

Quiz #1  
Introduction to Algorithms  
601.433/633  
Spring 2020

You must submit your solutions by Wednesday, April 8th, 10am. Late submissions will NOT be accepted. You may submit handwritten answers – you will have to scan/photograph them, convert it to a pdf and upload it to Gradescope.

**1 Problem 1 (50 points)**

For each statement below explain if it is true or false and in a couple of sentences provide an explanation for your answer. Be as mathematically precise as you can in your explanation. The base of log is 2 unless stated otherwise.

1.  $2^{n^2} = \Theta(3^{n+\sqrt{n}})$
2.  $n! = \omega(2^n)$
3. Let  $f$  be positive function. Then  $f(n) = O((f(n))^2)$ .

## 2 Problem 2 (50 points)

Resolve the following recurrences in terms of a big- $\Theta$  bound. You may assume that  $T(0) = T(1) = 1$ . Provide a proof for the bound you give. If appropriate, you may invoke the Master Theorem (and the appropriate case).

1.  $T(n) = T(n - 1) + 2T(n - 2)$

2.  $T(n) = 7T(n/15) + n^5$

### 3 Problem 3 (50 points)

A sequence  $a_1, a_2, \dots, a_n$  has a special element if more than half of the elements in the sequence are the same. For example, 3 is a special element in the sequence 7, 3, 3, 3, 1, 3, 3, 4, 5, 3. On the other hand, the sequence 5, 4, 1, 1, 2, 3, 2, 3, 6 has no special element. Give a divide and conquer algorithm that runs in time  $O(n \log n)$  and returns a special element in a sequence of  $n$  numbers or returns *None* if no such element exists. Prove the correctness of your algorithm and prove that its runtime is  $O(n \log n)$ . (Note: there exists an  $O(n)$  algorithm to solve this problem that doesn't make use of divide and conquer if you figure it out, you may prove its correctness and runtime instead.)

## 4 Problem 4 (50 points)

You are given  $n$  tasks for a machine. Task  $i$  is described by  $l_i = [a_i, b_i]$  on the real line, where  $a_i, b_i$  are real numbers,  $a_i \leq b_i$  and  $1 \leq i \leq n$ . Give an algorithm that computes the total times of this set of tasks, that is, the length of  $\cup_{i=1}^n l_i$  in  $O(n \log n)$  time.

For example, for the set of tasks  $\{[1, 3], [2, 4.5], [6, 9], [7, 8]\}$ , the total time is  $(4.5 - 1) + (9 - 6) = 6.5$ .

Make sure to prove the correctness and running time of your algorithm.

## 5 Problem 5 (50 points)

Suppose that you have a set of  $n$  integers,  $A = \{a_1, \dots, a_n\}$ , each of them is between 0 and  $K$  (inclusive). Your goal is to find a partition of  $A$  into two sets  $S_1$  and  $S_2$  (so  $S_1 \cup S_2 = A$  and  $S_1 \cap S_2 = \emptyset$ ) that minimizes  $|W(S_1) - W(S_2)|$  where  $W(S)$  denote the sum of integers in  $S$ . Your algorithm's running time should be polynomial in  $n$  and  $K$ .

Make sure to prove the correctness (mainly the optimal substructure property) and running time.