

Homework #3  
Introduction to Algorithms/Algorithms 1  
600.363/463  
Spring 2014

**Due on:** Tuesday, Feb 18th, 5pm

**Late submissions:** will NOT be accepted

**Format:** Please start each problem on a new page.

**Where to submit:** On blackboard, under student assessment

Please type your answers; handwritten assignments will not be accepted.

To get full credit, your answers must be explained clearly,  
with enough details and rigorous proofs.

February 11, 2014

## 1 Problem 1 (20 points)

Given an array  $A$  of  $n$  numbers, call index  $i$ , with  $1 \leq i \leq n$ , a *strong* index if for all  $j$  for which  $1 \leq j < i$ , we have  $A[i] > A[j]$ . For example, for the list 1, 5, 2, 3, 6, indices 1, 2 and 5 are strong indices, while indices 3 and 4 are not. A set of  $n$  distinct numbers is randomly permuted into the array  $A[1..n]$ . Show that the expected number of strong indices in array  $A$  is  $O(\log n)$ . Hint: a review of CLRS appendix C and material on the harmonic numbers will be helpful.

## 2 Problem 2 (20 points)

### 2.1 (10 points)

Resolve the following recurrences. Use the master theorem, if applicable. In all examples assume that  $T(1) = 1$ . To simplify your analysis, you may assume that  $n = a^k$  for some  $a, k$ .

1.  $T(n) = 3T(n/2) + 1$

2.  $T(n) = T(n/2) + 2\sqrt{n}$
3.  $T(n) = 16T(n/16) + n^{\frac{3}{2}}$
4.  $T(n) = 28T(n/3) + n^3$
5.  $T(n) = nT(n/2)$
6.  $T(n) = 2T(n-1) + 1$
7.  $T(n) = 8T(n/2) + n^3$
8.  $T(n) = T(n/2) + n \log n$

## 2.2 (10 points)

A sequence  $a_1, a_2, \dots, a_n$  has a *dominant element* if more than half of the elements in the sequence are the same. For example, 3 is a dominant element in the sequence 7, 3, 3, 3, 1, 3, 3, 4, 5, 3. On the other hand, the sequence 5, 4, 1, 1, 2, 3, 2, 3, 6 has no dominant element. Give a divide and conquer algorithm that runs in time  $O(n \log n)$  and finds and returns a dominant element in a sequence of  $n$  numbers or returns `None` if no such element exists. Prove the correctness of your algorithm and prove that its runtime is  $O(n \log n)$ . (Note: there exists an  $O(n)$  algorithm to solve this problem that doesn't make use of divide and conquer— if you figure it out, you may prove its correctness and runtime instead.)

## 3 Optional Exercises

Solve the following problems and exercises from CLRS: 4-3, 4-1, 7-3.