

Homework #5  
Introduction to Algorithms/Algorithms 1  
600.363/463  
Spring 2013

**Due on:** Tuesday, March 4th, 5pm

**Late submissions:** will NOT be accepted

**Format:** Please start each problem on a new page.

**Where to submit:** On blackboard, under student assessment

Please type your answers; handwritten assignments will not be accepted.

To get full credit, your answers must be explained clearly,  
with enough details and rigorous proofs.

February 26, 2014

## 1 Problem 1 (20 points)

Suppose we have a complete directed binary tree  $T = (V, E)$  with  $n$  nodes, in which all leaves are at the same depth (a complete binary tree is one in which every node has either two children or none). Then we must have  $n = 2^d - 1$  for some positive integer  $d$ . Suppose further that we have a function  $f : V \mapsto \mathbb{R}$ , and  $f(u) \neq f(v)$  for all  $u, v \in V$  for which  $u \neq v$ . A node  $u \in V$  is *golden* if  $f(u)$  is such that  $f(u) > f(v)$  for all  $v \in V$  for which  $(u, v) \in E$  or  $(v, u) \in E$ . That is, node  $u$  is golden if the value assigned to  $u$  by function  $f$  is strictly greater than the value assigned to the parent of  $u$  (if it exists) and to the children of  $u$  (if they exist). Given any node  $u \in V$ , we can find out the value of  $f(u)$  by *querying*  $u$ . Devise an algorithm that finds a golden node in  $T$  using  $O(\log n)$  queries. Prove the correctness of your algorithm, and prove that it does indeed require  $O(\log n)$  queries.

## 2 Problem 2 (20 points)

Devise an algorithm that takes an undirected graph  $G = (V, E)$  encoded as an adjacency list and returns `True` if graph  $G$  is a tree and `False` if not. Prove the correctness of your algorithm and prove its runtime. Recall that a graph  $G$  is a tree if it is connected and contains no cycles. Your algorithm should run in  $O(|V| + |E|)$  time for full credit. You may assume that the input graph  $G$  does not contain any double edges and does not contain any self-loops— that is, there is at most one edge between any pair of nodes  $u$  and  $v$ , and there are no edges of the form  $\{u, u\}$ .

### Optional exercises

Solve the following problems and exercises from CLRS: 22.2-5, 22.2-8, 22.4-5, 22.5-4, 22-4, 22-1.