← mid to n-1 right[i-mid] = A[i] Mergesort (left) Mergesort (right) , Merge (left, right, A)

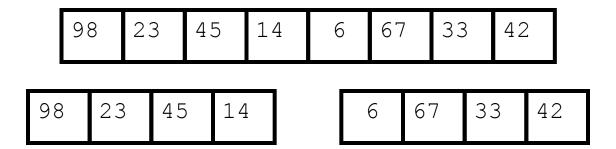
1 2 3 4 5 6 7 8

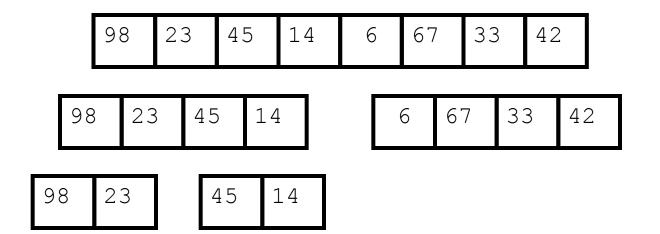
Mergesort (A) { ne length (A) base <- if (n < 2) return Condition mid = n/2 left array of size (mid) right array of size (n-mid) for it o to mid-1 reft[i] - A[i] for i < mid to n-1 Recursive right[i-mid] = A[i] Merge Sorted E Merge Sort (left)

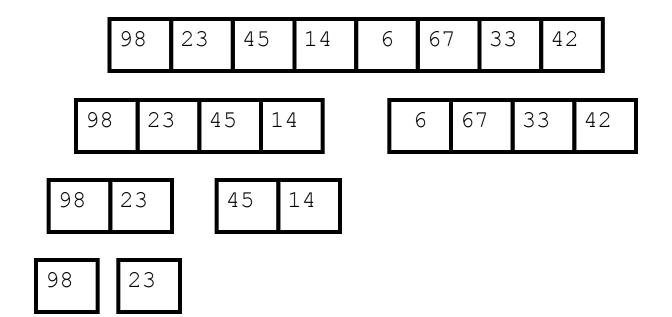
Merge Sorted E Merge Sort (right)

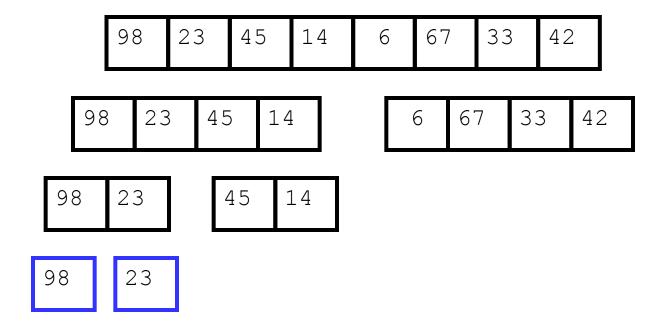
halves Merge (left, right, A)

| 98 | 23 | 45 | 14 | 6 | 67 | 33 | 42 |
|----|----|----|----|---|----|----|----|
|----|----|----|----|---|----|----|----|

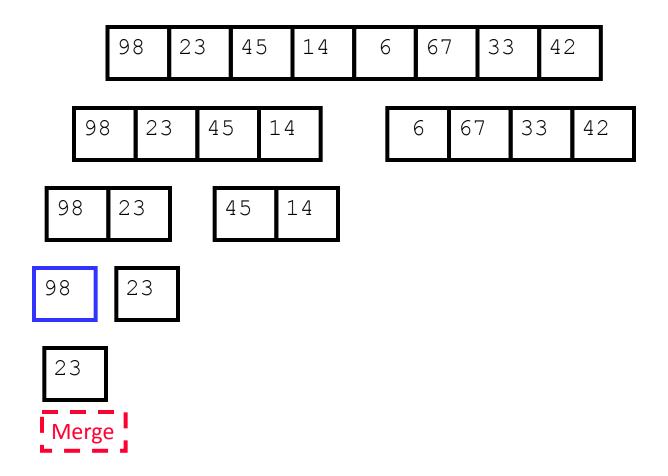


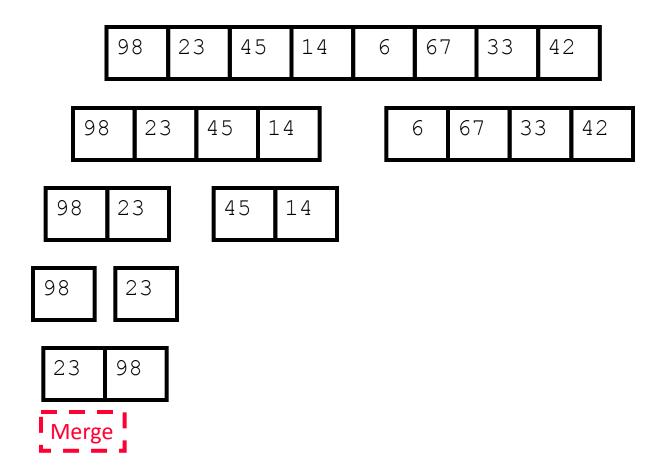


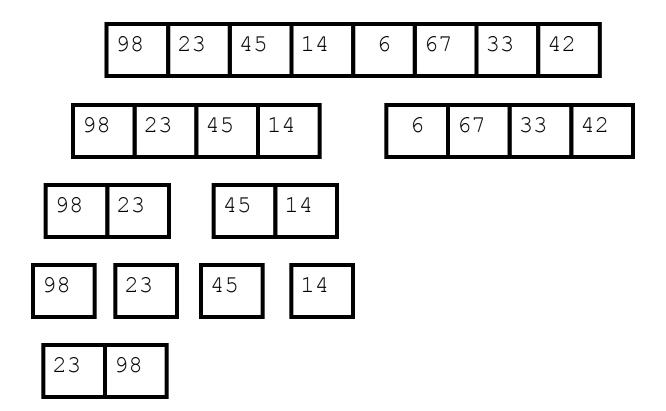


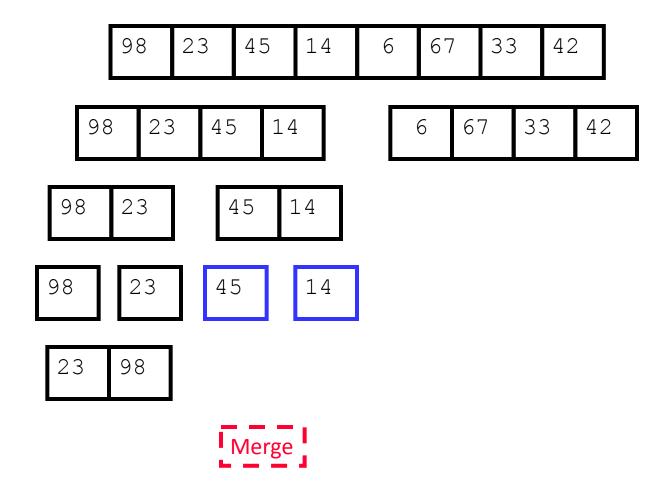


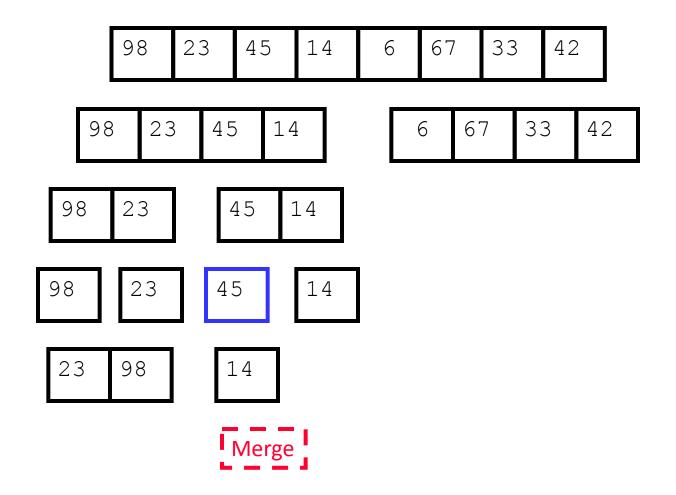


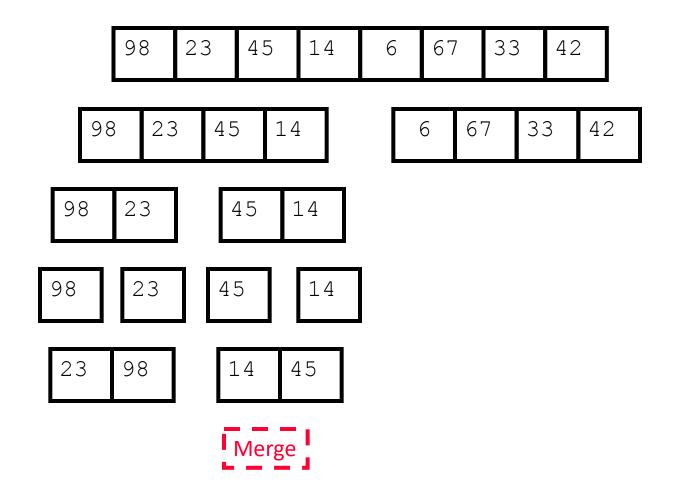


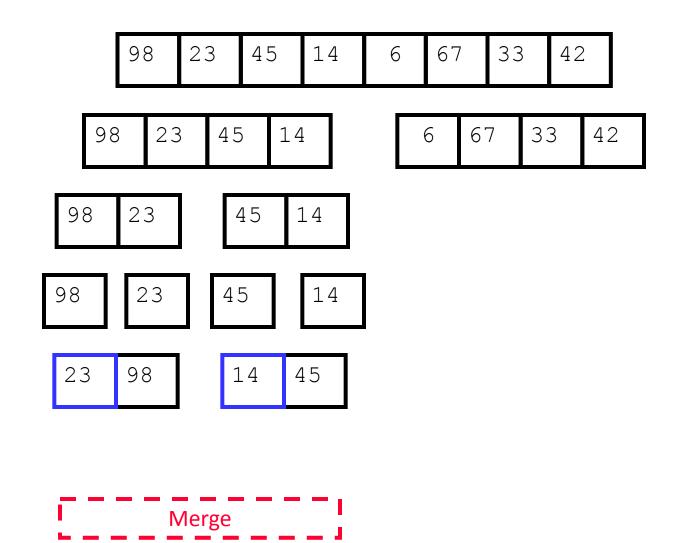


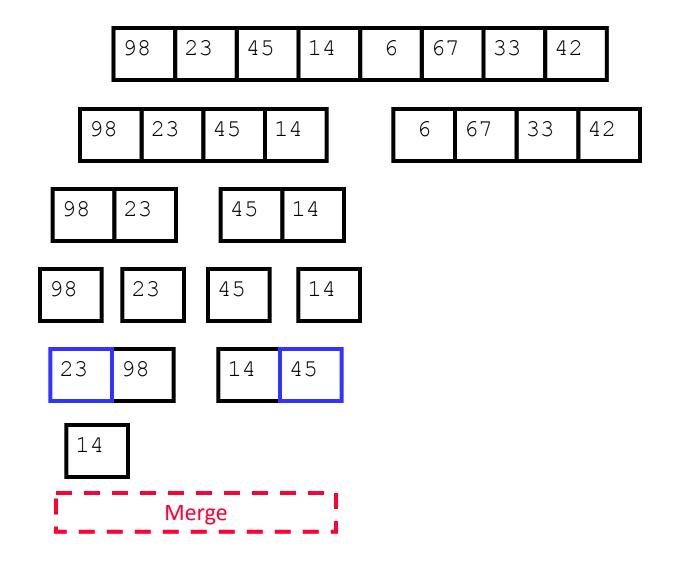


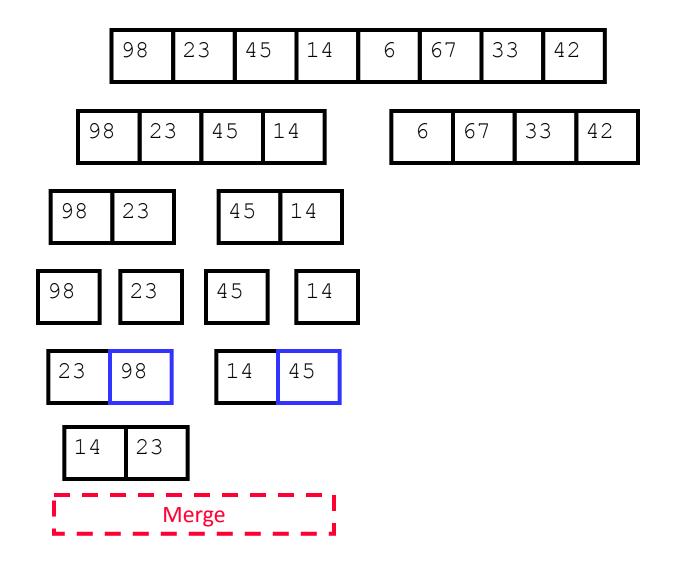


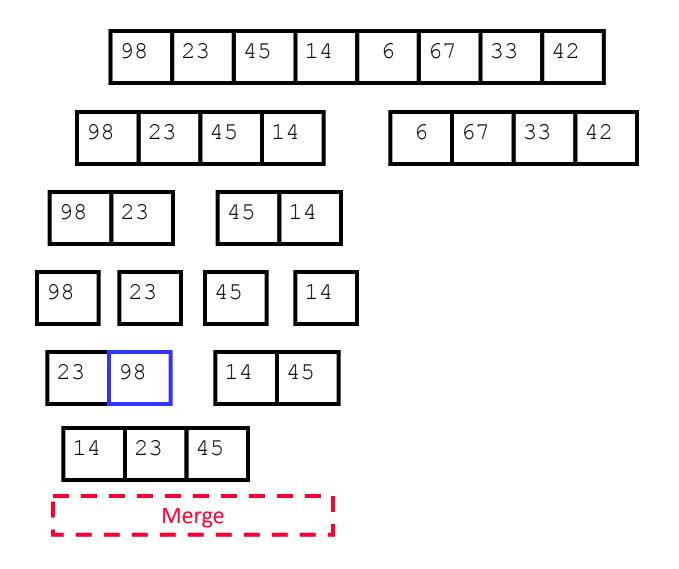


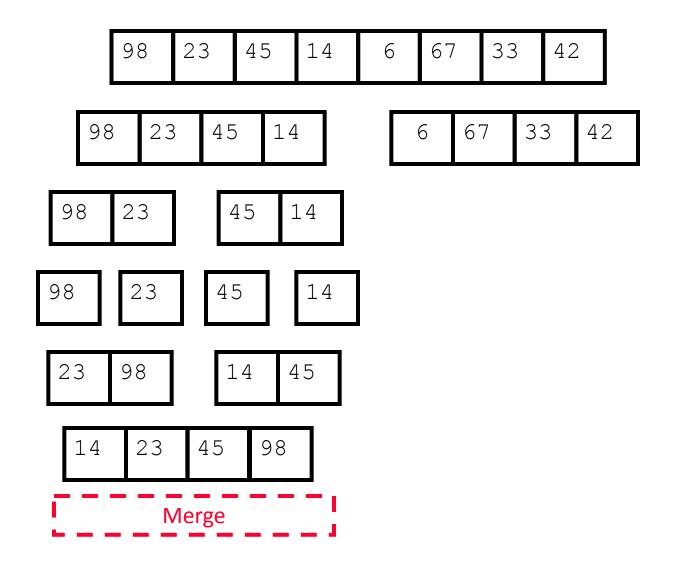


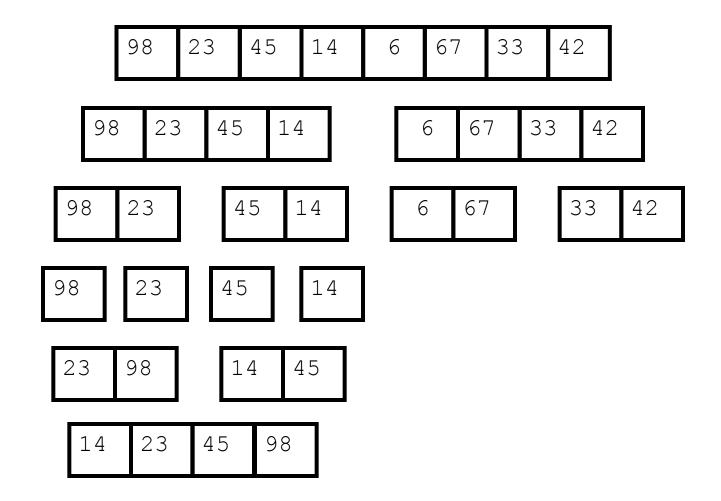


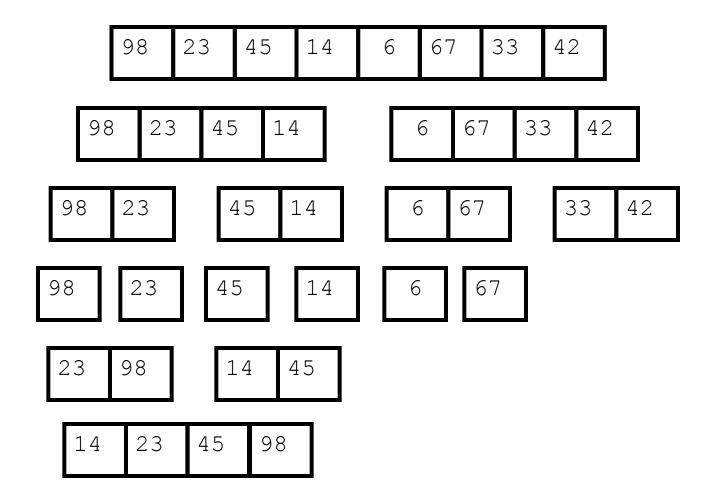


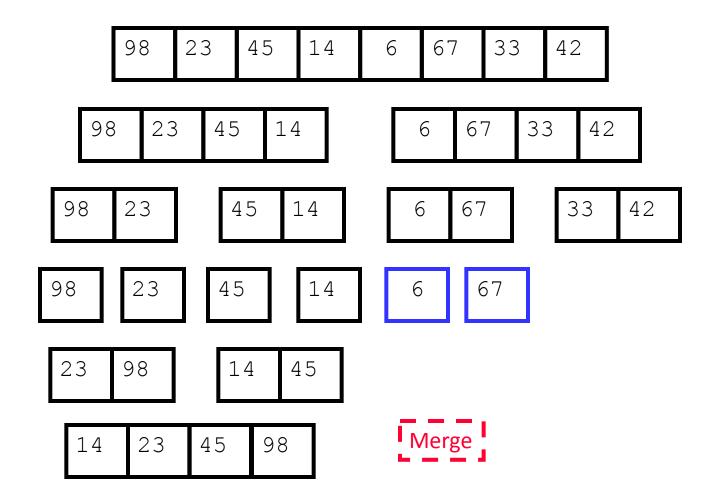


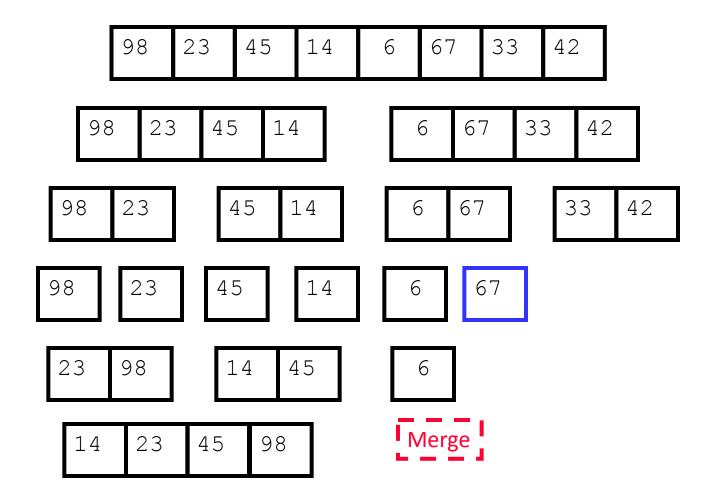


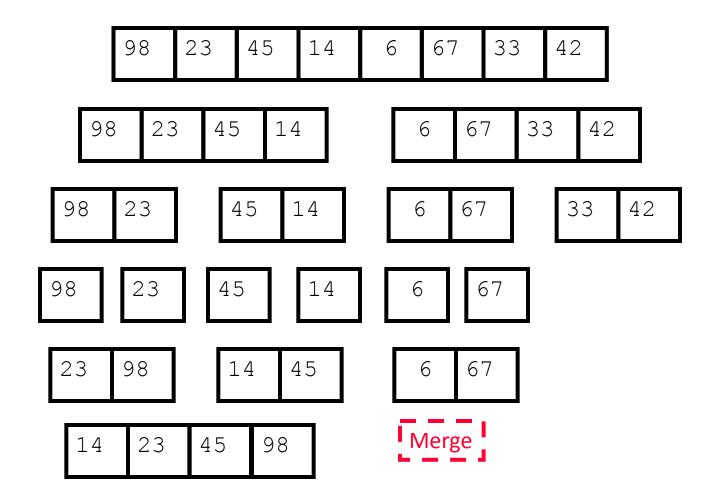


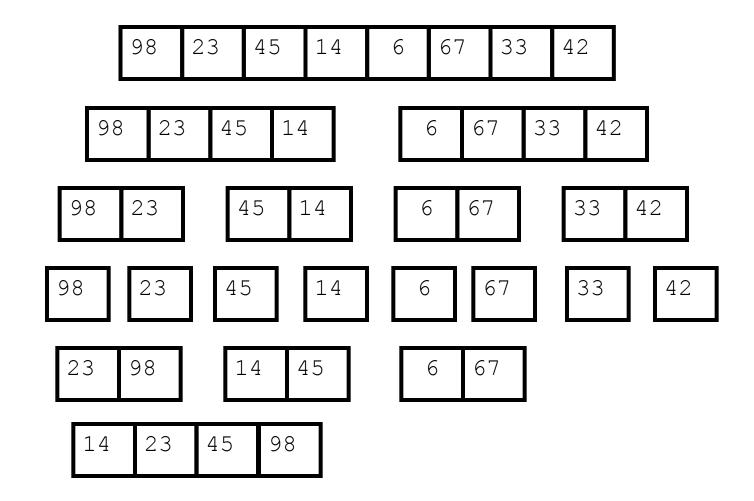








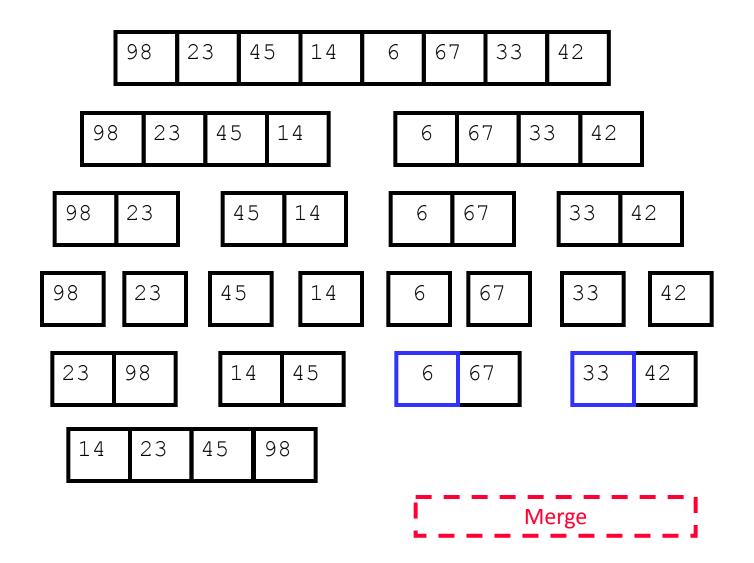


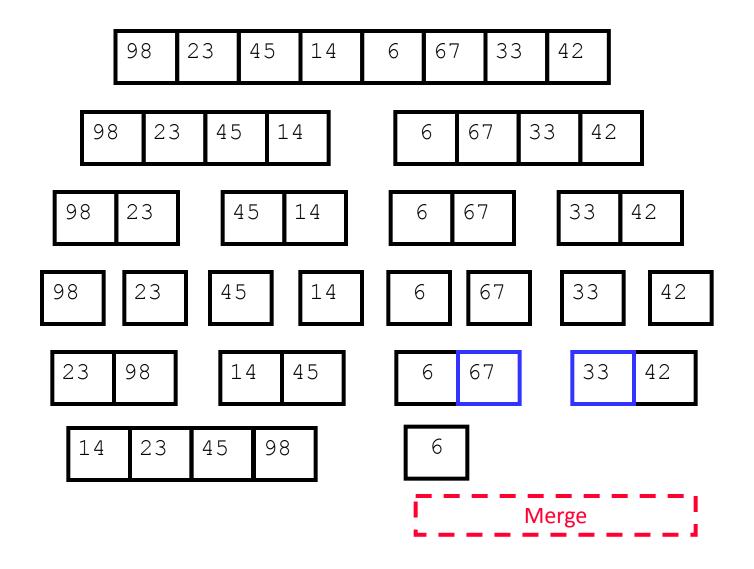


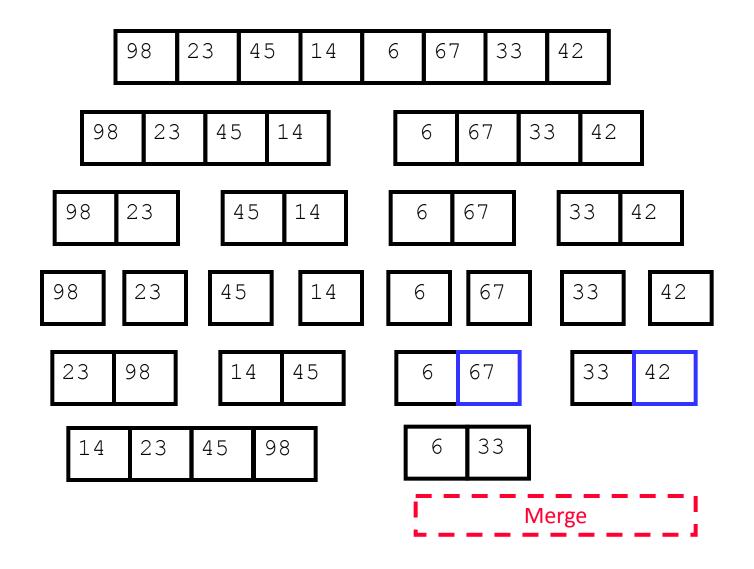


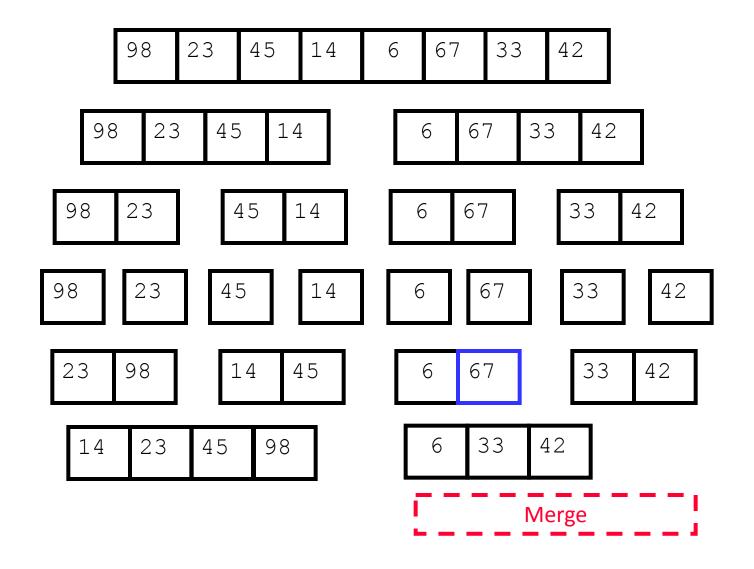


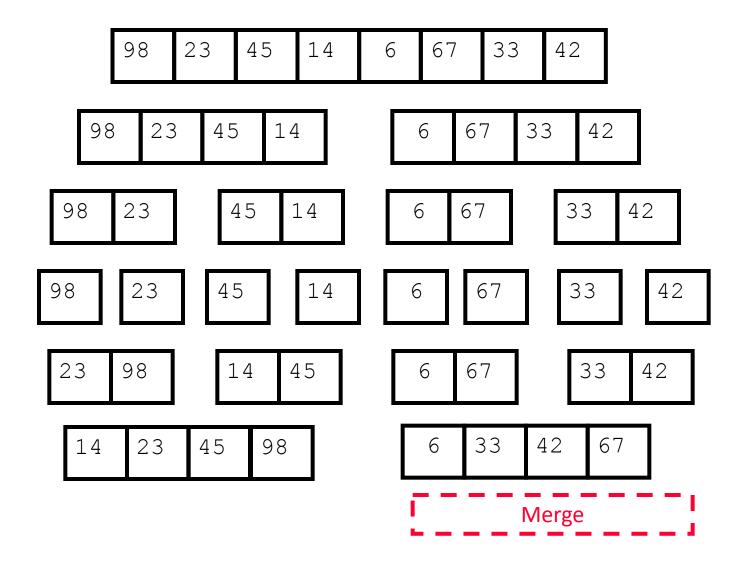


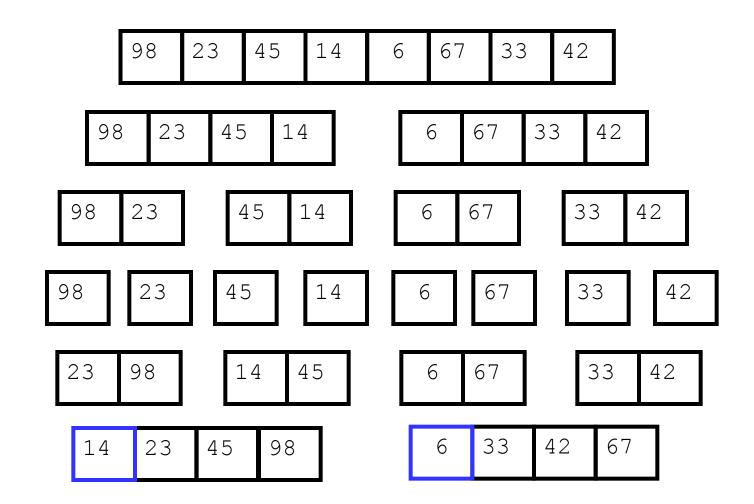


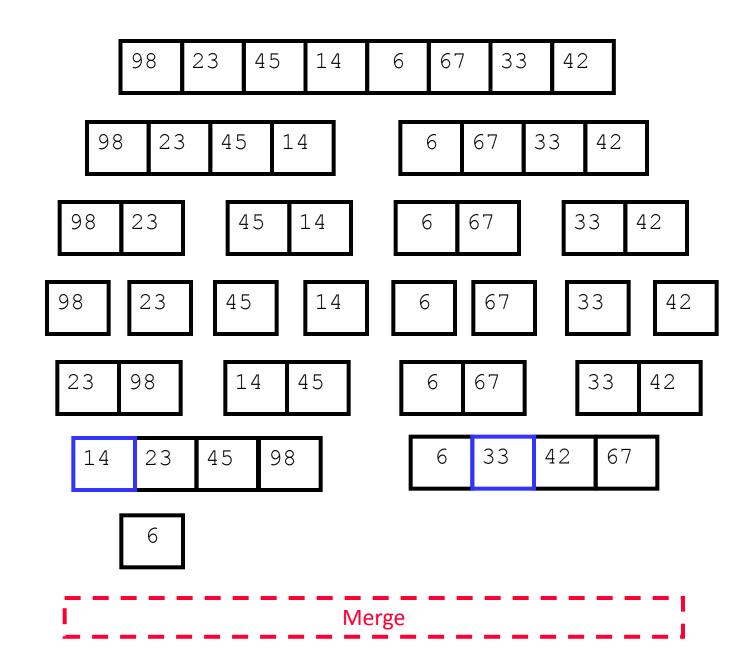


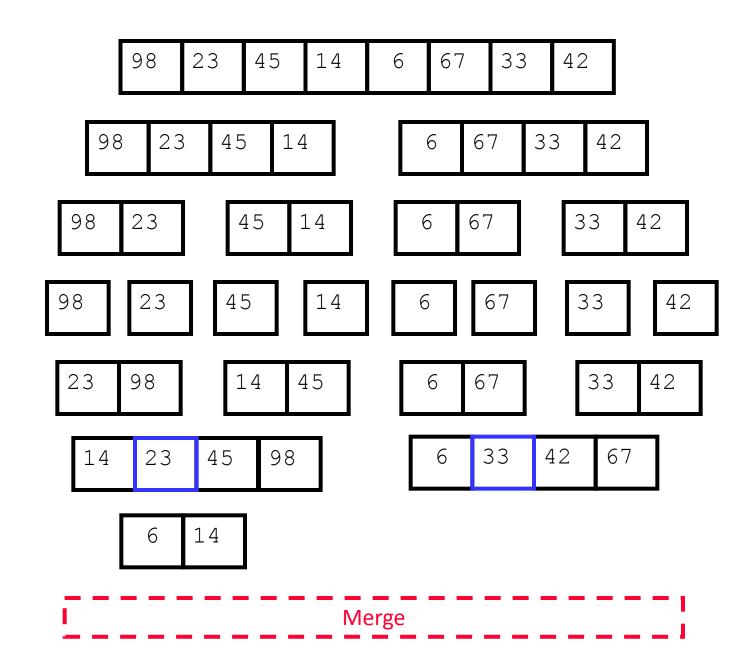


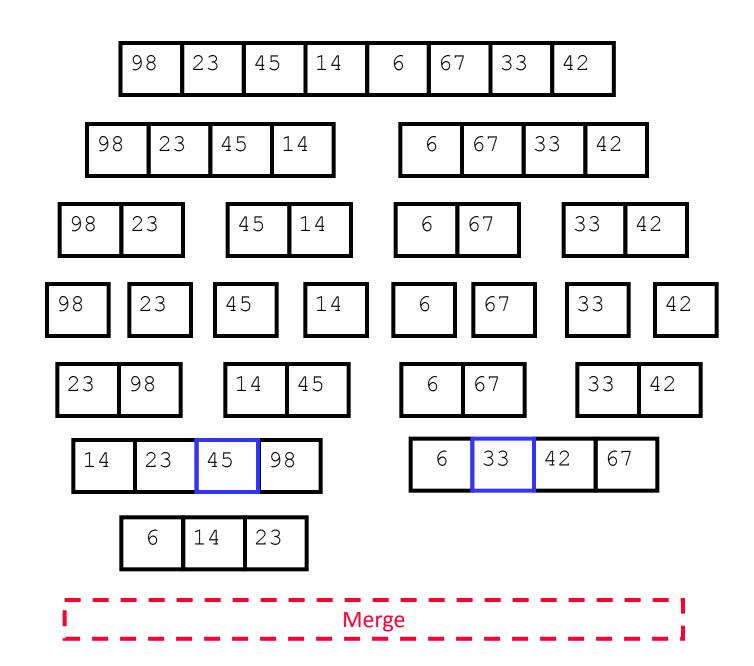


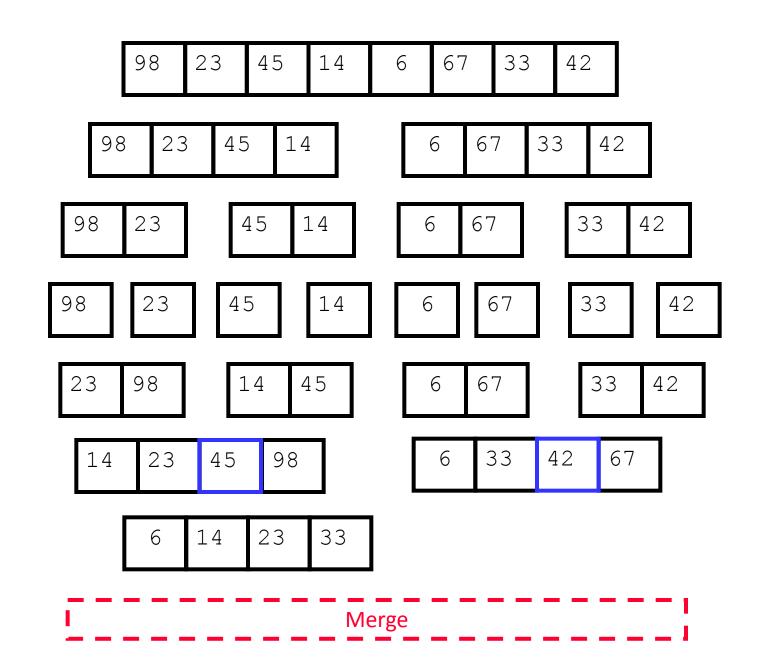


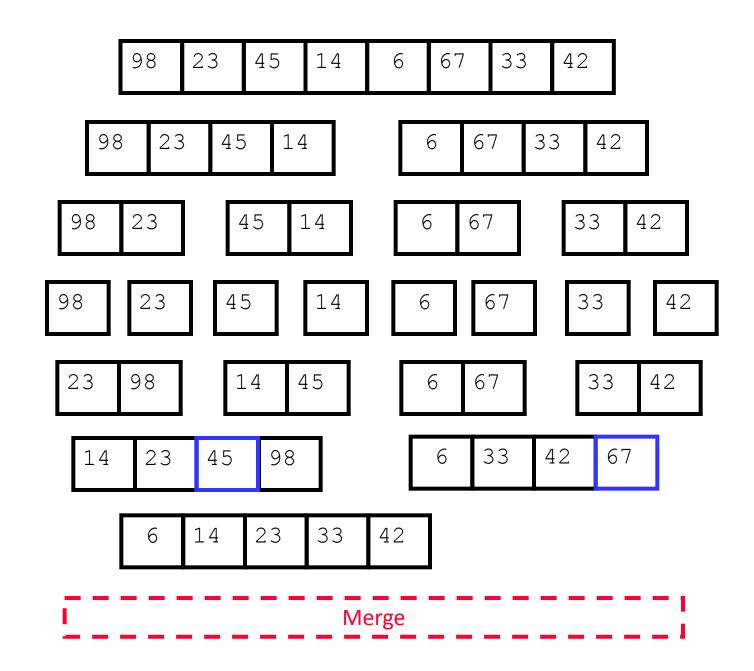


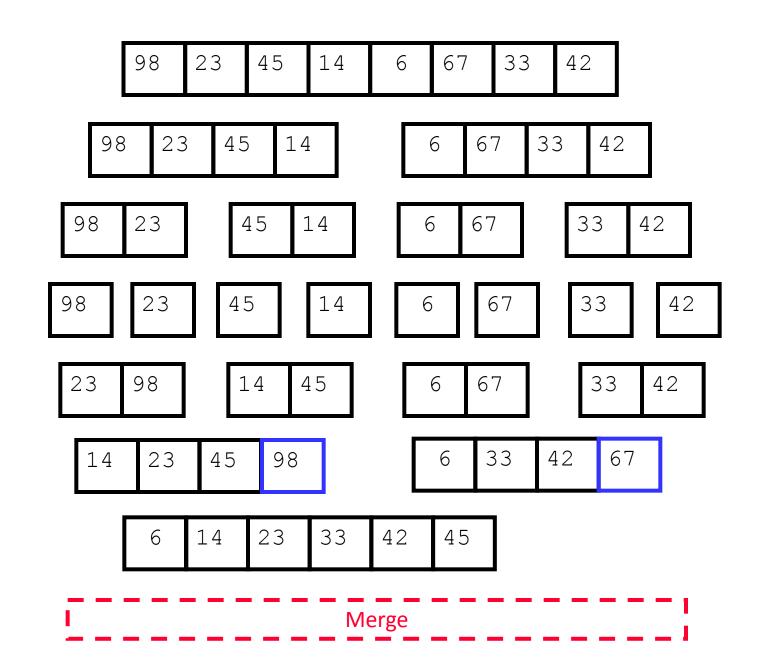


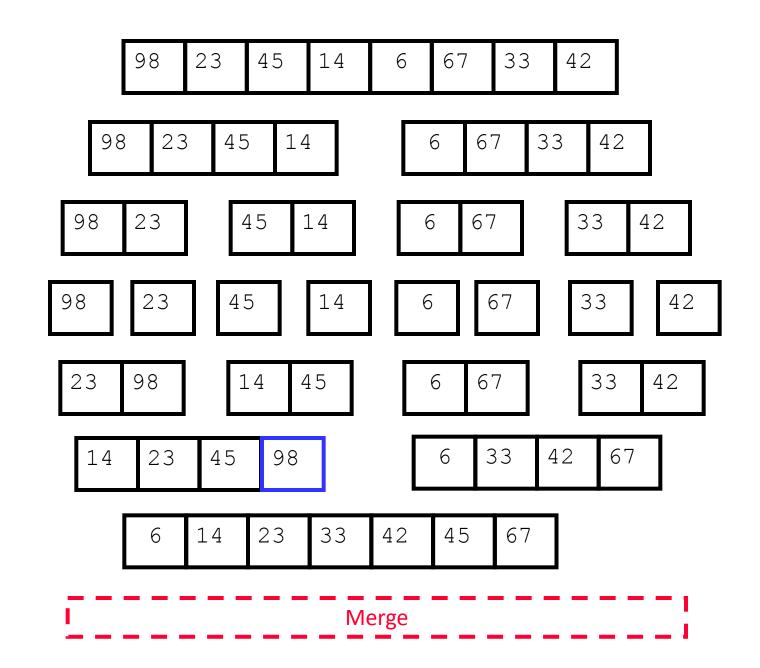


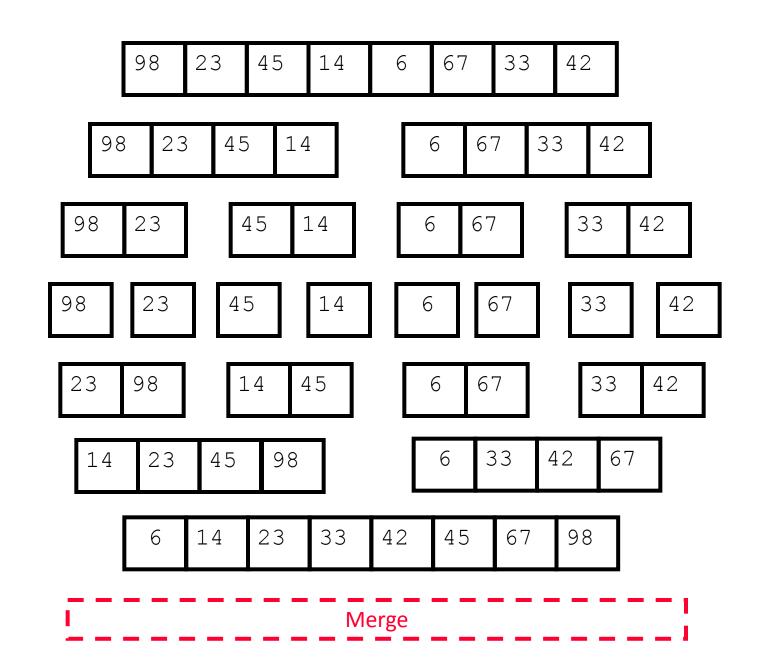


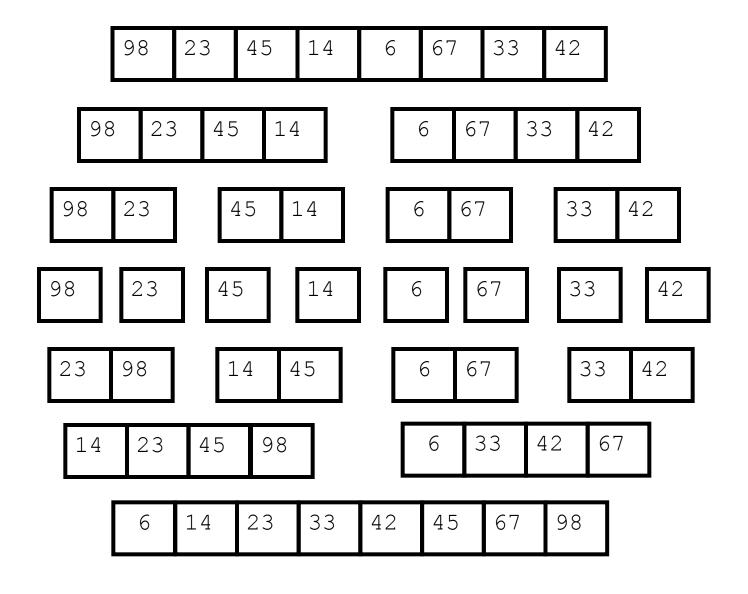


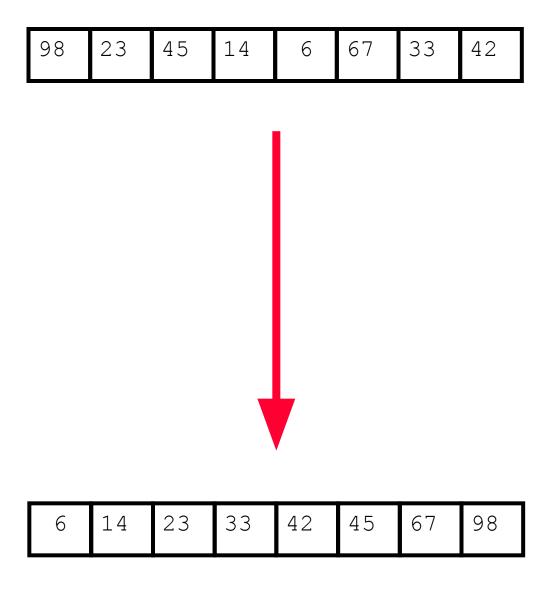












## Textbooks & Web References

- Text Book (Chapter 2)
- Reference book i (Chapter 2)
- Reference book ii (Chapter 5)
- www.visualgo.net

## Thank you & Any question?